Chance-constrained optimization problems in the operation of the power grid

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Michigan 2013

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Part I: Robust Optimal Power Flow with Uncertain Renewables

with Michael Chertkov (LANL) and Sean Harnett (Columbia)

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THE ENERGY CHALLENGE Wind Energy Bumps Into Power Grid's Limits



Mike Groll/Associated Press

The Maple Ridge Wind farm near Lowville, N.Y. It has been forced to shut down when regional electric lines become congested.

By MATTHEW L. WALD Published: August 26, 2008

When the builders of the Maple Ridge Wind farm spent \$320 million to put nearly 200 wind turbines in upstate New York, the idea was to get paid for producing electricity. But at times, regional electric lines have been so congested that Maple Ridge has been forced to shut down even with a brisk wind blowing.

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Optimal power flow (economic dispatch, tertiary control)



- Used periodically to handle the next time window (e.g. 15 minutes, one hour)
- Choose generator outputs
- Minimize cost (quadratic)
- Satisfy demands, meet generator and network constraints
- Constant load (demand) estimates for the time window

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OPF:

s.t.

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min c(p) (a quadratic) $B\theta = p - d$ (1)

$$|y_{ij}(\theta_i - \theta_j)| \leq u_{ij}$$
 for each line ij (2)

$$P_g^{min} \leq p_g \leq P_g^{max}$$
 for each bus g (3)

Notation:

 $p = \text{vector of generations} \in \mathcal{R}^n, \quad d = \text{vector of loads} \in \mathcal{R}^n$ $B \in \mathcal{R}^{n \times n}, \quad (\text{bus susceptance matrix})$ $\forall i, j: \quad B_{ij} = \begin{cases} -y_{ij}, & ij \in \mathcal{E} \text{ (set of lines)} \\ \sum_{k; \{k, j\} \in \mathcal{E}} y_{kj}, & i = j \\ 0, & \text{otherwise} \end{cases}$

min c(p) (a quadratic)

s.t.

$$\begin{array}{lll} \mathcal{B}\theta = p - d \\ |y_{ij}(\theta_i - \theta_j)| &\leq u_{ij} & \text{for each line } ij \\ P_g^{min} &\leq p_g &\leq P_g^{max} & \text{for each bus } g \end{array}$$

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min c(p) (a quadratic)
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s.t.
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 $\begin{array}{lll} B\theta = p - d \\ |y_{ij}(\theta_i - \theta_j)| &\leq u_{ij} \quad \text{for each line } ij \\ P_g^{min} &\leq p_g &\leq P_g^{max} \quad \text{for each bus } g \end{array}$

How does OPF handle short-term fluctuations in **demand** (d)? **Frequency control:**

- Automatic control: primary, secondary
- The output of special generators varies up or down proportionally to aggregate change

How does OPF handle short-term fluctuations in renewable output? **Answer:** Same mechanism, now used to handle aggregate wind power change

CIGRE -International Conference on Large High Voltage Electric Systems '09

- Large unexpected fluctuations in wind power can cause additional flows through the transmission system (grid)
- Large power deviations in renewables must be balanced by other sources, which may be far away
- Flow reversals may be observed control difficult
- A solution expand transmission capacity! Difficult (expensive), takes a long time
- Problems already observed when renewable penetration high

Image: A math a math

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CIGRE -International Conference on Large High Voltage Electric Systems '09

- "Fluctuations" 15-minute timespan
- Due to turbulence ("storm cut-off")
- Variation of the same order of magnitude as mean
- Most problematic when renewable penetration starts to exceed 20 30%

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Many countries are getting into this regime

Need to model variation in wind power between dispatches Wind at farm attached to bus *i* of the form $\mu_i + \mathbf{w_i}$ – Weibull distribution?



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Wind model

From CIGRE report, aggregated over Germany:



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Experiment

Bonneville Power Administration data, Northwest US

- data on wind fluctuations at planned farms
- with standard OPF, 7 lines exceed limit \geq 8% of the time



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If power flow in a line exceeds its limit, the line becomes compromised and may 'trip'. But process is complex and time-averaged:

- Thermal limit is most common
- Thermal limit may be in terms of terminal equipment, not line itself
- Wind strength and wind direction contributes to line temperature
- In medium-length lines (\sim 100 miles) the line limit is due to voltage drop, not thermal reasons
- In long lines, it is due to phase angle change (stability), not thermal reasons

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In 2003 U.S. blackout event, many critical lines tripped due to thermal reasons, but well short of their line limit summary: exceeding limit for too long is bad, but complicated

