

Maximum Flows

- A **flow network** $G = (V, E)$ is a directed graph in which each edge $(u, v) \in E$ has a nonnegative **capacity** .
- If $(u, v) \notin E$, we assume that $c(u, v) = 0$.
- We distinguish two vertices in a flow network: a **source** s and a **sink** t .

A **flow** in G is a real-valued function $f : V \times V \rightarrow R$ that satisfies the following three properties:

Capacity constraint: For all $u, v \in V$, we require $0 \leq f(u, v) \leq c(u, v)$.

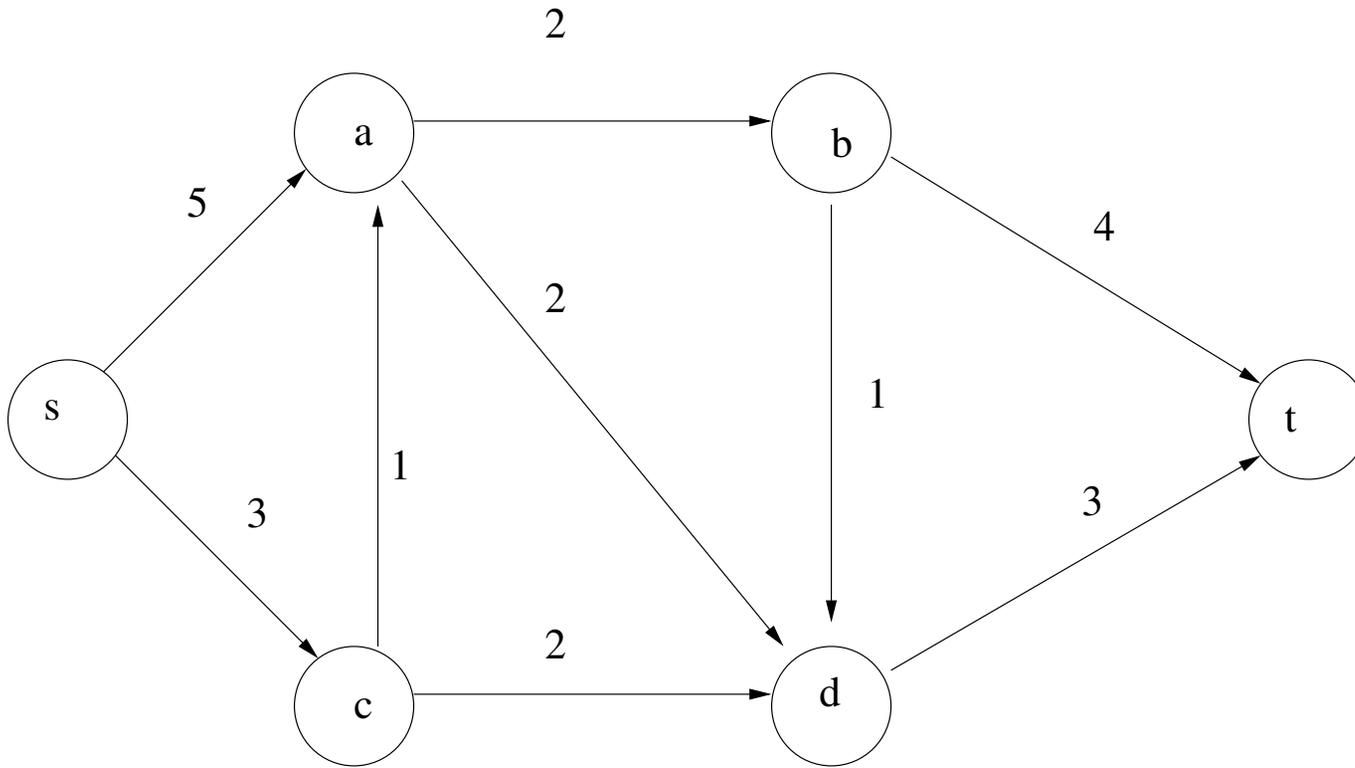
Flow conservation: For all $u \in V - \{s, t\}$, we require

