Controlling variability in power systems

Daniel Bienstock

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APAM Nov 17 2017

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A simple example:





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A simple example:



Only one solution:





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A simple example:



Only one solution:



But what if the red node suddenly injects power?

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Only one solution:





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Only one solution:



If red node suddenly injects power, offset using blue node:



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Better example: red node is cheap but unreliable source



Base case solution:



If the red node suddenly injects power, offset using blue node:



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Examples:



Any combination with X + Y = 25 "works" so long as $Y \le 15$.



Any combination with X + Y = 12 "works" so long as Y < -3.

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AC Power Flows

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Real-time:



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AC Power Flows

Real-time:



- Voltage at bus k: $v_k(t) = V_k^{max} \cos(\omega t + \theta_k^V)$
- Current injected at k into km: $i_{km}(t) = I_{km}^{max} \cos(\omega t + \theta_{km}^{I})$.

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AC Power Flows

Real-time:



- Voltage at bus k: $v_k(t) = V_k^{max} \cos(\omega t + \theta_k^V)$
- Current injected at k into km: $i_{km}(t) = I_{km}^{max} \cos(\omega t + \theta_{km}^{l})$.
- Power injected at k into km: $p_{km}(t) = v_k(t)i_{km}(t)$.

Averaged over period T:

$$p_{km} \doteq \frac{1}{T} \int_0^T p(t) = \frac{1}{2} V_k^{max} I_{km}^{max} \cos(\theta_k^V - \theta_{km}^I).$$

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$$V_{k} \doteq \frac{V_{k}^{max}}{\sqrt{2}} e^{j\theta_{k}^{V}}, \quad I_{km} \doteq \frac{I_{km}^{max}}{\sqrt{2}} e^{j\theta_{mk}^{I}}$$

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$$V_k \doteq rac{V_k^{max}}{\sqrt{2}} e^{j heta_k^V}, \quad I_{km} \doteq rac{I_{km}^{max}}{\sqrt{2}} e^{j heta_{mk}^I}$$

$$p_{km} = |V_k||I_{km}|\cos(\theta_k^V - \theta_{km}^I) = \mathcal{R}e(V_k I_{km}^*)$$

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$$p_{km} = |V_k| |I_{km}| \cos(\theta_k^V - \theta_{km}^I) = \mathcal{R}e(V_k I_{km}^*)$$

$$q_{km} \doteq Im(V_{km}I_{km}^*) \quad \text{and} \quad S_{km} \doteq p_{km} + jq_{km}$$

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$$V_{k} \doteq \frac{V_{k}^{max}}{\sqrt{2}} e^{j\theta_{k}^{V}}, \quad I_{km} \doteq \frac{I_{km}^{max}}{\sqrt{2}} e^{j\theta_{mk}^{I}} \text{ (voltage, current)}$$

$$p_{km} = \mathcal{R}e(V_{k}I_{km}^{*}), \quad q_{km} = Im(V_{km}I_{km}^{*})$$

$$(1)$$

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$$V_{k} \doteq \frac{V_{k}^{max}}{\sqrt{2}} e^{j\theta_{k}^{V}}, \quad I_{km} \doteq \frac{I_{km}^{max}}{\sqrt{2}} e^{j\theta_{mk}^{I}} \text{ (voltage, current)}$$

$$p_{km} = \mathcal{R}e(V_{k}I_{km}^{*}), \quad q_{km} = Im(V_{km}I_{km}^{*}) \tag{1}$$

$$I_{km} = \mathbf{y}_{\{\mathbf{k},\mathbf{m}\}}(V_k - V_m),$$

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Controlling variability in power systems

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$$V_{k} \doteq \frac{V_{k}^{max}}{\sqrt{2}} e^{j\theta_{k}^{V}}, \quad I_{km} \doteq \frac{I_{km}^{max}}{\sqrt{2}} e^{j\theta_{mk}^{I}} \text{ (voltage, current)}$$

$$p_{km} = \mathcal{R}e(V_{k}I_{km}^{*}), \quad q_{km} = Im(V_{km}I_{km}^{*}) \tag{1}$$

$$I_{km} = \mathbf{y}_{\{k,m\}}(V_k - V_m), \ \mathbf{y}_{\{k,m\}} = admittance of \ km.$$
(2)