Image: A mathematical states and a mathem

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want: "fraction time a line exceeds its limit is small"

proxy: prob(violation on line i) < ϵ for each line i

Control

For each generator i, two parameters:

- $\overline{p_i}$ = mean output
- α_i = response parameter

Real-time output of generator *i*:

$$p_i = \overline{p}_i - \alpha_i \sum_j \Delta \omega_j$$

where $\Delta \omega_j$ = change in output of renewable *j* (from mean).

$$\sum_{i} \alpha_{i} = 1$$

$$\sim$$
 primary + secondary control

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wind power at bus *i*: $\mu_i + \mathbf{w}_i$

DC approximation

$$B\theta = \overline{p} - d$$

+(\mu + \mu - \alpha \sum_{i \in G} \mu_i)
$$\theta = B^+(\overline{p} - d + \mu) + B^+(I - \alpha e^T) \mu$$

flow is a linear combination of bus power injections:

$$\mathbf{f_{ij}} = y_{ij}(\boldsymbol{\theta}_i - \boldsymbol{\theta}_j)$$

Image: A mathematical states and a mathem

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Computing line flows

$$\mathbf{f}_{ij} = y_{ij} \left((B_i^+ - B_j^+)^T (\bar{p} - d + \mu) + (A_i - A_j)^T \mathbf{w} \right),$$
$$A = B^+ (I - \alpha e^T)$$

Given distribution of wind can calculate moments of line flows:

•
$$Ef_{ij} = y_{ij}(B_i^+ - B_j^+)^T(\bar{p} - d + \mu)$$

•
$$var(\mathbf{f}_{ij}) := s_{ij}^2 \ge y_{ij}^2 \sum_k (A_{ik} - A_{jk})^2 \sigma_k^2$$

(assuming independence)

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Chance constraints to deterministic constraints

• chance constraint: $P(\mathbf{f}_{ij} > f_{ij}^{max}) < \epsilon_{ij}$ and $P(\mathbf{f}_{ij} < -f_{ij}^{max}) < \epsilon_{ij}$

 from moments of f_{ij}, can get conservative approximations using e.g. Chebyshev's inequality

Image: A math a math

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Chance constraints to deterministic constraints

• chance constraint: $P(\mathbf{f}_{ij} > f_{ij}^{max}) < \epsilon_{ij}$ and $P(\mathbf{f}_{ij} < -f_{ij}^{max}) < \epsilon_{ij}$

- from moments of f_{ij}, can get conservative approximations using e.g. Chebyshev's inequality
- $\hfill for Gaussian wind, can do better, since <math display="inline">f_{ij}$ is Gaussian :

$$|E\mathbf{f}_{\mathbf{ij}}| + var(\mathbf{f}_{\mathbf{ij}})\phi^{-1}(1-\epsilon_{ij}) \leq f_{ij}^{max}$$

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Formulation:

Choose mean generator outputs and control to minimize expected cost, with the probability of line overloads kept small.

$$\begin{split} \min_{\overline{p},\alpha} \mathbb{E}[c(\overline{p})] \\ \text{s.t.} \sum_{i \in G} \alpha_i &= 1, \ \alpha \ge 0 \\ B\delta &= \alpha, \delta_n = 0 \\ \sum_{i \in G} \overline{p}_i + \sum_{i \in W} \mu_i &= \sum_{i \in D} d_i \\ \overline{f}_{ij} &= y_{ij}(\overline{\theta}_i - \overline{\theta}_j), \\ B\overline{\theta} &= \overline{p} + \mu - d, \ \overline{\theta}_n = 0 \\ s_{ij}^2 &\geq y_{ij}^2 \sum_{k \in W} \sigma_k^2 (B_{ik}^+ - B_{jk}^+ - \delta_i + \delta_j)^2 \\ |\overline{f}_{ij}| &+ s_{ij} \phi^{-1} (1 - \epsilon_{ij}) \le f_{ij}^{max} \end{split}$$

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Data errors?