$$\sum_{v \in V} f(v, u) = \sum_{v \in V} f(u, v) .$$

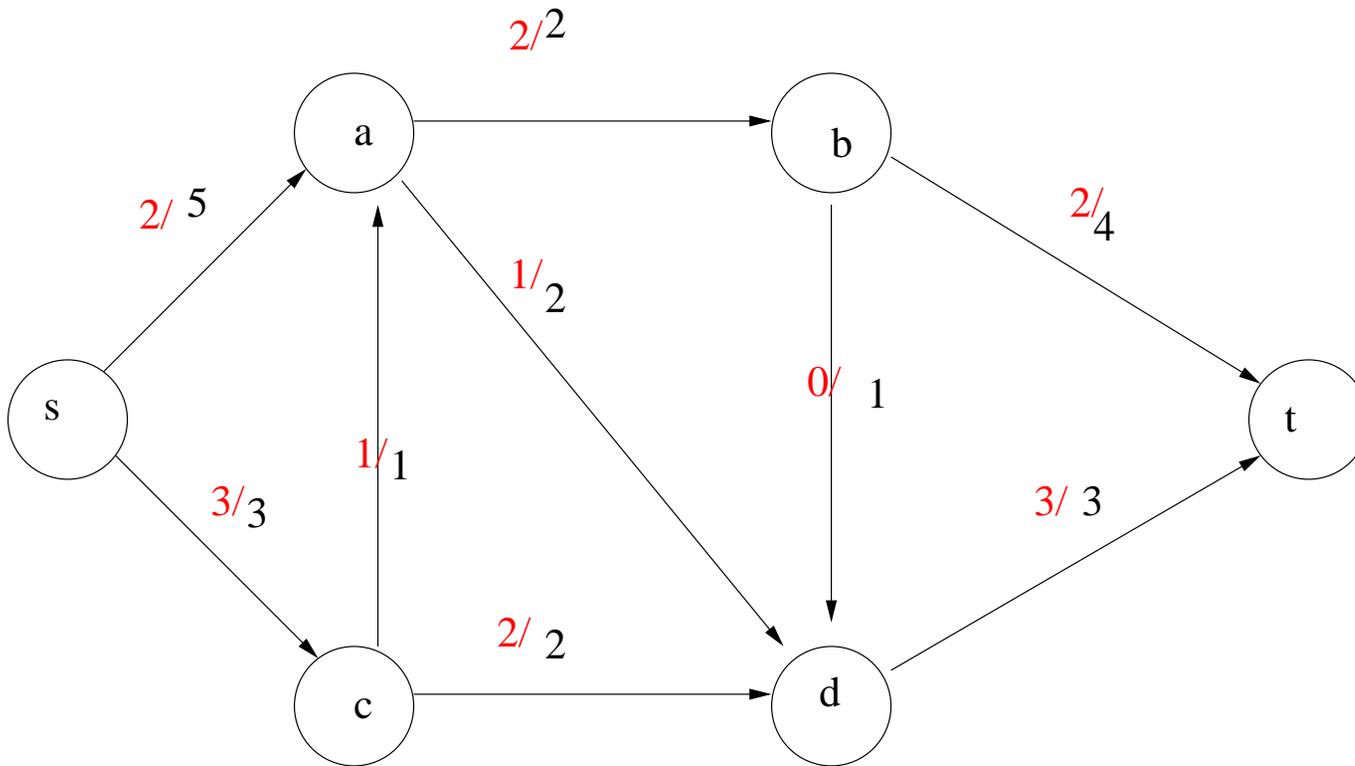
The **value** of a flow f is defined as

$$|f| = \sum_{v \in V} f(s, v) - \sum_{v \in V} f(v, s) , \tag{1}$$

Example



Solutions



Ford Fulkerson

Ford-Fulkerson-Method(G, s, t)

- 1 initialize flow f to 0
- 2 while there exists an augmenting path p
- 3 augment flow f along p
- 4 return f