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$$V_{k} \doteq \frac{V_{k}^{max}}{\sqrt{2}} e^{j\theta_{k}^{V}}, \quad I_{km} \doteq \frac{I_{km}^{max}}{\sqrt{2}} e^{j\theta_{mk}^{I}} \text{ (voltage, current)}$$

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(2)

Network Equations



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$$V_{k} \doteq \frac{V_{k}^{max}}{\sqrt{2}} e^{j\theta_{k}^{V}}, \quad I_{km} \doteq \frac{I_{km}^{max}}{\sqrt{2}} e^{j\theta_{mk}^{I}} \text{ (voltage, current)}$$

$$p_{km} = \mathcal{R}e(V_{k}I_{km}^{*}), \quad q_{km} = Im(V_{km}I_{km}^{*})$$

$$(3)$$

$$I_{km} = \mathbf{y}_{\{k,m\}}(V_k - V_m), \ \mathbf{y}_{\{k,m\}} = admittance \text{ of } km.$$
(4)

Network Equations

$$\sum_{km\in\delta(k)}p_{km} = \hat{P}_k, \quad \sum_{km\in\delta(k)}q_{km} = \hat{Q}_k \quad \forall k \qquad (5)$$

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$$V_{k} \doteq \frac{V_{k}^{max}}{\sqrt{2}} e^{j\theta_{k}^{V}}, \quad I_{km} \doteq \frac{I_{km}^{max}}{\sqrt{2}} e^{j\theta_{mk}^{I}} \text{ (voltage, current)}$$

$$p_{km} = \mathcal{R}e(V_{k}I_{km}^{*}), \quad q_{km} = Im(V_{km}I_{km}^{*})$$

$$(3)$$

$$I_{km} = \mathbf{y}_{\{k,m\}}(V_k - V_m), \ \mathbf{y}_{\{k,m\}} = admittance \text{ of } km.$$
(4)

Network Equations

$$\sum_{km\in\delta(k)}p_{km} = \hat{P}_k, \quad \sum_{km\in\delta(k)}q_{km} = \hat{Q}_k \quad \forall k \qquad (5)$$

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Generator: \hat{P}_k , $|V_k|$ given. Other buses: \hat{P}_k , \hat{Q}_k given.

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Optimal power flow (economic dispatch, tertiary control)



- Used periodically to handle the next time window (e.g. 15 minutes, one hour)
- Choose generator outputs
- Minimize cost (quadratic)
- Satisfy demands, meet generator and network constraints
- Constant load (demand) estimates for the time window

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DC-OPF:

s.t.

min
$$c(p)$$
 (convex piecewise-linear or quadratic)

$$B\theta = p - d \tag{6}$$

$$|y_{ij}(\theta_i - \theta_j)| \leq u_{ij}$$
 for each line ij (7)

$$P_g^{min} \leq p_g \leq P_g^{max}$$
 for each bus g (8)

Notation:

 $p = \text{vector of generations} \in \mathcal{R}^n, \quad d = \text{vector of loads} \in \mathcal{R}^n$ $B \in \mathcal{R}^{n \times n}, \quad (\text{bus susceptance matrix})$ $\forall i, j: \quad B_{ij} = \begin{cases} -y_{ij}, & ij \in \mathcal{E} \text{ (set of lines)} \\ \sum_{k; \{k, j\} \in \mathcal{E}} y_{kj}, & i = j \\ 0, & \text{otherwise} \end{cases}$

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Managing changing demands





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What happens when there is a generation/load mismatch



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What happens when there is a generation/load mismatch



Frequency response:

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What happens when there is a generation/load mismatch



Frequency response:

mismatch ΔP

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What happens when there is a generation/load mismatch



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Frequency response:

mismatch $\Delta P \Rightarrow$ frequency change $\Delta \omega \approx -c \Delta P$

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$$H\dot{\omega} = p_m(t) - p_e(t) - D\omega$$

•
$$\omega = \omega(t) =$$
 frequency

- $p_m(t)$ = mechanical power supplied by motor
- $p_e(t)$ = electrical power supplied by motor
- *D* > 0 (tamping)

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Primary frequency control. Handles instantaneous (small) changes.



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Primary frequency control. Handles instantaneous (small) changes. Agent: physics.



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- Primary frequency control. Handles instantaneous (small) changes. Agent: physics.
- 2 Secondary control. Handles changes that span more than a few seconds.