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$$\begin{split} s_{ij}^2 &\geq y_{ij}^2 \sum_{k \in W} \sigma_k^2 (B_{ik}^+ - B_{jk}^+ - \delta_i + \delta_j)^2 \\ |\overline{f}_{ij}| &+ s_{ij} \phi^{-1} \left(1 - \epsilon_{ij}\right) \leq f_{ij}^{max} \end{split}$$

(the \overline{f}_{ij} implicitly incorporate the μ_i)

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Data errors?

$$\begin{split} s_{ij}^2 &\geq y_{ij}^2 \sum_{k \in W} \sigma_k^2 (B_{ik}^+ - B_{jk}^+ - \delta_i + \delta_j)^2 \\ |\overline{f}_{ij}| &+ s_{ij} \phi^{-1} \left(1 - \epsilon_{ij}\right) \leq f_{ij}^{max} \end{split}$$

(the \overline{f}_{ij} implicitly incorporate the μ_i) What if the μ_i or the σ_k are incorrect? ... What happens to

 $Prob(\mathbf{f_{ij}} > u_{ij})?$

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Theorem: Suppose there are parameters M > 0, V > 0 such that

$$|\bar{\mu}_i - \mu_i| < M\mu_i$$
 and $|\bar{\sigma}_i^2 - \sigma_i| < V\sigma_i$

for all *i*. Then:

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Theorem: Suppose there are parameters M > 0, V > 0 such that

$$|\bar{\mu}_i - \mu_i| < M\mu_i$$
 and $|\bar{\sigma}_i^2 - \sigma_i| < V\sigma_i$

for all *i*. Then:

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$$Prob(f_{ij} > f_{ij}^{max}) < \epsilon_{ij} + O(V) + O(M)$$

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Theorem: Suppose there are parameters M > 0, V > 0 such that

$$|\bar{\mu}_i - \mu_i| < M\mu_i$$
 and $|\bar{\sigma}_i^2 - \sigma_i| < V\sigma_i$

for all *i*. Then:

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$$Prob(f_{ij} > f_{ij}^{max}) < \epsilon_{ij} + O(V) + O(M)$$

Here, the O() "hides" some constants dependent on e.g. reactances

Image: A math a math

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Can we guarantee that $Prob(f_{ij} > f_{ij}^{max})$ is small even under data errors?

Polyhedral data error model:

$$|\tilde{\sigma}_i^2 - \sigma_i^2| \le \gamma_i \ \forall i, \ \sum_i \frac{|\tilde{\sigma}_i^2 - \sigma_i^2|}{\gamma_i} \le \Gamma.$$

Ellipsoidal data error model:

$$(\tilde{\sigma}^2 - \sigma^2)^T A(\tilde{\sigma}^2 - \sigma^2) \leq b$$

Here $A \succeq 0$ and b > 0 are parameters.

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chance constraints

Nominal case:

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Nominal case: $|E \mathbf{f}_{ij}| + var(\mathbf{f}_{ij})\phi^{-1}(1 - \epsilon_{ij}) \leq f_{ij}^{max}$

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Nominal case: $|E \mathbf{f}_{ij}| + var(\mathbf{f}_{ij})\phi^{-1}(1 - \epsilon_{ij}) \leq f_{ij}^{max}$

 \rightarrow a conic constraint

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Nominal case: $|E \mathbf{f}_{ij}| + var(\mathbf{f}_{ij})\phi^{-1}(1 - \epsilon_{ij}) \le f_{ij}^{max}$ \rightarrow a conic constraint

Robust case: $\max_{\mathcal{E}} \left\{ |E \mathbf{f}_{ij}| + var(\mathbf{f}_{ij})\phi^{-1}(1 - \epsilon_{ij}) \right\} \le f_{ij}^{max}$ (\mathcal{E} : data error model)

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Theorem. The robust problem is a convex optimization problem and can be solved in polynomial time in the polyhedral and ellipsoidal data cases.

Image: A math a math

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An "ambiguous chance-constrained problem"

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Toy example

1 What if no line limits?

2 What if tight limit on line connecting generators?



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Answer 1

What if no line limits?



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Answer 2

What if small limit on line connecting generators?



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Experiment

How much wind penetration can we handle? And how much money does this save?

39-bus New England system from MATPOWER



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Experiment

'standard' OPF solution with 10% buffer on line limits feasible only up to 5% penetration (right)



Cost 1,275,000 - almost 5 times greater than chance-constrained

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Polish system - winter 2003-04 evening peak



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Big cases

Polish 2003-2004 winter peak case

- 2746 buses, 3514 branches, 8 wind sources
- 5% penetration and $\sigma = .3\mu$ each source



CPLEX: the optimization problem has

- 36625 variables
- 38507 constraints, 6242 conic constraints
- 128538 nonzeros, 87 dense columns

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Big cases

CPLEX:

- total time on 16 threads = 3393 seconds
- "optimization status 6"
- solution is wildly infeasible

Gurobi:

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- time: 31.1 seconds
- "Numerical trouble encountered"

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Candidate solution violates conic constraint



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Separate: find a linear constraint also violated



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Solve again with linear constraint



