

Thm Top Sort Alg is correct.  
i.e.  $(x, y) \in E \Rightarrow f(x) > f(y)$ .

Pf  
 $x \rightarrow y$

i)  $d(x) < d(y)$

Then at  $d(x)$ ,  $y$  is white

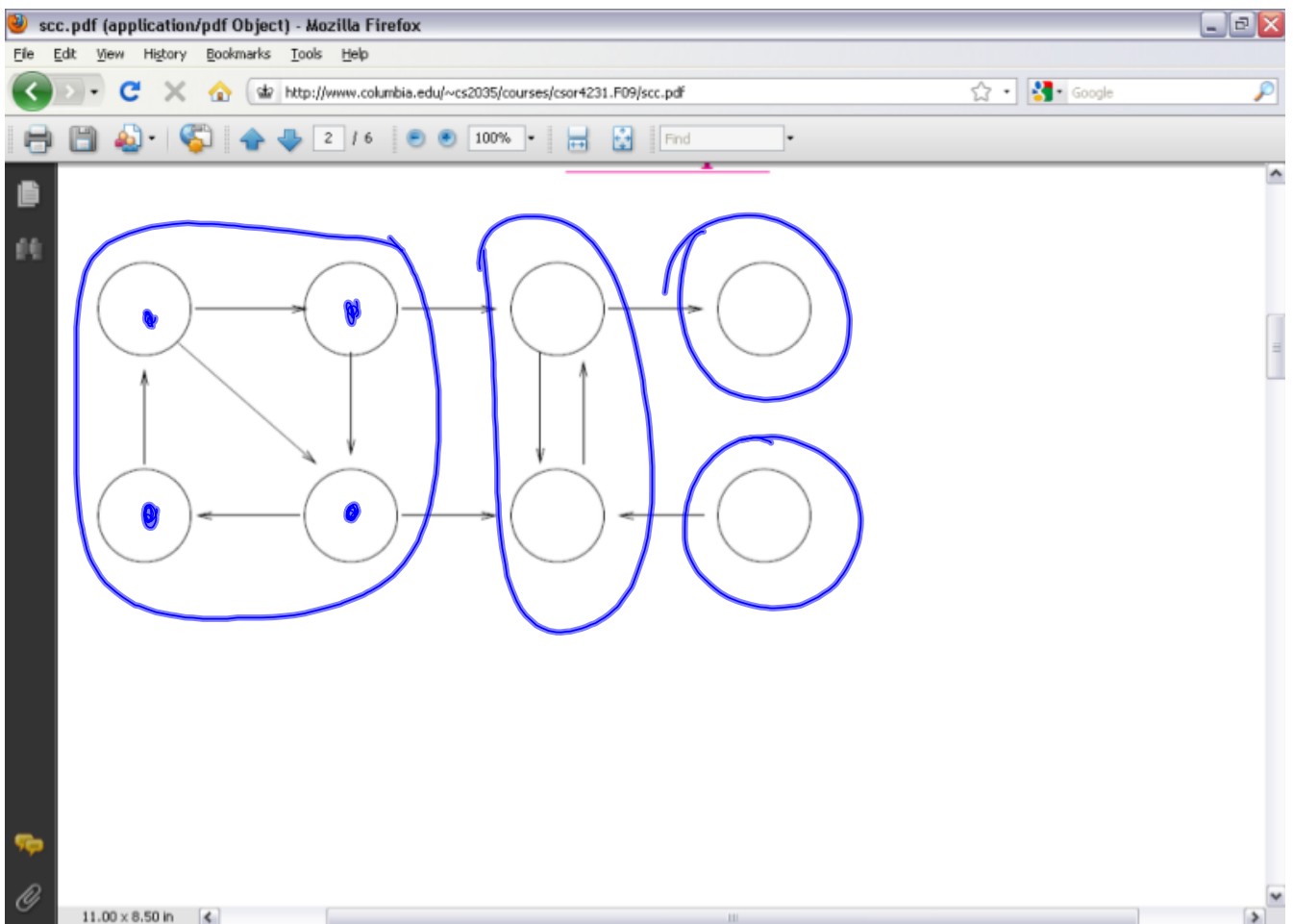
$$d(x) < d(y) < \underline{f(y) < f(x)}$$