Image: A math a math

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- Primary frequency control. Handles instantaneous (small) changes. Agent: physics.
- Secondary control. Handles changes that span more than a few seconds. Agent: algorithms, pre-set controls.

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Controlling variability in power systems

- Primary frequency control. Handles instantaneous (small) changes. Agent: physics.
- Secondary control. Handles changes that span more than a few seconds. Agent: algorithms, pre-set controls.
- 3 "Tertiary" control: OPF (Optimal power flow). Manages longer lasting changes. Run every few minutes. Goal: economic generation that meets demands while maintaining feasibility (stability).

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- Primary frequency control. Handles instantaneous (small) changes. Agent: physics.
- Secondary control. Handles changes that span more than a few seconds. Agent: algorithms, pre-set controls.
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- Once (?) a day: unit commitment problem. Chooses which generators will operate in the next day or half-day.

Image: A math a math

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- Once (?) a day: unit commitment problem. Chooses which generators will operate in the next day or half-day. Agent: algorithms, humans.

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THE ENERGY CHALLENGE Wind Energy Bumps Into Power Grid's Limits



Mike Groll/Associated Press

The Maple Ridge Wind farm near Lowville, N.Y. It has been forced to shut down when regional electric lines become congested.

By MATTHEW L. WALD Published: August 26, 2008

When the builders of the Maple Ridge Wind farm spent \$320 million to put nearly 200 wind turbines in upstate New York, the idea was to get paid for producing electricity. But at times, regional electric lines have been so congested that Maple Ridge has been forced to shut down even with a brisk wind blowing.

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CIGRE -International Conference on Large High Voltage Electric Systems '09

- Large unexpected fluctuations in wind power can cause additional flows through the transmission system (grid)
- Large power deviations in renewables must be balanced by other sources, which may be far away
- Flow reversals may be observed control difficult
- A solution expand transmission capacity! Difficult (expensive), takes a long time
- Problems already observed when renewable penetration high

Image: A math a math

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CIGRE -International Conference on Large High Voltage Electric Systems '09

- "Fluctuations" 15-minute timespan
- Due to turbulence ("storm cut-off")
- Variation of the same order of magnitude as mean
- Most problematic when renewable penetration starts to exceed 20 30%
- Many countries are getting into this regime

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Controlling variability in power systems

Control model

1 $\bar{w}_i + w_i$ = output of renewable at bus *i*. \bar{w}_i = forecast, w_i = error (uncertain). 2 δ_j = response at bus *j*.

Generic linear control:

$$\delta_{m{j}} = -\sum_{m{i}} \, oldsymbol{lpha_{ji}} \, oldsymbol{w_i}$$

• \mathcal{A} : matrix of all values α_{jj} ; $(\mathcal{A}, \ldots, \mathcal{A}) \in \mathcal{K}$

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Control model

1 $\bar{w}_i + w_i$ = output of renewable at bus *i*. \bar{w}_i = forecast, w_i = error (uncertain). 2 δ_j = response at bus *j*.

Generic linear control:

$$\delta_{m{j}} = -\sum_{m{i}} \, oldsymbol{lpha_{ji}} \, oldsymbol{w_i}$$

A: matrix of all values α_{ji}; (A,..., A) ∈ K
 e.g. ∑_j(1 − α_{ji}) = 0 ∀i for "full-dimensional" uncertainty set.

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Image: A math a math

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Chance-constrained problems (one period)

Optimization Problem

$$\min_{P^g,\mathcal{A}} \sum_k c_k(P^g_k)$$

- s.t. the following system is feasible:
 - \rightarrow Flow balance:



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Safety-constrained problems (one period)

Optimization Problem

$$\min_{P^g,\mathcal{A}} \quad \sum_k c_k(P^g_k)$$

s.t. the following system is feasible:

 \rightarrow Flow balance:

$$B \theta = P^{g} + \overbrace{\bar{w} + w}^{renewables} - \overbrace{\mathcal{A} w}^{linear \ control} - P^{d}$$

$$\rightarrow \text{ Line limits:}$$

$$|\mathbf{E}(f_{km})| + \nu_{km} \operatorname{Std}(f_{km}) \leq f_{km}^{\max} \quad \forall \quad km$$