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New solution still violates conic constraint



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Chance-constrained optimization problems in the operation of the power grid

Separate again



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We might end up with many linear constraints



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... which approximate the conic constraint



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conic constraint:

$$\sqrt{x_1^2 + x_2^2 + \ldots + x_k^2} = \|x\|_2 \le y$$

candidate solution:

$$(x^*, y^*)$$

cutting-plane (linear constraint):

$$||x^*||_2 + \frac{x^{*T}}{||x^*||_2}(x - x^*) = \frac{x^{*T}x}{||x^*||_2} \le y$$

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Polish 2003-2004 case CPLEX: "opt status 6" Gurobi: "numerical trouble"



Example run of cutting-plane algorithm:

Iteration	Max rel. error	Objective
1	1.2e-1	7.0933e6
4	1.3e-3	7.0934e6
7	1.9e-3	7.0934e6
10	1.0e-4	7.0964e6
12	8.9e-7	7.0965e6

Total running time: 32.9 seconds

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Back to motivating example

BPA case

- standard OPF: cost 235603, 7 lines unsafe \geq 8% of the time
- CC-OPF: cost 237297, every line safe \geq 98% of the time
- run time = 9.5 seconds (one cutting plane!)



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Back to motivating example

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Part II: Modeling line temperature

with Jose Blanchet and Juan Li (Columbia)

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Chance-constrained optimization problems in the operation of the power grid



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2003 North American blackout: initiated by several line trips



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- 2003 North American blackout: initiated by several line trips
- When a power line overheats it becomes exposed to several risk factors

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- When a power line overheats it becomes exposed to several risk factors
- If the line overheats enough, it may sag and experience a contact/arc, which will cause a trip
- If overheating is detected, and is deemed risky, the line will may be preemptively tripped

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- What is risky?

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- If overheating is detected, and is deemed risky, the line will may be preemptively tripped
- What is risky? What is a critical temperature?
- 2003 event: critical temperatures estimates were sometimes incorrect.

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 A comprehensive method for determining the temperature of a power line,

Chance-constrained optimization problems in the operation of the power grid

 A comprehensive method for determining the temperature of a power line, as a function of current and pause physical properties of the conductor

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Chance-constrained optimization problems in the operation of the power grid

- A comprehensive method for determining the temperature of a power line, as a function of current and pause physical properties of the conductor.
- It attempts to account for:

Chance-constrained optimization problems in the operation of the power grid

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Chance-constrained optimization problems in the operation of the power grid

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Chance-constrained optimization problems in the operation of the power grid

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Chance-constrained optimization problems in the operation of the power grid

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- It also relies on the heat equation for a "static" calculation.

Chance-constrained optimization problems in the operation of the power grid
IEEE Standard 738

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Note: power lines can be more than 100 miles long.

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- It also relies on the heat equation for a "static" calculation.
- Note: power lines can be more than 100 miles long.
- How can we account for data uncertainty, errors, unavailability?

Image: A math a math

- Line modeled as one-dimensional object parameterized by x, 0 ≤ x ≤ L.
- Time domain: $[0, \tau]$



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Chance-constrained optimization problems in the operation of the power grid

- Line modeled as one-dimensional object parameterized by x, 0 ≤ x ≤ L.
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Chance-constrained optimization problems in the operation of the power grid

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Chance-constrained optimization problems in the operation of the power grid

- Line modeled as one-dimensional object parameterized by x, 0 ≤ x ≤ L.
- Time domain: [0, *τ*] (for example: OPF intervals)
- I(t) = current at time t, T(x, t) = temperature at x at time t.
- Heat equation:

$$\frac{\partial T(x,t)}{\partial t} = \kappa \frac{\partial^2 T(x,t)}{\partial x^2} + \alpha I^2(t) - \nu (T(x,t) - T^{ext}(x,t)),$$

where $\kappa \ge 0$, $\alpha \ge 0$ and $\nu \ge 0$ are (line dependent) constants, and $T^{ext}(x, t)$ is the ambient temperature at (x, t)

Image: A math the second se

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Heat equation:

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IEEE 738, other authors:

$$\frac{\partial T(x,t)}{\partial t} = \alpha I^2(t) - \nu (T(x,t) - T^{ext}(x,t)).$$

This paper:

$$\frac{\partial T(x,t)}{\partial t} = \alpha I^2 - \nu (T(x,t) - G(\mathbf{h}(\mathbf{x}))).$$

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h(x) = a random variable, at x.