- The **residual capacity** of (u, v) , is

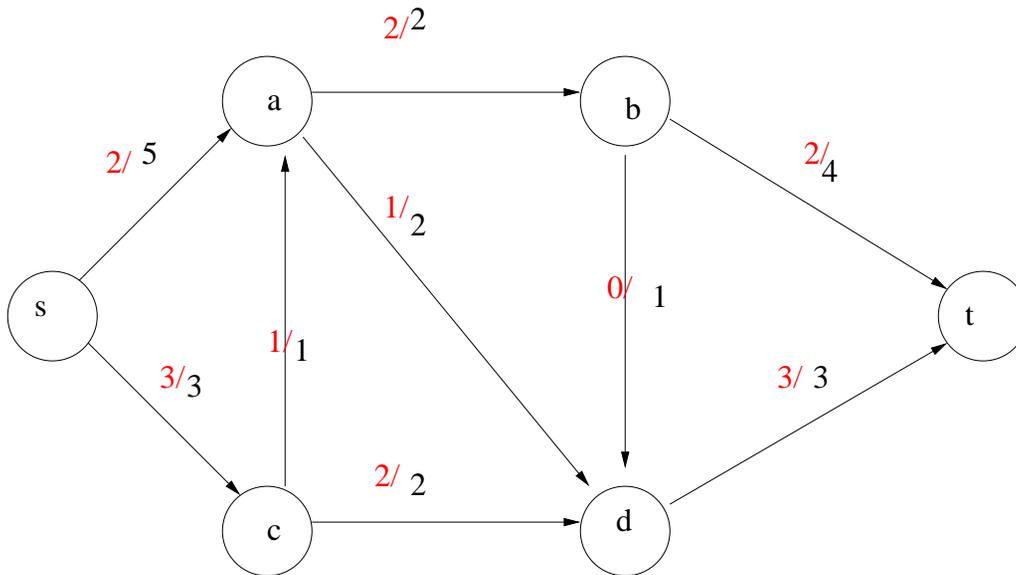
$$c_f(u, v) = \begin{cases} c(u, v) - f(u, v) & \text{if } (u, v) \in E , \\ f(v, u) & \text{if } (v, u) \in E , \\ 0 & \text{otherwise .} \end{cases} \quad (2)$$

- The residual graph G_f is the graph consisting of edges with positive residual capacity
- Flows add

s-t Cuts

An *s-t* cut satisfies

- $s \in S, t \in T$
- $S \cup T = V, S \cap T = \emptyset$



$$c(S, T) = \sum_{u \in S, v \in T} c(u, v)$$

$$f(S, T) = \sum_{u \in S, v \in T} f(u, v)$$

- For all cuts (S, T) and all feasible flows f , $f(S, T) \leq c(S, T)$
- For all pairs of cuts (S_1, T_1) and (S_2, T_2) , and all feasible flows f , $f(S_1, T_1) = f(S_2, T_2)$.

Max-flow min-cut theorem

If f is a flow in a flow network $G = (V, E)$ with source s and sink t , then the following conditions are equivalent:

1. f is a maximum flow in G .
2. The residual network G_f contains no augmenting paths.
3. $|f| = c(S, T)$ for some cut (S, T) of G .

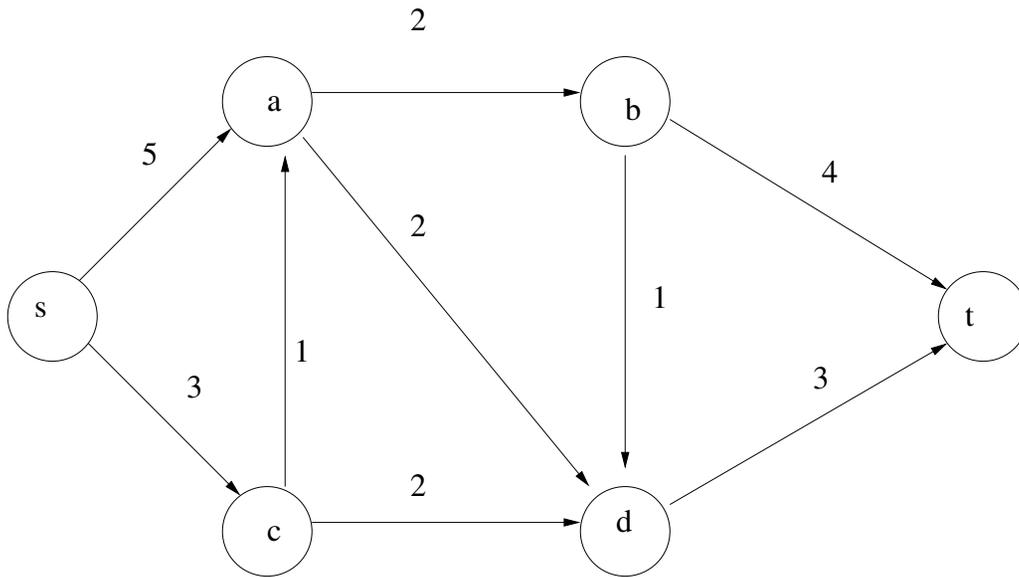
Ford Fulkerson expanded

Ford – Fulkerson(G, s, t)