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$$\min_{P^g,\mathcal{A}} \quad \sum_k c_k(P^g_k)$$

s.t. the following system is feasible:

→ Flow balance:

 $\begin{array}{rcl} & \mathbf{renewables} & \underset{\mathbf{w} \to \mathbf{w}}{\mathsf{linear control}} \\ B \ \theta \ = \ P^g \ + & \overbrace{\mathbf{w} + \mathbf{w}} \ - & \overbrace{\mathcal{A} \mathbf{w}} \ - \ P^d \\ \rightarrow & \mathsf{Line limits,} & \forall \ km \\ & \pi_{km}^T y_{km} | \mathsf{E}(P^g \ + \ \bar{w} \ + \ \mathbf{w} \ - \ \mathcal{A} \ \mathbf{w} \ - \ P^d) | \ + \ \mathbf{\nu}_{km} \, \mathsf{Std}(f_{km}) \le \ f_{km}^{\mathsf{max}}, \\ & \pi_{km}^T \ = \ \text{"shift factors"} \ \rightarrow \ \mathsf{from row-differences of pseudo-inverse of} \ \mathbf{B} \end{array}$

$$\min_{P^g,\mathcal{A}} \quad \sum_k c_k(P^g_k)$$

s.t. the following system is feasible:

→ Flow balance:

 $\begin{array}{rcl} & \mathbf{F} \mathbf{e} = P^{g} + \overbrace{\bar{w} + \mathbf{w}}^{\mathbf{renewables}} & \lim_{\mathcal{A} \to \mathbf{w}}^{\mathbf{linear \ control}} \\ & \rightarrow \mathbf{Line \ limits}, \quad \forall \ km \\ & \pi_{km}^{T} y_{km} |\mathbf{E}(P^{g} + \bar{w} + \mathbf{w} - \mathcal{A} \ \mathbf{w} - P^{d})| + \nu_{km} \operatorname{Std}(f_{km}) \leq f_{km}^{\max}, \\ & \rightarrow \text{ from row-differences of pseudo-inverse of } \mathbf{B} \end{array}$

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$$\min_{P^g,\mathcal{A}} \quad \sum_k c_k(P^g_k)$$

s.t. the following system is feasible:

→ Flow balance:

 $B \theta = P^{g} + \underbrace{\vec{w} + w}_{\vec{w}} - \underbrace{\mathcal{A} w}_{\mathcal{A} w} - P^{d}$ $\rightarrow \text{Line limits, } \forall km$ $\pi_{km}^{T} y_{km} |\mathbf{E}(P^{g} + \bar{w} + w - \mathcal{A} w - P^{d})| + \nu_{km} \operatorname{Std}(f_{km}) \leq f_{km}^{\max},$ $\pi_{km}^{T} = \text{"shift factors"} \rightarrow \text{from row-differences of pseudo-inverse of } B$

Image: A math a math

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$$\min_{P^g,\mathcal{A}} \quad \sum_k c_k(P^g_k)$$

s.t.

 \rightarrow Flow balance:

$$\sum_{j}(1-\alpha_{ji}) = 0 \quad \forall i; \quad \sum_{i}(P_i^g + \bar{w}_i - P_i^d) = 0$$

 $\rightarrow \text{ Line limits, } \forall km$ $\pi_{km}^T y_{km} |\mathbf{E}(P^g + \bar{w} + w - \mathcal{A} w - P^d)| + \nu_{km} \operatorname{Std}(f_{km}) \leq f_{km}^{\max}$

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$$\bullet \mathbf{E}(P^g + \bar{w} + w - \mathcal{A} w - P^d) = P^g + \bar{w} - P^d$$

• Var
$$(f_{ij}) = y_{ij}^2 \pi_{ij}^T (I - \mathcal{A}) \Omega (I - \mathcal{A}^T) \pi_{ij};$$



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$$\bullet \mathbf{E}(P^g + \bar{w} + w - \mathcal{A} w - P^d) = P^g + \bar{w} - P^d$$

$$Var(f_{ij}) = y_{ij}^2 \pi_{ij}^T (I - \mathcal{A}) \Omega (I - \mathcal{A}^T) \pi_{ij};$$

• Yields SOCP formulation for safety-constrained problem.

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$$\bullet \mathbf{E}(P^g + \bar{w} + w - \mathcal{A} w - P^d) = P^g + \bar{w} - P^d$$

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• Yields SOCP formulation for safety-constrained problem.