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This paper: stochasticity in the spatial domain (x)

CDC '13: stochasticity in the time domain (t)

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Image: A mathematical states and a mathem

Chance-constrained optimization problems in the operation of the power grid

This paper: stochasticity in the spatial domain (x)

CDC '13: stochasticity in the time domain (t)

The goal: algorithm- and data-driven estimates for "safe" current/temperature limits

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Chance-constrained optimization problems in the operation of the power grid

$$\frac{\partial T(x,t)}{\partial t} = \alpha I^2 - \nu (T(x,t) - G(\mathbf{h}(\mathbf{x})))$$

Recall: $x \in [0, L], t \in [0, \tau]$

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$$\frac{\partial T(x,t)}{\partial t} = \alpha I^2 - \nu (T(x,t) - G(\mathbf{h}(\mathbf{x}))).$$

Recall: $x \in [0, L]$, $t \in [0, \tau]$

Integrate and divide by L, get

$$\frac{1}{L}\int_0^L \frac{\partial T(x,t)}{\partial t}dx = \alpha l^2(t) - \frac{\nu}{L}\int_0^L T(x,t)dx + \frac{\nu}{L}\int_0^L G(\mathbf{h}(\mathbf{x}))dx.$$

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$$\frac{\partial T(x,t)}{\partial t} = \alpha I^2 - \nu (T(x,t) - G(\mathbf{h}(\mathbf{x}))).$$

Recall: $x \in [0, L]$, $t \in [0, \tau]$

Integrate and divide by L, get

$$\frac{1}{L} \int_0^L \frac{\partial T(x,t)}{\partial t} dx = \alpha l^2(t) - \frac{\nu}{L} \int_0^L T(x,t) dx + \frac{\nu}{L} \int_0^L G(\mathbf{h}(\mathbf{x})) dx.$$
$$\frac{1}{L} \int_0^L \frac{\partial T(x,t)}{\partial t} dx =$$

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$$\frac{1}{L} \int_{0}^{L} \frac{\partial T(x,t)}{\partial t} dx = \frac{d}{dt} \frac{1}{L} \int_{0}^{L} T(x,t) dx = \frac{d\mathbf{H}(\mathbf{t})}{dt}.$$
$$\mathbf{H}(\mathbf{t}) \triangleq \frac{1}{L} \int_{0}^{L} T(x,t) dx \quad (\text{average internal line temperature at t})$$

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 $\mathbf{R} \triangleq \frac{1}{L} \int_0^L G(\mathbf{h}(\mathbf{x})) dx$ (average ambient temperature,

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$$\frac{d\mathbf{H}(\mathbf{t})}{dt} = \alpha l^2(t) - \nu \mathbf{H}(\mathbf{t}) + \frac{\nu}{L} \int_0^L G(\mathbf{h}(\mathbf{x})) d\mathbf{x}.$$

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$$\frac{d\mathbf{H}(\mathbf{t})}{dt} = \alpha l^2(t) - \nu \mathbf{H}(\mathbf{t}) + \nu \mathbf{R}.$$

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$$\frac{d\mathbf{H}(\mathbf{t})}{dt} = \alpha I^2(t) - \nu \mathbf{H}(\mathbf{t}) + \nu \mathbf{R}.$$
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Chance-constrained optimization problems in the operation of the power grid

$$\frac{d\mathbf{H}(\mathbf{t})}{dt} = \alpha l^2(t) - \nu \mathbf{H}(\mathbf{t}) + \nu \mathbf{R}.$$
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Solution:

$$H(t) = \int_0^t e^{-\nu(t-s)} \alpha I^2(s) ds + R(1-e^{-\nu t}) + C e^{-\nu t},$$

where

$$C=\mathbf{H}(\mathbf{0})=\frac{1}{L}\int_0^L T(x,0)dx.$$

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Control goal: make I(t) "large",

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where

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Control goal: make I(t) "large", but with $P\left(\max_{t \in [0,\tau]} \mathbf{H}(\mathbf{t}) > k\right) \leq \epsilon$

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Chance-constrained optimization problems in the operation of the power grid

$$H(t) = \int_0^t e^{-\nu(t-s)} \alpha \bar{I}^2(s) ds + R(1-e^{-\nu t}) + Ce^{-\nu t},$$

where

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$$C = \mathbf{H}(\mathbf{0}) = \frac{1}{L} \int_0^L T(x, 0) dx.$$

Constant current \Rightarrow **H**(**t**) = $(\frac{\alpha}{\nu}\overline{l}^2 + \mathbf{R})(1 - e^{-\nu t}) + Ce^{-\nu t}$

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So, H'(t) > 0 for \overline{I} large enough,

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So, $P(\max_{t \in [0,\tau]} \mathbf{H}(\mathbf{t}) > k) \le \epsilon$ equivalent to $P(\mathbf{H}(\tau) > k) \le \epsilon$.

