

## Min Cost Cycle Cancelling

1. Find a feasible flow  $f$  (max flow problem)

2. While  $\exists$  a neg cost cycle in  $G_f$

- let  $X$  be the min. cost cycle

- let  $\delta = \min_{(u,w) \in X} u_f(u,w)$

- send  $\delta$  units of flow around  $X$ .  
(maintain  $\pi$ 's at nodes)

- $\epsilon$ -optimality

-flows are always feasible (in this dg.)

A flow is  $\epsilon$ -opt. if  $\exists \pi$  s.t.

$$c^\pi(v,w) \geq -\epsilon \quad \forall (v,w) \in G_f.$$

Lemma

1) Any flow is  $C$ -optimal

2) If  $\epsilon < \frac{1}{n}$ , then an  $\epsilon$ -optimal flow is optimal.

Pf

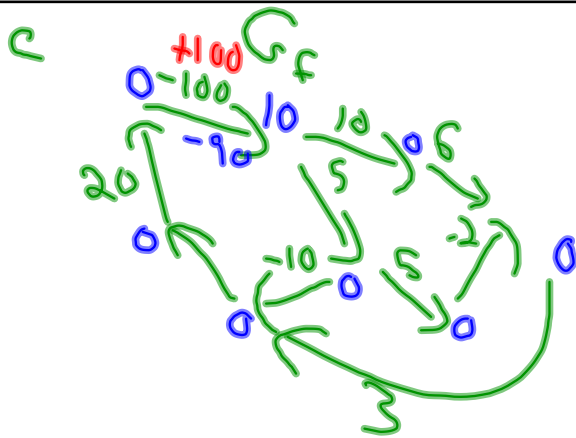
1) Set  $\bar{\pi} = 0$

2) Consider a cycle  $X$  in  $G_f$ .  $\left. \begin{array}{l} c^{\bar{\pi}}(v,w) \geq -\epsilon > -\frac{1}{n} \\ |X| \leq n \end{array} \right\}$

$$\sum_{(v,w) \in X} c^{\bar{\pi}}(v,w) > \left(-\frac{1}{n}\right) |X| > -\frac{1}{n}(n) > -1. \geq 0$$

$$\Rightarrow \sum_{(v,w) \in X} c(v,w) \geq 0 \Rightarrow f \text{ is opt.}$$

$$\begin{array}{l} n = 99 \\ \frac{1}{100} \text{-opt.} \\ > -\frac{1}{100} \quad \left( \frac{99}{100} \right) \\ > -\frac{99}{100} \end{array}$$



$$\Sigma = 100$$

$$\Sigma = 90$$

$$\Sigma = 0$$

Given  $\pi, f$  let  $\epsilon^\pi(f) = \min_{v, w \in G_f} \{c^\pi(v, w)\}$   
 . What is the best  $\pi$ ?

$$\text{let } \epsilon(f) = \min_{\pi} \epsilon^\pi(f).$$

let  $\nu(f) = \min_{\text{value}} \text{mean cycle in } G_f$

Thm  $\epsilon(f) = -\mu(f)$ .

Pf 1)  $\epsilon(f) \geq -\mu(f)$

$f$  is  $\epsilon(f)$ -opt.

$$\Rightarrow \forall \text{ cycles } X \quad \sum_{(v,w) \in X} c(v,w) = \sum_{(v,w) \in X} c^{\pi}(v,w) \geq -\epsilon(f) |X|$$

$$\text{For min mean cycle } X' \quad \sum_{(v,w) \in X'} c(v,w) = \mu(f) |X'|$$
$$\mu(f) \geq -\epsilon(f)$$
$$-\mu(f) \leq \epsilon(f)$$

$$2) \epsilon(f) \leq \nu(f)$$

Let  $X$  be the <sup>a</sup> min mean cycle in  $G_f$

$$\text{let } c'(v, w) = c(v, w) - \nu(f)$$

w/rt  $c'$ , all cycles have non-neg cost.