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Caution!

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$$\bullet \mathbf{E}(P^g + \bar{w} + w - \mathcal{A} w - P^d) = P^g + \bar{w} - P^d$$

• Var
$$(f_{ij}) = y_{ij}^2 \pi_{ij}^T (I - \mathcal{A}) \Omega (I - \mathcal{A}^T) \pi_{ij};$$

- Yields SOCP formulation for safety-constrained problem.
- Caution! SOCP, but not easy in larger cases.
 - Should use sparse formulation.
 - Should use first-order or outer-envelope method.

Image: A mathematical states and a mathem

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Previous work on chance-constrained OPF (review)

Bienstock, Chertkov, Harnett

Roald, Andersson, several coauthors



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Previous work on chance-constrained OPF (review)

- Bienstock, Chertkov, Harnett
- Roald, Andersson, several coauthors
- 1 Chance-constrained DC-OPF with linear control

$$\delta_{m{j}} = -\sum_{i} \; oldsymbol{lpha}_{m{j}m{i}} \; oldsymbol{w}_{m{i}}$$

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can be implemented as convex optimization problem under suitable assumptions

Previous work on chance-constrained OPF (review)

- Bienstock, Chertkov, Harnett
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- 1 Chance-constrained DC-OPF with linear control

$$\delta_{m{j}} = -\sum_{m{i}} \; oldsymbol{lpha}_{m{j}m{i}} \; oldsymbol{w}_{m{i}}$$

can be implemented as convex optimization problem under suitable assumptions

However such convex problems (SOCPs) beyond solvers But first-order methods fast and accurate

An extreme example of variability



Quantity k is large. Bus b has a load of L units.

Stochastic injection at bus **b** = ω, **E**(ω) = μ, Var(ω) = σ².

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An extreme example of variability



- Quantity k is large. Bus b has a load of L units.
- Stochastic injection at bus $\boldsymbol{b} = \boldsymbol{\omega}$, $\mathbf{E}(\boldsymbol{\omega}) = \mu$, $\mathbf{Var}(\boldsymbol{\omega}) = \sigma^2$. $2\sigma > \mu$.

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• Linear generation cost function at $i (0 \le i \le k+1)$: $c_i p_i$.

 $c_0 < c_1 = c_2 \ldots = c_k < c_{k+1}.$

- Safety parameters equal to 3.
- (specify line limits later)



•
$$P_0^g = L - \mu - 3\sigma$$
.
• For $1 \le i \le k$: $\alpha_i = 1/k$ and $P_i^g = 3\sigma/k$.
• $\alpha_{k+1} = P_{k+1}^g = 0$.

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• Stochastic flow on $ab = 3\sigma - \mu - \omega$, with variance σ^2 .

Image: A mathematical states and a mathem

Controlling variability in power systems



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• Stochastic flow on $2h = 3\sigma$, $\mu = 4$, with variance σ^2 .

Stochastic flow on $ab = 3\sigma - \mu - \omega$, with variance σ^2 .

Image: A math and A

Controlling variability in power systems



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Suppose we want to reduce variance on ab by 50%. Then $\alpha_{k+1} = 1 - \sqrt{.5} \approx .293$

Controlling variability in power systems



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• Stochastic flow on $ab = 3\sigma - \mu - \omega$, with variance σ^2 .

• Suppose we want to reduce variance on **ab** by 50%. Then $\alpha_{k+1} = 1 - \sqrt{.5} \approx .293$ sum of line flow variances $> (.5 + (.293)^2(D+1))\sigma^2$.

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Unique optimal safety-constrained solution:

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• $\alpha_{k+1} = P_{k+1}^g = 0$.
• Stochastic flow on $ab = 3\sigma - \mu - \omega$, with variance σ^2 .

• Suppose we want to reduce variance on ab by 50%. Then $\alpha_{k+1} = 1 - \sqrt{.5} \approx .293$ $D = 10, \rightarrow$ sum of line flow variances $\approx 1.35 \sigma^2$.