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Solution:

$$\overline{I}^2 \leq rac{
u}{lpha} rac{k - C e^{-
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ho_\epsilon (1 - e^{-
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So, H'(t) > 0 for \overline{I} large enough, (and of constant sign for any \overline{I}).

So, $P(\max_{t \in [0,\tau]} \mathbf{H}(\mathbf{t}) > k) \le \epsilon$ equivalent to $P(\mathbf{H}(\tau) > k) \le \epsilon$.

Solution:

$$\overline{I}^2 \leq \frac{\nu}{\alpha} \frac{k - Ce^{-\nu\tau} - \rho_\epsilon (1 - e^{-\nu\tau})}{1 - e^{-\nu\tau}} = L(\tau, k)$$

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Chance-constrained optimization problems in the operation of the power grid

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Simplification:

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R is a discrete random variable. $P(\mathbf{R} = r_i) = p_i, i = 1, 2, ..., n$ (known).



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1. At time $\tau = 0$, we compute values I_1 , and $I_{2,i}$ for i = 1, 2, ..., n. These values are used as follows:

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- 3. At time $\tau/2$, we observe the value of **R**. Assuming **R** = r_i , then for all $t \in [\tau/2, \tau]$, we set $I(t) = I_{2,i}$.

Goals:

(a)
$$P(\mathbf{H}(\boldsymbol{\tau}) > k) < \epsilon$$
.

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Goals:

(a) $P(\mathbf{H}(\tau) > k) < \epsilon$. k smaller than critical temperature (b) $l_1 \le L(\tau/2)$.

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Simplification:

R is a discrete random variable. $P(\mathbf{R} = r_i) = p_i, i = 1, 2, ..., n$ (known).

- 1. At time $\tau = 0$, we compute values I_1 , and $I_{2,i}$ for i = 1, 2, ..., n. These values are used as follows:
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Goals:

- (a) P(H(τ) > k) < ε. k smaller than critical temperature
 (b) l₁ ≤ L(τ/2).
- (c) What about performance?

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We want to maximize:

• "Total" current:
$$\frac{\tau}{2} I_1 + \frac{\tau}{2} I_{2,i}$$
 ?

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Chance-constrained optimization problems in the operation of the power grid

We want to maximize:

• "Total" current:
$$\frac{\tau}{2} I_1 + \frac{\tau}{2} I_{2,i}$$
 ?

• "Average" current? Square current?



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Chance-constrained optimization problems in the operation of the power grid

We want to maximize:

• "Total" current:
$$\frac{\tau}{2} I_1 + \frac{\tau}{2} I_{2,i}$$
 ?

• "Average" current? Square current?

 $F(I_1, I_2)$: a monotonely increasing function of I_1 , I_2

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Chance-constrained optimization problems in the operation of the power grid

Adaptive control

Simplification:

R is a discrete random variable. $P(\mathbf{R} = r_i) = p_i, i = 1, 2, ..., n$ (known).

- 1. At time $\tau = 0$, we compute values l_1 , and $l_{2,i}$ for i = 1, 2, ..., n. These values are used as follows:
- 2. For all $t \in [0, \tau/2]$, we set $I(t) = I_1$.
- 3. At time $\tau/2$, we observe the value of **R**. Assuming **R** = r_i , then for all $t \in [\tau/2, \tau]$, we set $I(t) = I_{2,i}$.

Goals:

(a) P(H(τ) > k) < ε. k smaller than critical temperature
(b) I₁ ≤ L(τ/2).
(c) Maximize:

$$\sum_{i=1}^n F(I_1,I_{2,i})p_i$$

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 $\begin{array}{ll} \max & \sum_{i=1}^{n} F(I_1, I_{2,i}) p_i \\ \text{s.t.} & P(\mathbf{H}(\boldsymbol{\tau}) > k) \leq \epsilon \end{array}$

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max
$$\sum_{i=1}^{n} F(l_{1}, l_{2,i}) p_{i}$$

s.t.
$$P(\mathbf{H}(\tau) > k) \le \epsilon$$
$$\mathbf{H}(\tau) \le u$$

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Chance-constrained optimization problems in the operation of the power grid

$$\begin{array}{ll} \max & \sum_{i=1}^{n} F(l_{1}, l_{2,i}) p_{i} \\ \text{s.t.} & P(\mathbf{H}(\boldsymbol{\tau}) > k) \leq \epsilon \\ & \mathbf{H}(\boldsymbol{\tau}) \leq u \quad (k < u < \text{ critical temp}) \end{array}$$