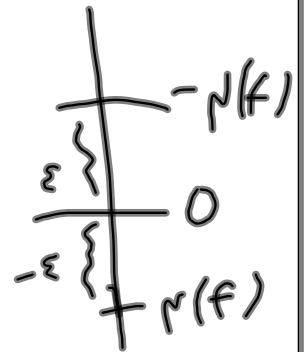
Pick  $v_1$ , let  $d(v)$  be the shortest path distance from  $v_1$  to  $v$  w/rt  $c'$ .

$$\Rightarrow d(w) \leq d(v) + c'(v, w) \\ = d(v) + c(v, w) - \nu(f).$$

$$\text{let } \pi(v) = d(v)$$

$$c^\pi(v, w) = c(v, w) - \pi(v) + \pi(w) \\ = c(v, w) + d(v) - d(w) \\ \geq \nu(f)$$

$$\epsilon(f) \leq -\nu(f).$$

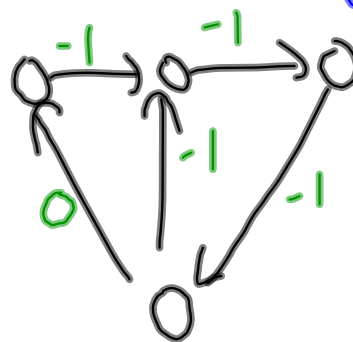
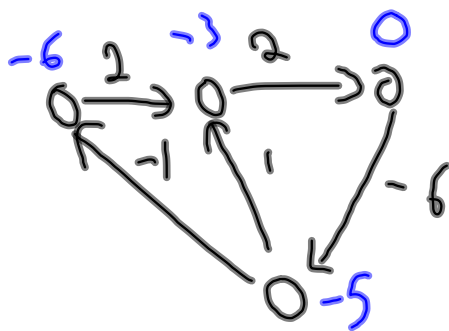


Lemma Let  $f$  be non-opt. Then  $\exists \pi$  s.t.

$$C^\pi(v, w) = p(f) = -\xi(f) \quad \forall (v, w) \in X$$

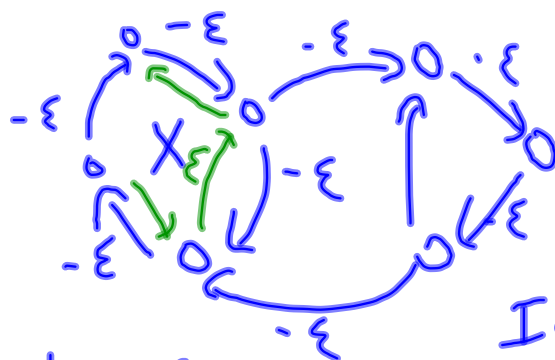
the min mean cycle in  $G_f$ .

(Previous proof  $w$ ) = on min. mean cycle.



Lemma Let  $f$  be a non-opt. flow, suppose we push flow around a min. mean cycle  $X$  to obtain  $f'$ . Then  $\epsilon(f') \leq \epsilon(f) = \epsilon$

Pf Draw  $G_f$  w/ costs  $c_\pi$



$c_\pi \geq \epsilon$   
 To get  $f'$ ,  
 I delete some  
 edges w/  $c_\pi = -\epsilon$ ,  
 I add some edges

Any new cycles in  $G_{f'}$  have at least one edge w/  $c_\pi = \epsilon > 0$ , and other edges  $c_\pi \geq \epsilon$ , and therefore cannot have mean value  $\leq -\epsilon$ .

—  $-\epsilon \cdot \epsilon \cdot \epsilon$

Lemma After  $m$  iterations of min. mean cycle  
 $\varepsilon(f') \leq \left(1 - \frac{1}{n}\right) \varepsilon(f)$ .  $f$  before  
 $f'$  after

Pf Fix  $\pi$   
 Case 1: All  $m$  cycles have all edges w/  
 $c^\pi(v,w) \leq 0$ .



Each iteration removes one edge from  $G_f$   
 & replaces it w/ an edge w/  $c^\pi(v,w) \geq 0$

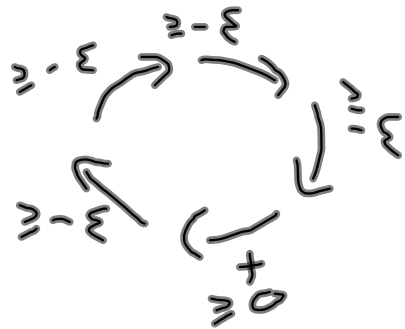
Total # of edges w/  $c^\pi(v,w) < 0$  has decreased by  $m$

$\therefore$  all edges in  $G_{f'}$  have  $c^\pi(v,w) \geq 0$ ,  
 $\therefore f'$  is optimal.



Case 2: Some min mean cycle has an edge  $w$

$$C^\pi(v, w) > 0$$



$$\begin{aligned} \mu(f) &\geq \frac{-\epsilon(|X|-1) + o(1)}{|X|} \\ &\geq -\epsilon \left(1 - \frac{1}{n}\right) \end{aligned}$$

$$\epsilon(f) \leq \epsilon \left(1 - \frac{1}{n}\right)$$