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$$\min_{P^{g},\mathcal{A}} \sum_{k} c_{k}(P_{k}^{g}) + \Delta(\mathbf{v}^{2})$$
s.t.
$$\sum_{j} (1 - \alpha_{ji}) = 0 \quad \forall i; \quad \sum_{i} (P_{i}^{g} + \bar{w}_{i} - P_{i}^{d}) = 0$$

$$\pi_{km}^{T} y_{km} |\mathbf{E}(f_{km})| + \nu_{km} \operatorname{Std}(f_{km}) \leq f_{km}^{\max} \quad \forall \ km \quad (9)$$

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• $v^2 \doteq$ vector with entries $Var(f_{ij})$. • $\Delta =$ "variance metric".

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• $v^2 \doteq$ vector with entries $Var(f_{ij})$.

• $\Delta =$ "variance metric".

• Special case: $\Delta(Var(f)) = \sum_{ij \in \mathcal{F}} \Delta_{ij}(Var(f_{ij}))$

 Δ_{ij} convex nondecreasing,

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- $\Delta =$ "variance metric".
- Special case: $\Delta(Var(f)) = \sum_{ij \in \mathcal{F}} \Delta_{ij}(Var(f_{ij}))$

 Δ_{ij} convex nondecreasing, but \mathcal{F} could depend on solution

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1 \mathcal{F} = all lines. Var-aware SCOPF is a convex optimization problem.



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- **1** \mathcal{F} = all lines. Var-aware SCOPF is a convex optimization problem.
- **2** \mathcal{F} = set of **N** lines with largest flow variance, **N** fixed. E.g. **N** = 50.



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Image: A math a math

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- 4 Logarithmic barrier function.

$$\Delta_{ij} = -\rho \log(f_{km}^{\max} - |\mathbf{E}(f_{km})| - \boldsymbol{\nu_{km}} \mathbf{Std}(f_{km}))$$

where $\rho > 0$

Image: A math a math

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$$\min_{P^{g},\mathcal{A}} \sum_{k} c_{k}(P_{k}^{g}) + \sum_{ij \in \mathcal{F}} \Delta_{ij}(\operatorname{Var}(f_{ij}))$$
s.t.
$$\sum_{j} (1 - \alpha_{ji}) = 0 \quad \forall i; \quad \sum_{i} (P_{i}^{g} + \bar{w}_{i} - P_{i}^{d}) = 0$$

$$\pi_{km}^{T} y_{km} |\mathbf{E}(f_{km})| + \nu_{km} \operatorname{Std}(f_{km}) \leq f_{km}^{\max} \quad \forall \quad km (10)$$

$$\min_{P^g,\mathcal{A}} \sum_{k} c_k(P^g_k) - \rho \log(f_{km}^{\max} - |\mathbf{E}(f_{km})| - \nu_{km} \operatorname{Std}(f_{km}))$$

s.t.
$$\sum_{j} (1 - \alpha_{ji}) = 0 \quad \forall i; \quad \sum_{i} (P^g_i + \bar{w}_i - P^d_i) = 0$$

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$$\min_{P^g, \mathcal{A}} \sum_{k} c_k(P^g_k) + \Delta(\mathbf{v}^2)$$

s.t. $\sum_{j} (1 - \alpha_{ji}) = 0 \quad \forall i; \quad \sum_{i} (P^g_i + \bar{w}_i - P^d_i) = 0$
 $\pi^T_{km} y_{km} |\mathbf{E}(f_{km})| + \nu_{km} \operatorname{Std}(f_{km}) \leq f^{\max}_{km} \quad \forall \ km$

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Correction vs Formal Variance-Aware Optimization

Several groups of authors:

- Outer-approximation algorithm for chance-constrained DC-OPF converges in **few** iterations.
- 2 Only a few lines are "at risk" in realistic cases.

Correction vs Formal Variance-Aware Optimization

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- **3** Low-hanging fruit:

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Correction vs Formal Variance-Aware Optimization

Several groups of authors:

- Outer-approximation algorithm for chance-constrained DC-OPF converges in **few** iterations.
- 2 Only a few lines are "at risk" in realistic cases.
- **I Low-hanging fruit:** instead of variance-aware optimization, first solve SC-OPF problem, and then **correct** or **adjust** to reduce variability without increasing cost (by much).

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GENERIC CORRECTION TEMPLATE

Input: an instance of the safety-constrained problem and a variance metric.

Step I. Solve safety-constrained OPF problem, with solution $(\bar{P}^g, \bar{\mathcal{A}})$.