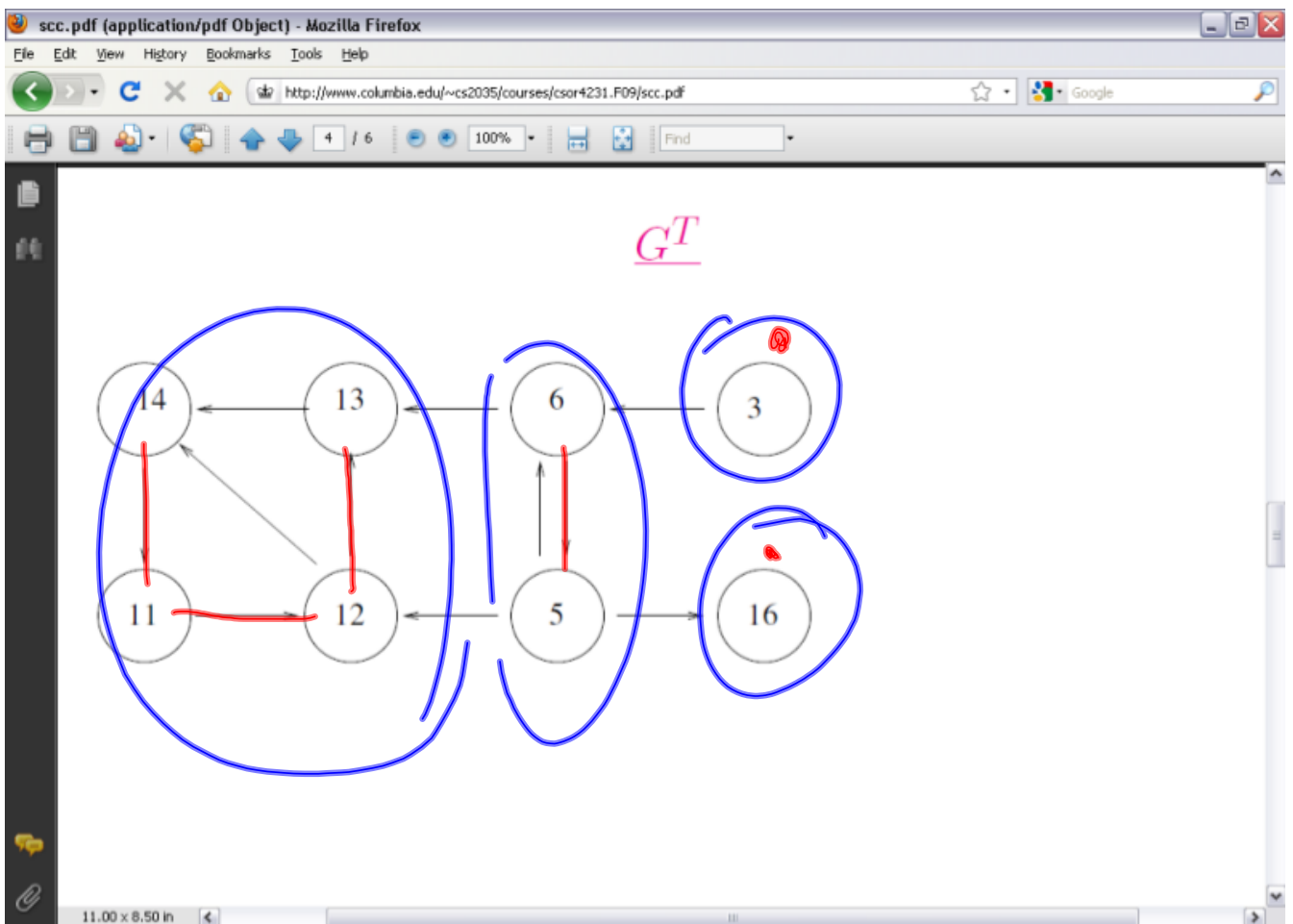
ii)  $d(x) > d(y)$

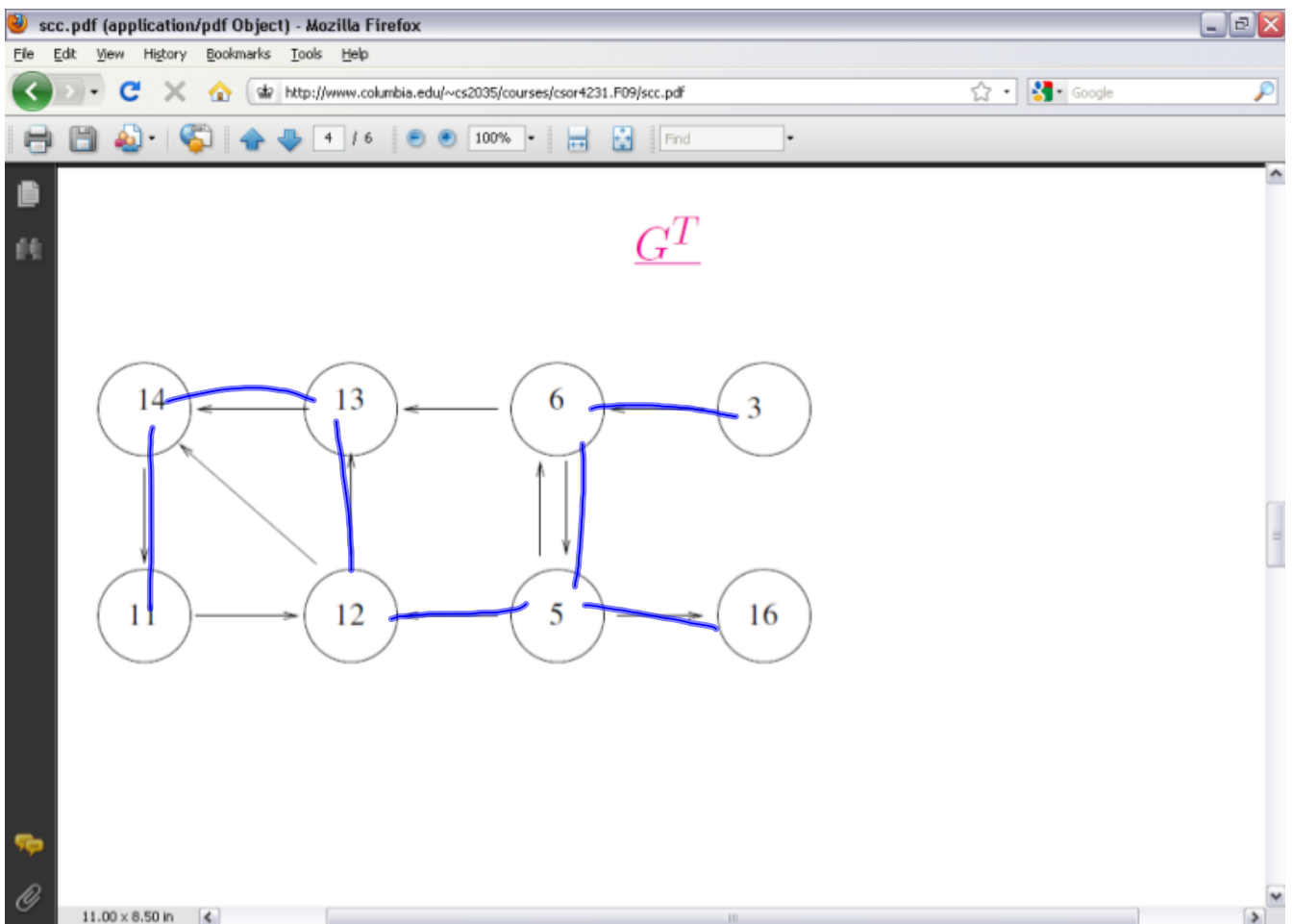
Because  $G$  is a DAG, no path from  
 $y$  to  $x \Rightarrow \underline{f(y) < d(x)} < \underline{f(x)}$   $\square$