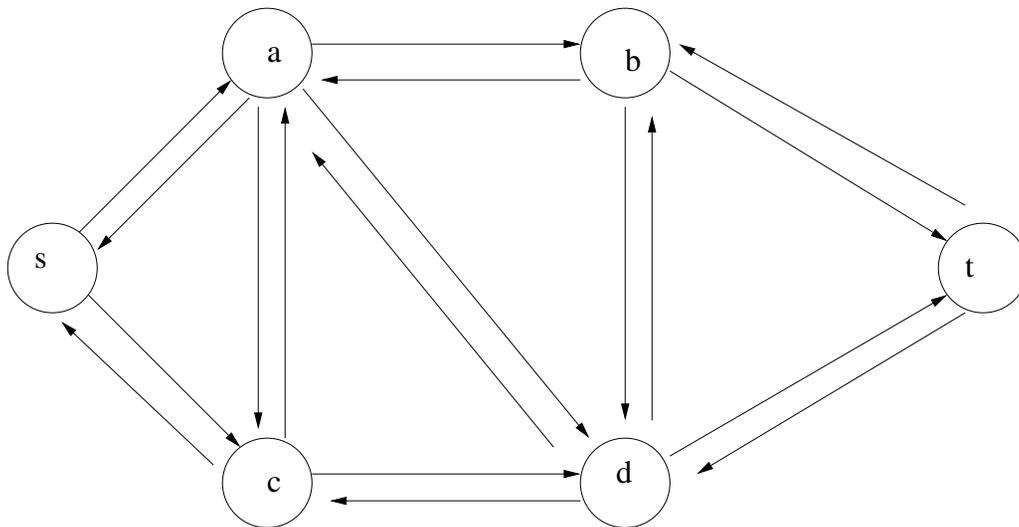
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1  for each edge  $(u, v) \in G.E$ 
2       $(u, v).f = 0$ 
3  while there exists a path  $p$  from  $s$  to  $t$  in the residual network  $G_f$ 
4       $c_f(p) = \min\{c_f(u, v) : (u, v) \text{ is in } p\}$ 
5      for each edge  $(u, v)$  in  $p$ 
6          if  $(u, v) \in E$ 
7               $(u, v).f = (u, v).f + c_f(p)$ 
8          else  $(v, u).f = (v, u).f - c_f(p)$ 
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Algorithm