Step II. Perform a small number of adjustment iterations which shift (\bar{P}^g, \bar{A}) to a new feasible solution that attains an improved value of the variance metric, while at the same time increasing generation cost in a moderate manner.



• $\hat{\mathcal{A}}$: a given control matrix



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Reroute $(au, \hat{\mathcal{A}})$

• $\hat{\mathcal{A}}$: a given control matrix

- So $\operatorname{Var}(f_{km}) = y_{km}^2 \pi_{km}^T (I \hat{\mathcal{A}}) \Omega (I \hat{\mathcal{A}}^T) \pi_{km} \doteq \hat{\mathcal{V}}_{km}$ is fixed for every km
- **0** < *τ* < **1**.



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Reroute $(au, \hat{\mathcal{A}})$

• $\hat{\mathcal{A}}$: a given control matrix

So
$$\operatorname{Var}(f_{km}) = y_{km}^2 \pi_{km}^T (I - \hat{\mathcal{A}}) \Omega (I - \hat{\mathcal{A}}^T) \pi_{km} \doteq \hat{V}_{km}$$

is fixed for every km

 $\bullet \ 0 < \tau < 1.$

$$\min_{P^g} \sum_k c_k(P^g_k)$$

s.t.
$$\sum_{i} (P_i^g + \bar{w}_i - P_i^d) = 0$$

 $\pi_{km}^{T} y_{km} |\mathbf{E}(f_{km})| + \nu_{km} \sqrt{\hat{\mathbf{V}}_{km}} \leq (1 - \tau) f_{km}^{\max} \quad \forall \ km$

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Here, f_{km} is an implicit function of P^g and \hat{A}_{a}

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• \vec{P}^g : generation vector. \vec{f} : corresponding flow vector.

$$\min_{\mathcal{A}} \sum_{km \in \mathcal{F}(\bar{f})} \Delta_{km} \left(y_{km}^2 \pi_{km}^{\mathcal{T}} (I - \mathcal{A}) \Omega (I - \mathcal{A}^{\mathcal{T}}) \pi_{km} \right)$$
constraints on $\mathcal{A}: \sum_{j} (1 - \alpha_{jj}) = 0 \quad \forall i$

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 $\rightarrow\,$ Keep power flows fixed, improve on control

Correction procedure

Input: Feasible solution $(P^{g,0}, \mathcal{A}_0), \quad 0 < \tau < 1.$

For k = 1, 2, ..., K perform iteration k:

- 1. Run **Reroute** $(\mathcal{A}_{k-1}, \tau)$. If infeasible STOP. Else let $P^{g,k}$ be optimal.
- **2.** Solve *VShift*($P^{g,k}$), with solution $\hat{\mathcal{A}}_k$.
- 3. Set $A_k \leftarrow (1 \lambda)A_{k-1} + \lambda \hat{A}_k$. $0 < \lambda < 1$ chosen so that $P^{g,k}, A_k$ feasible

4. If $\Delta(\mathcal{A}_k) \geq \Delta(\mathcal{A}_{k-1})$, STOP.

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Thm.: In convex case if we stop in Step 4, then Δ minimum.

Image: Image:

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- 2746 buses, 3514 branches, 520 generators
- 22 stochastic injection sites, sum of mean injections 4611.57 (penetration 18.5% penetration)
- average ratio of standard deviation to mean $\approx 30\%$.

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 - The 100 lines with largest flow magnitude
 - Lines ij where $|\bar{f}_{ij}| + \nu_{ij} \mathsf{Std}(f_{ij}) \ge (1 \tau) f_{ij}^{\max}$

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• Initial variance metric: $\approx 6.3 \times 10^{04}$.



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- Initial variance metric: $\approx 6.3 \times 10^{04}$.
- After one iteration: $\approx 4.65 \times 10^{04}$.



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- Initial variance metric: $\approx 6.3 \times 10^{04}$.
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- Cost nearly constant after two iterations

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- Initial variance metric: $\approx 6.3 \times 10^{04}$.
- After one iteration: $\approx 4.65 \times 10^{04}$.
- After two iterations: $\approx 4.50 \times 10^{04}$. Approx. 40% reduction relative to original.
- Cost nearly constant after two iterations
- \blacksquare Variance-shifting SOCP has approximately 1.4×10^{05} variables and constraints and 1×10^{06} nonzeros

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Solution times of a few seconds

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