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Chance-constrained optimization problems in the operation of the power grid

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Chance-constrained optimization problems in the operation of the power grid

$$\begin{array}{ll} \max & \sum_{i=1}^{n} F(l_{1}, l_{2,i}) p_{i} \\ \text{s.t.} & P(\mathbf{H}(\boldsymbol{\tau}) > k) \leq \epsilon \\ & \mathbf{H}(\boldsymbol{\tau}) \leq u \quad (k < u < \text{ critical temp}) \\ & l_{1} \leq L(\tau/2, k) \\ & \text{other constraints.} \end{array}$$

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Chance-constrained optimization problems in the operation of the power grid

$$\begin{array}{ll} \max & \sum_{i=1}^{n} F(l_{1}, l_{2,i}) p_{i} \\ \text{s.t.} & P(\mathbf{H}(\boldsymbol{\tau}) > k) \leq \epsilon \\ & \mathbf{H}(\boldsymbol{\tau}) \leq u \quad (k < u < \text{ critical temp}) \\ & l_{1} \leq L(\tau/2, k) \\ & \text{other constraints.} \end{array}$$

$$\mathbf{H}(\tau) = \int_0^\tau e^{-\nu(\tau-s)} \alpha l^2(s) ds + \mathbf{R}(1-e^{-\nu\tau}) + C e^{-\nu\tau},$$

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$$\begin{array}{ll} \max & \sum_{i=1}^{n} F(l_{1}, l_{2,i}) p_{i} \\ \text{s.t.} & P(\mathbf{H}(\tau) > k) \leq \epsilon \\ & \mathbf{H}(\tau) \leq u \quad (k < u < \text{ critical temp}) \\ & l_{1} \leq L(\tau/2, k) \\ & \text{other constraints.} \end{array}$$

$$\mathbf{H}(\tau) = \int_0^\tau e^{-\nu(\tau-s)} \alpha l^2(s) ds + \mathbf{R}(1-e^{-\nu\tau}) + Ce^{-\nu\tau}, \\ = v_1 l_1^2 + v_2 l_{2,i}^2 + r_i(1-e^{-\nu\tau}) + Ce^{-\nu\tau} \quad \text{in state i}$$

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$$\begin{array}{ll} \max & \sum_{i=1}^{n} F(I_{1}, I_{2,i}) p_{i} \\ \text{s.t.} & P(\mathbf{H}(\tau) > k) \leq \epsilon \\ & \mathbf{H}(\tau) \leq u \quad (k < u < \text{ critical temp}) \\ & I_{1} \leq L(\tau/2, k) \\ & \text{other constraints.} \end{array}$$

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So chance constraint is of the form:

$$\sum_{i=1}^{n} \mathbb{I}\{v_{1} \ l_{1}^{2} + v_{2} \ l_{2,i}^{2} > u - r_{i}(1 - e^{-\nu\tau}) - Ce^{-\nu\tau}\}p_{i} \leq \epsilon.$$

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$$\begin{array}{ll} \max & \sum_{i=1}^{n} F(l_{1}, l_{2,i}) p_{i} \\ \text{s.t.} & P(\mathbf{H}(\tau) > k) \leq \epsilon \\ & \mathbf{H}(\tau) \leq u \quad (k < u < \text{ critical temp}) \\ & l_{1} \leq L(\tau/2, k) \\ & \text{other constraints.} \end{array}$$

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So chance constraint s of the form:

$$\sum_{i=1}^{n} \mathbb{I}\{\underbrace{v_{1} \ l_{1}^{2}}_{z_{1}} + \underbrace{v_{2} \ l_{2,i}^{2}}_{z_{2}(i)} > \underbrace{u \ - \ r_{i}(1 - e^{-\nu\tau}) \ - \ Ce^{-\nu\tau}}_{w_{i}}\}p_{i} \leq \epsilon.$$

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$$\max \qquad \sum_{i=1}^{n} \tilde{F}(z_1, z_2(i)) p_i$$

s.t.
$$\sum_{i=1}^{n} \mathbb{I}\{z_1 + z_2(i) > w_i\} p_i \leq \epsilon$$
$$z_1 + z_2(i) \leq u_i \quad (w_i < u_i)$$
$$z_1 \leq \bar{k}$$

other constraints.



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Lemma: At optimality,

$$z_1 + z_2(i) = w_i$$
 or u_i , all i

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other constraints.