Graph

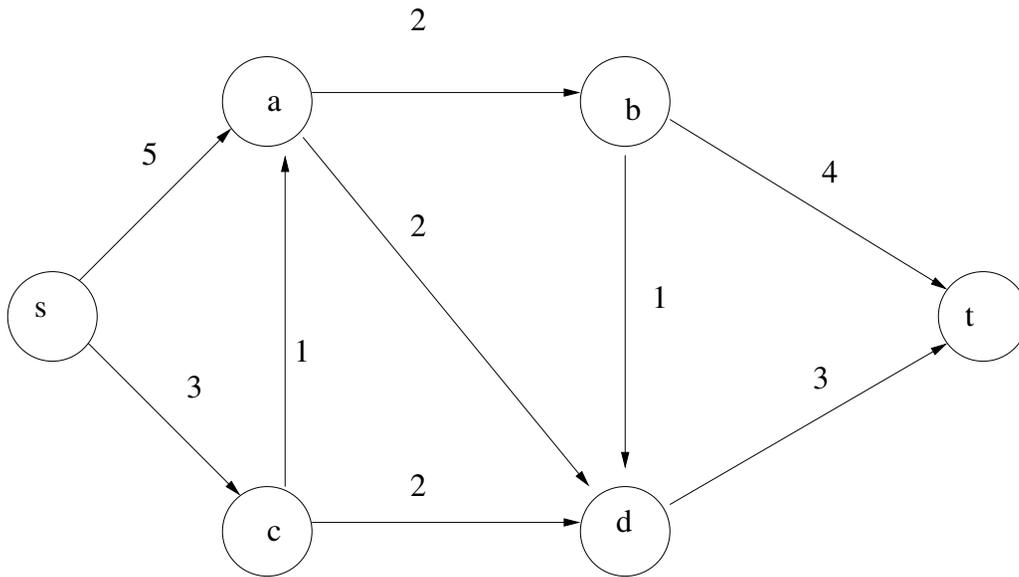


Residual graph

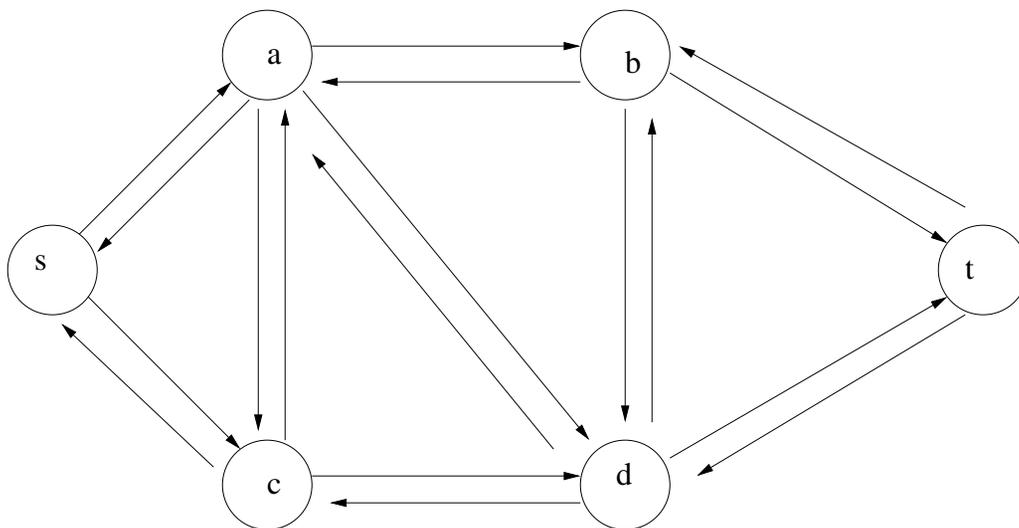


Algorithm

Graph

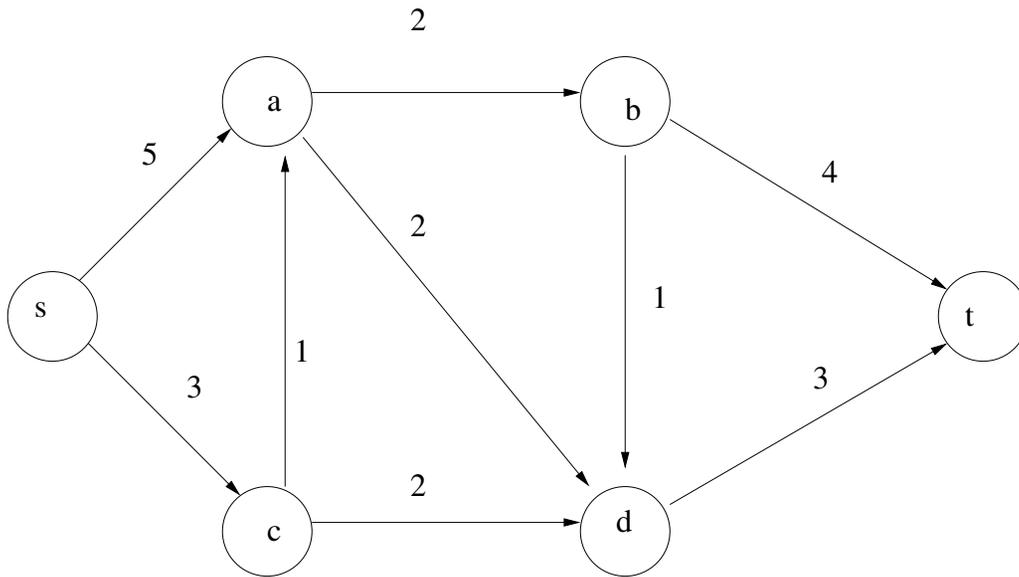


Residual graph



Algorithm

Graph



Residual graph