Connectivity  
Given an undirected graph, find connected components

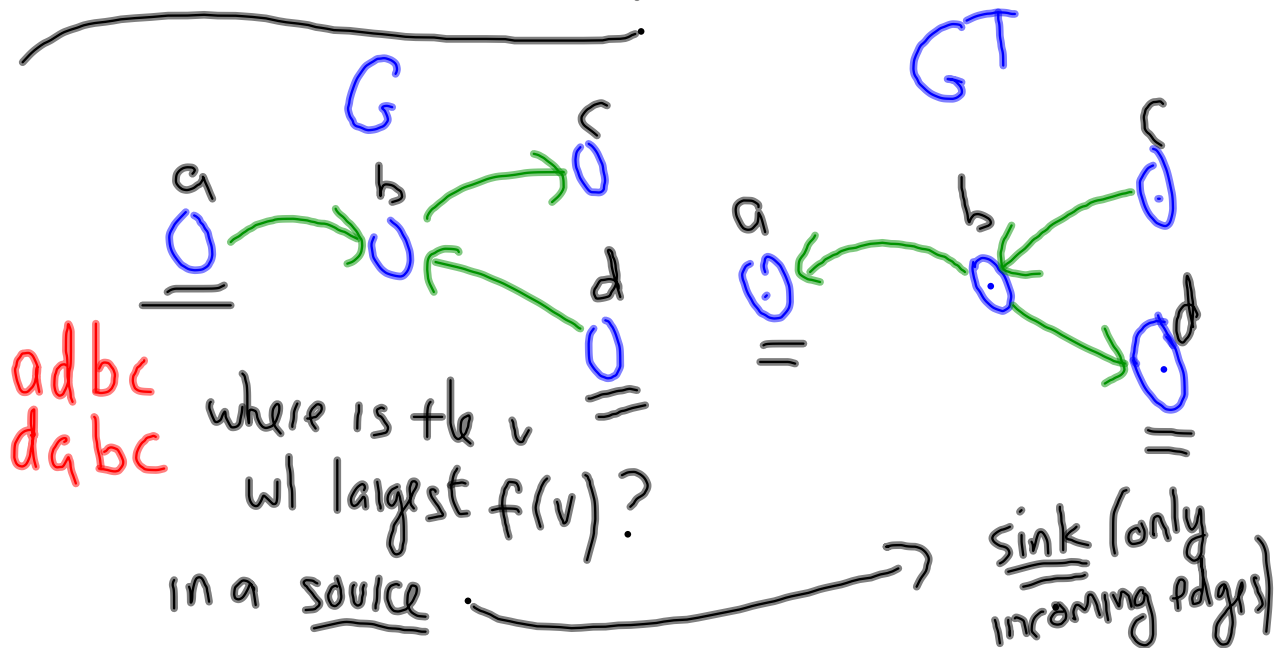








## Acyclic Component Graph



at each step, we explore a component<sup>(node)</sup> of the component graph, & we do so in topological order of  $G$ ,  $\therefore$  the component has no outgoing edges in  $G^T$  to an unexplored component.

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1 / 6 100% Find

## Strongly Connected Components

**Definition** A strongly connected component of a directed graph  $G$  is a maximal set of vertices  $C \subseteq V$  such that for every pair of vertices  $u$  and  $v$ , there is a directed path from  $u$  to  $v$  and a directed path from  $v$  to  $u$ .

Strongly-Connected-Components( $G$ )

- 1 call DFS( $G$ ) to compute finishing times  $f[u]$  for each vertex  $u$   $O(E)$
- 2 compute  $G^T$
- 3 call DFS( $G^T$ ), but in the main loop of DFS, consider the vertices  $O(E)$   
in order of decreasing  $f[u]$  (as computed in line 1)  $O(E)$
- 4 output the vertices of each tree in the depth-first forest formed in line 3  
separate strongly connected component  $O(E)$

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4 / 10 79.1% Find

- A **cut** in a graph is a partition of the vertices into two sets  $S$  and  $T$ .
- An edge  $(u, v)$  with  $u \in S$  and  $v \in T$  is said to **cross the cut**.

The graph consists of 9 vertices and 10 edges. The edges and their weights are: (1,2) weight 6, (1,3) weight 14, (2,3) weight 5, (2,4) weight 4, (3,5) weight 10, (3,6) weight 3, (4,7) weight 9, (5,6) weight 2, (6,8) weight 8, (7,8) weight 15. A red line represents a cut that partitions the vertices into two sets:  $S = \{1, 2, 3, 5, 6, 8\}$  and  $T = \{4, 7, 8\}$ . The edges crossing the cut are (4,7), (5,6), and (6,8).

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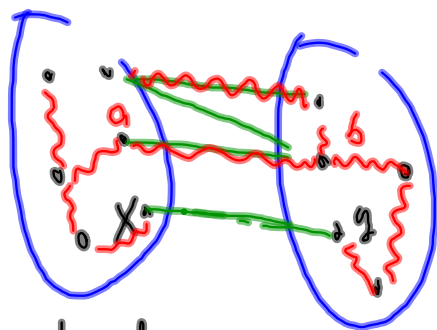
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5 / 10 79.1% Find

- A **cut** in a graph is a partition of the vertices into two sets  $S$  and  $T$ .
- An edge  $(u, v)$  with  $u \in S$  and  $v \in T$  is said to **cross the cut**.

Claim For any cut  $(S, T)$ , the minimum wt. edge crossing the cut is in the min. spanning tree.

Pf Assume not. There must be some cut for which the min wt. edge crossing the cut is not in the minimum spanning tree.



red is alleged MST  
there is at least one edge in T crossing the cut

$(x, y)$  is the min. wt edge crossing the cut ( $\notin T$ ).

Now let  $(a, b)$  be the edge in  $T$  on the path from  $x$  to  $y$  that crosses the cut. Consider  $T' = T \cup \{(x, y)\} - (a, b)$ .