Lemma: At optimality,

$$z_1 + z_2(i) = w_i$$
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 \rightarrow Use **binary** variable

$$y_i = \begin{cases} 0 & \text{when } z_1 + z_2(i) = w_i \\ 1 & \text{when } z_1 + z_2(i) = u_i \end{cases}$$

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$$\max \sum_{i=1}^{n} \tilde{F}(z_{1}, w_{i} - z_{1})p_{i}(1 - y_{i}) + \tilde{F}(z_{1}, u_{i} - z_{1})p_{i}y_{i}$$

s.t.
$$\sum_{i=1}^{n} u_{i}p_{i}y_{i} \leq \epsilon$$
$$0 \leq z_{1} \leq \bar{k}$$
$$y_{i} = 0 \text{ or } 1, \text{ all } i.$$

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 $\max_{z_1\in[0,\bar{k}]}M(z_1)$



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Chance-constrained optimization problems in the operation of the power grid

 $\max_{z_1\in[0,\bar{k}]}M(z_1)$

$$M(z_1) \triangleq \sum_{i=1}^{n} \tilde{F}(z_1, w_i - z_1) p_i (1 - y_i) + \tilde{F}(z_1, u_i - z_1) p_i y_i$$

s.t.
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Practicable!

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Practicable! Grid over *z*₁

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s.t.
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Practicable! Grid over z_1 + knapsack for given z_1

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Chance-constrained optimization problems in the operation of the power grid

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Theorem. (B. and Mc Closky 2010) Consider a 0 - 1 knapsack problem on *N* variables

$$\max \sum_{j=1}^{N} p_j x_j, \qquad \text{s.t. } \sum_{j=1}^{n} a_j x_j \le b,$$
$$x_j = 0 \text{ or } 1, \text{ all } j.$$

For each fixed tolerance 0 $<\delta<1$ there is a linear program LP with the following properties

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Chance-constrained optimization problems in the operation of the power grid

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For each fixed tolerance $0<\delta<1$ there is a linear program LP with the following properties

• The number of variables and constraints in LP is $O(N^2)$.

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Chance-constrained optimization problems in the operation of the power grid

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Image: A math a math

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For each fixed tolerance $0<\delta<1$ there is a linear program LP with the following properties

- The number of variables and constraints in LP is $O(N^2)$.
- The x_j are among the variables of LP (a "lifted" formulation)
- The solution of LP, together with a simple rounding for the x_j variables yields a (binary) solution for the knapsack that is guaranteed to be within δ of the optimum.

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$$\begin{array}{ll} \max & \sum_{i=1}^{n} F(l_{1}, l_{2,i}) p_{i} \\ \text{s.t.} & P(\mathbf{H}(\tau) > k) \leq \epsilon \\ & \mathbf{H}(\tau) \leq u \quad (k < u < \text{ critical temp}) \\ & l_{1} \leq L(\tau/2, k), \quad \text{etc.} \end{array}$$

$$\mathbf{H}(\tau) = v_1 l_1^2 + v_2 l_{2,i}^2 + r_i(1 - e^{-\nu\tau}) + C e^{-\nu\tau} \text{ in state i} \\ v_1 = \int_0^{\tau/2} e^{-\nu(\tau-s)} \alpha \, ds, \quad v_2 = \int_{\tau/2}^{\tau} e^{-\nu(\tau-s)} \alpha \, ds$$

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$$\begin{array}{ll} \max & \sum_{i=1}^{n} F(l_{1}, l_{2,i}) p_{i} \\ \text{s.t.} & P(\mathbf{H}(\tau) > k) \leq \epsilon \\ & \mathbf{H}(\tau) \leq u \quad (k < u < \text{ critical temp}) \\ & l_{1} \leq L(\tau/2, k), \quad \text{etc.} \end{array}$$

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What if $F(I_1, I_{2,i}) = v_1 I_1^2 + v_2 I_{2,i}^2$?

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What if $F(I_1, I_{2,i}) = v_1 I_1^2 + v_2 I_{2,i}^2$? Lemma: (again) At optimality, $v_1 I_1^2 + v_2(i) I_{2,i}^2 = w_i$ or u_i , all i

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A straight knapsack problem

$$\max \sum_{i=1}^{n} \tilde{F}(w_i) p_i (1 - y_i) + \tilde{F}(u_i) p_i y_i$$

s.t.
$$\sum_{i=1}^{n} u_i p_i y_i \leq \epsilon$$
$$y_i = 0 \text{ or } 1, \text{ all } i.$$

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