

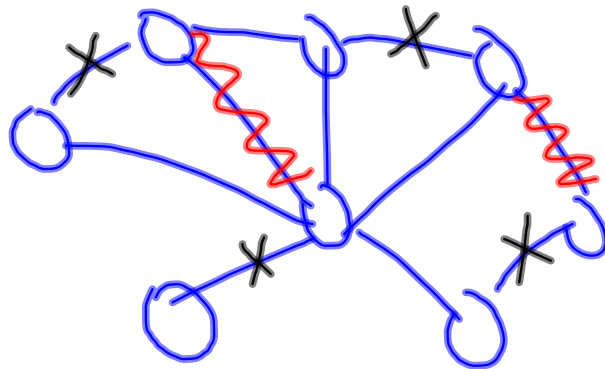
Matchings

Given a graph G , a matching M is a subset of the edges $M \subseteq E$, s.t. each vertex $v \in V$ is incident to at most one edge from M .

M is maximal

M is maximum

M is perfect



Variants

- weighted / unweighted
- bipartite / non-bipartite

maximum matching
max wt. matching

$$\max |M|$$

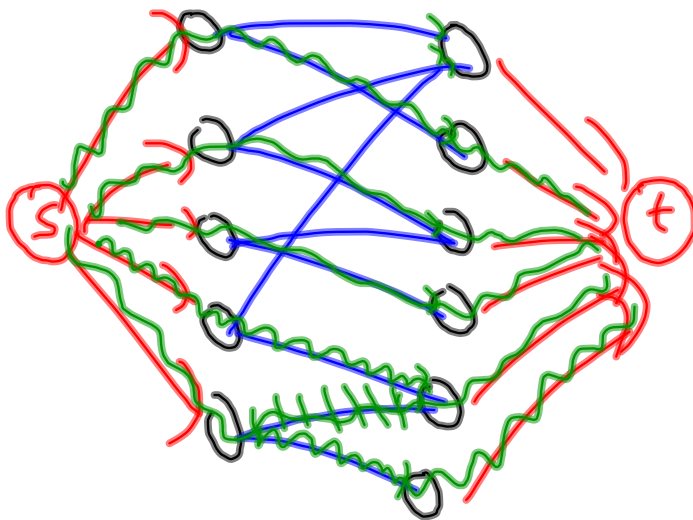
$$\max w(M) = \sum_{uv \in M} w(uv)$$

perfect matching

$$|M| = |V|/2$$

maximal matching is a matching M s.t. $\forall e \notin M$,
 $M \cup \{e\}$ is not a matching.

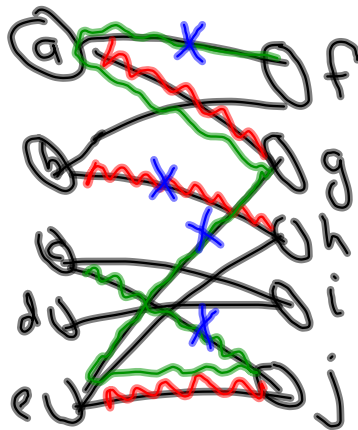
Bipartite Graphs



flows $\equiv$ matchings

$O(nm)$ time

a, g, e, j
matched.
b, c, d, f, h, i
free



alternating path

path in which the
edges alternate
between edges in M
& edges not in M .

ag
bh (aug)
fa, ag
c, je, eg, ga; a f (aug)

matched vertices v (v incident to an edge in M) $v \in M$
free vertices = $\neg$ (matched vertices)

augmenting path = alternating path in which the first & last
edges are not in M .

and it starts & ends at a free vertex

shortest aug. path

$$A \oplus B = (A \cup B) - (A \cap B)$$

Fact

M matching

P aug. path (relative to M)

Then $M \oplus P$ is a matching
& $|M \oplus P| = |M| + 1$

Fact

M matching, $P_1, \dots, P_k$ are vertex disjoint
aug. paths, then
 $M \oplus (P_1 \cup \dots \cup P_k)$ is a matching
of cardinality $|M| + k$

HK alg

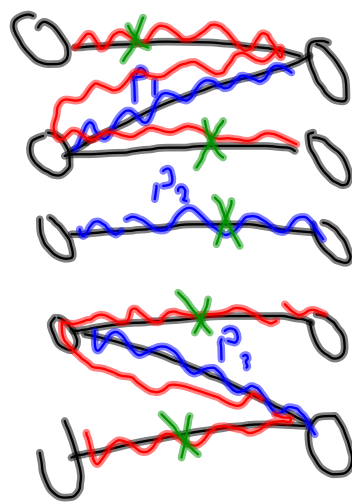
$M = \emptyset$

Repeat

- Let $P = (P_1 \cup \dots \cup P_k)$ be a maximal set of vertex disjoint shortest aug. paths wrt M .

• $M = M \oplus (P_1 \cup \dots \cup P_k)$
until $P = \emptyset$

$O(nm)$



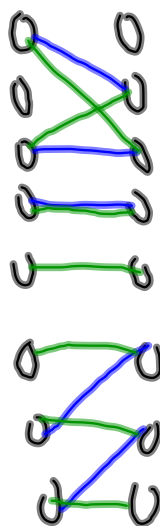
You can find a
maximal set of
vertex disjoint shortest
aug. paths in
 $O(m)$ time
via BFS.

After each iteration, the length
of the SAP. increases

Facts Given two matchings M, M^*

Consider $G' = (V, M \oplus M^*)$

S T
 M M^*

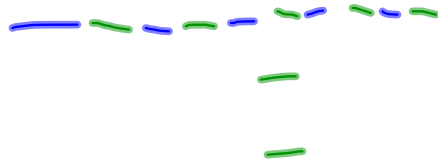


- $\deg(v) \leq 2$
- alternating paths & alternating cycles
- IF $|M| < |M^*|$, then G' has at least $|M^*| - |M|$ vertex disjoint augmenting paths w/rt M .

7 green edges

5 blue edges

Draw alt. paths & cycles
minimize the # of paths
starting and ending with green



Lemma let $l = \text{length of s.a.p. w.r.t } M$

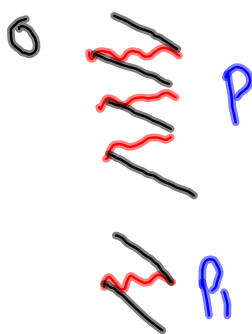
let $P_1 \dots P_k$ be a maximal set of vertex-disjoint ^{shortest-} augmenting paths

let $M' = M \oplus (P_1 \cup \dots \cup P_k)$

Let P be a s.a.p. w.r.t M'

Then $|P| > l$.

Pf case 1) P is vertex disjoint from $(P_1 \cup \dots \cup P_k) = S$
Then since S is a maximal set, P is not in S ,
so $|P| > l$.



$$\text{Case 2) } P \cap (P_1 \cup \dots \cup P_k) \neq \emptyset$$