$T'$  is a spanning tree. and

$$w(T') = w(T) + w(x, y) - w(a, b) < w(T).$$

which contradicts  $w(T)$  being a minimum spanning tree.  $\otimes$

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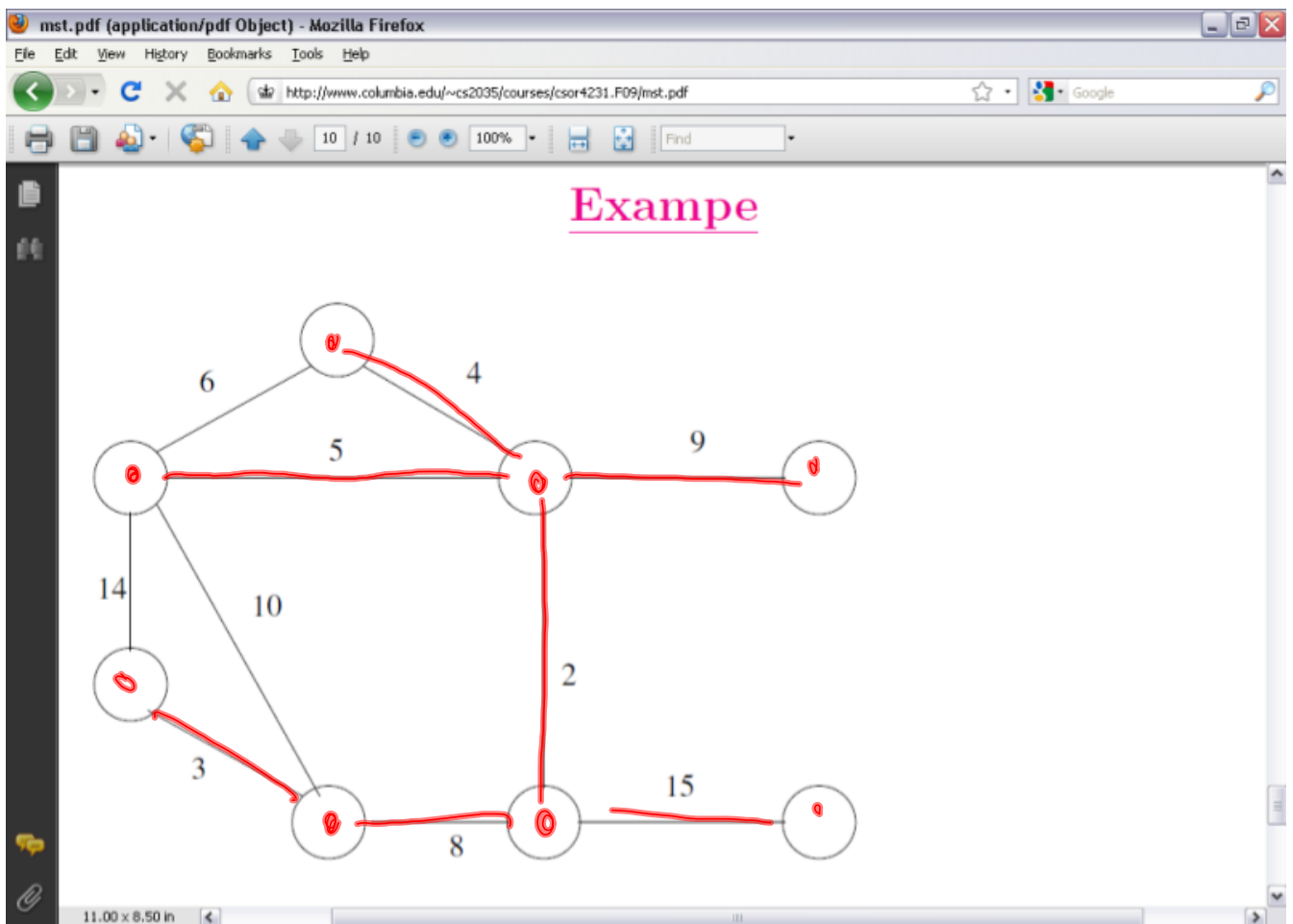
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- The total cost (weight) of a spanning tree  $T$  is defined as  $\sum_{e \in T} w(e)$
- A **minimum spanning tree** is a tree of minimum total weight.

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3 / 10 100% Find

- An edge  $(u, v)$  with  $u \in S$  and  $v \in T$  is said to **cross the cut**.

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1 / 10 100% Find

## Minimum Spanning Trees

- $G = (V, E)$  is an undirected graph with non-negative edge weights  $w : E \rightarrow \mathbb{R}$
- We assume wlog that edge weights are distinct
- A **spanning tree** is a tree with  $V - 1$  edges, i.e. a tree that connects all the vertices.
- The total cost (weight) of a spanning tree  $T$  is defined as  $w(T) = \sum_{e \in T} w(e)$
- A **minimum spanning tree** is a tree of minimum total weight.

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