

x_1 = # soldiers manufactured

x_2 = # trains

$$\begin{array}{ll} \max & 3x_1 + 2x_2 \\ \text{s.t.} & \end{array}$$

$$2x_1 + x_2 \leq 100$$

$$x_1 + x_2 \leq 80$$

$$x_1 \leq 40$$

$$x_1, x_2 \geq 0$$

-feas. sol'n
satisfies
constraints

-opt. sol'n,
feasible &
maximizes obj.
function

$$x_1=10 \ x_2=10 \ \text{obj}=50 \ / \ x_1=20 \ x_2=60 \ \text{obj}=180$$

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Fractional Knapsack

V_i pick some of each item, s.t.

w_i total wt. $\leq W$, value maximized.

W

x_i = fraction of item i included.

$$\begin{array}{ll} \text{linear} & \max \sum_{i=1}^n v_i x_i \\ \text{program} & \text{s.t.} \sum_{i=1}^n w_i x_i \leq W \\ & 0 \leq x_i \leq 1 \quad i=1 \dots n \end{array}$$

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Solving a LP.

n vars. m constraints.

- Algs to solve LP in polynomial time
(interior point, ellipsoid)

3 4

- Simplex (not known polynomial time)
Fast in practice (CPLEX,)

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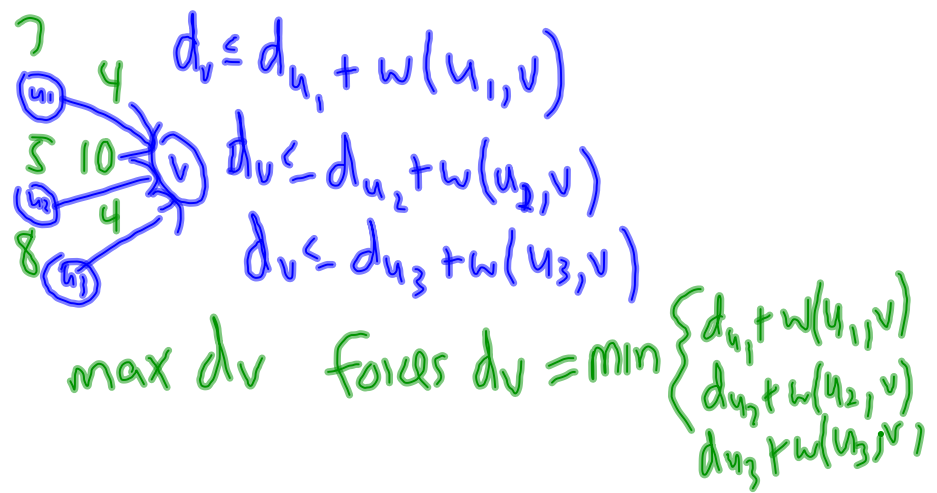
Integer Programming

LP w/ constraints on some
vars that

x_i is an integer.

Integer Programming is NP-complete.

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Max flow &

$$f(a, b) + f(b, c) \leq 2f(a, c)$$

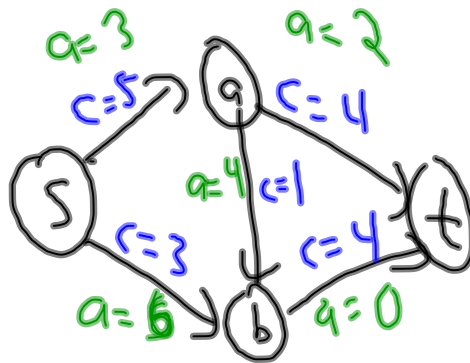
$$f(e, f) \leq 1$$

$$f(h, b) \geq 4$$

still a LP

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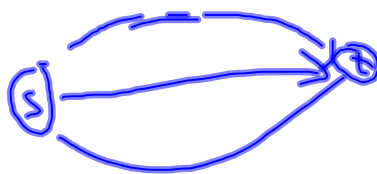
min cost flow



I want to route d units of flow from s to t as cheaply as possible.

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Flow



Nov 19-5:01 PM

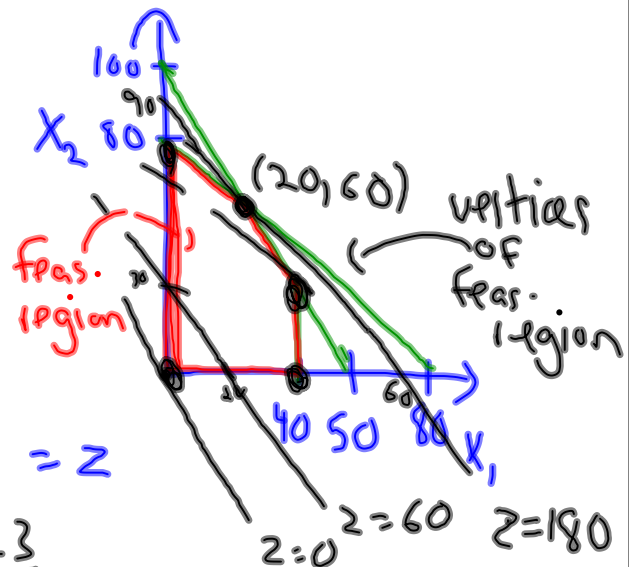
$$\begin{aligned} \max \quad & 3x_1 + 2x_2 \\ \text{s.t.} \quad & 2x_1 + x_2 \leq 100 \\ & x_1 + x_2 \leq 80 \\ & x_1 \leq 40 \\ & x_1, x_2 \geq 0 \end{aligned}$$

$(1, 180)$
 $(180, 1)$

$$\begin{aligned} \min \quad & 100y_1 + 80y_2 + 40y_3 \\ \text{s.t.} \quad & 2y_1 + y_2 + y_3 \geq 3 \\ & y_1 + y_2 \geq 2 \end{aligned}$$

$$3x_1 + 2x_2 = z$$

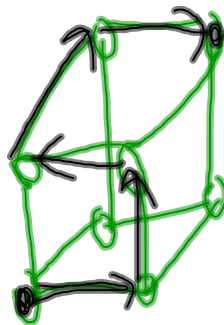
slope = $-\frac{3}{2}$



If feasible region is non-empty & bounded then
 $\exists$ an opt sol. which is a vertex of feasible region

Nov 19-5:07 PM

Simplex moves from vertex to adj.
 vertex in the feasible region,
 increasing (not decreasing) the
 obj. value at each step.



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