## **Average Completion Time on Multiple Machines**

$P    \Sigma C_j$				
j	$p_j$			
A	1			
В	3			
$\mathbf{C}$	<b>4</b>			
D	<b>5</b>			
$\mathbf{E}$	6			
$\mathbf{F}$	9			
G	<b>12</b>			
Η	<b>20</b>			
Ι	<b>50</b>			
$\mathbf{J}$	60			

What is the right algorithm?

## **Average Completion Time on Multiple Machines**

- $P|| \Sigma C_J SPT$  is optimal.
- $P|| \sum w_j C_j$  Is WSPT optimal?

Example

### **Average Completion Time on Multiple Machines**

- $P|| \sum C_J$  SPT is optimal.
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#### Example

- $P|| \sum w_j C_j$  is NP-complete.
- WSPT is a  $(1 + \sqrt{2})/2$  -approximation for  $P || \sum w_j C_j$

## $|R|| \operatorname{S} C_j$

- Can be solved as a matching problem.
- Left side node for each job j
- Right hand side node for the k th from last job on machine i

#### Example

	$J_1$	$J_2$	$J_3$	$J_4$
$M_1$	6	4	$\infty$	3
$M_2$	<b>7</b>	<b>5</b>	<b>2</b>	3
$M_3$	3	8	<b>5</b>	3

### LP for the matching problem

Variable  $x_{ijk} = 1$  if j is the k th from last job on  $M_i$ 

$$\begin{split} \min \sum_{i=1}^{m} \sum_{j=1}^{n} \sum_{k=1}^{n} k p_{ij} x_{ijk} \\ \textbf{s.t.} \\ \textbf{Each job runs} \\ \sum_{i=1}^{m} \sum_{k=1}^{n} x_{ijk} = 1 \qquad \qquad j = 1 \dots n \\ \textbf{Each machine/slot has at most 1 job} \\ \sum_{j=1}^{n} x_{ijk} \leq 1 \qquad \qquad i = 1 \dots m; k = 1 \dots n \\ x_{ijk} \in \{0, 1\} \quad i = 1 \dots m; j = 1 \dots n; k = 1 \dots n \end{split}$$

• Note that the may be unforced idleness e.g.

 $\begin{array}{cccc} & J_1 & J_2 \\ M_1 & {\bf 1} & {\bf 1} \\ M_2 & {\bf 10} & {\bf 10} \end{array}$ 

# $Q|\text{pmtn}| \Sigma C_j$

- Algorithm is SRPT-FM. Shortest Remaining Processing Time on the Fastest Machines.
- What about preemption in other models?
- $\bullet$  P doesn't help
- $\bullet \ R NP\text{-complete}$