

Mar 20-10:58 AM

Residual Graph

Residual

- Capacity is as for flow (now use $u_f(v,w)$ for residual capacity)
- If $(v,w) \in E$ and $(w,v) \in E_f$ then $c(w,v) = -c(v,w)$.

The diagram shows two nodes, v and w .
 - Top graph: Original flow. Edge $v \rightarrow w$ has $u=5, c=10$ and $f=3$. Residual capacity $u_f=2$ is shown below the edge.
 - Bottom graph: Residual graph. Edge $v \rightarrow w$ has $c=10$ and $u_f=3$. Edge $w \rightarrow v$ has $c=-10$.
 The word "Res." is written to the left of the bottom graph.

Mar 20-11:13 AM

Residual Graph

Residual

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Res.

Mar 20-11:13 AM

$f_{\text{optimal}} \iff G_f \text{ no neg. cost cycles}$

PF
 \Rightarrow if G_f has a neg. cost cycle then f is not optimal.

G_f

costs
 $u_f > 0$
 $-3 \cdot 1 + -1 \cdot 1 + 2 \cdot 1 = -2$
decreased the flow by 2
 $\Rightarrow f$ was not optimal

Mar 20-11:15 AM

If a G_f has no neg cycles, then f is optimal.

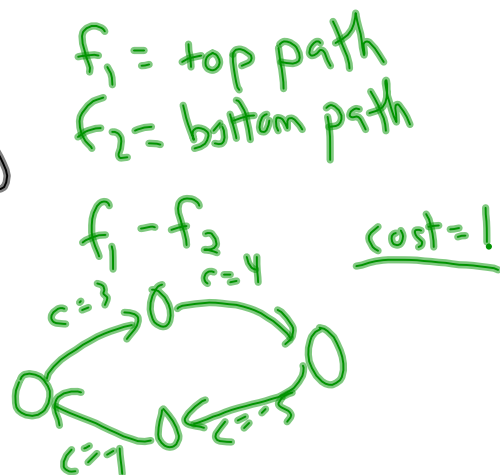
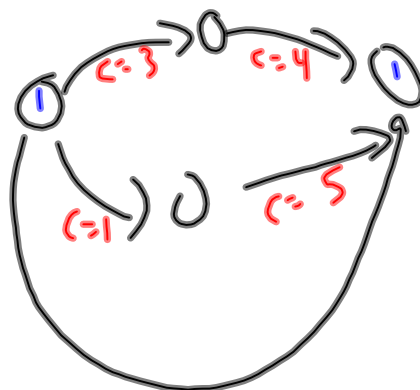
Think about f & f^* , the opt. flow.

can be decomposed $f - f^*$
into a collection of cycles, in G_f

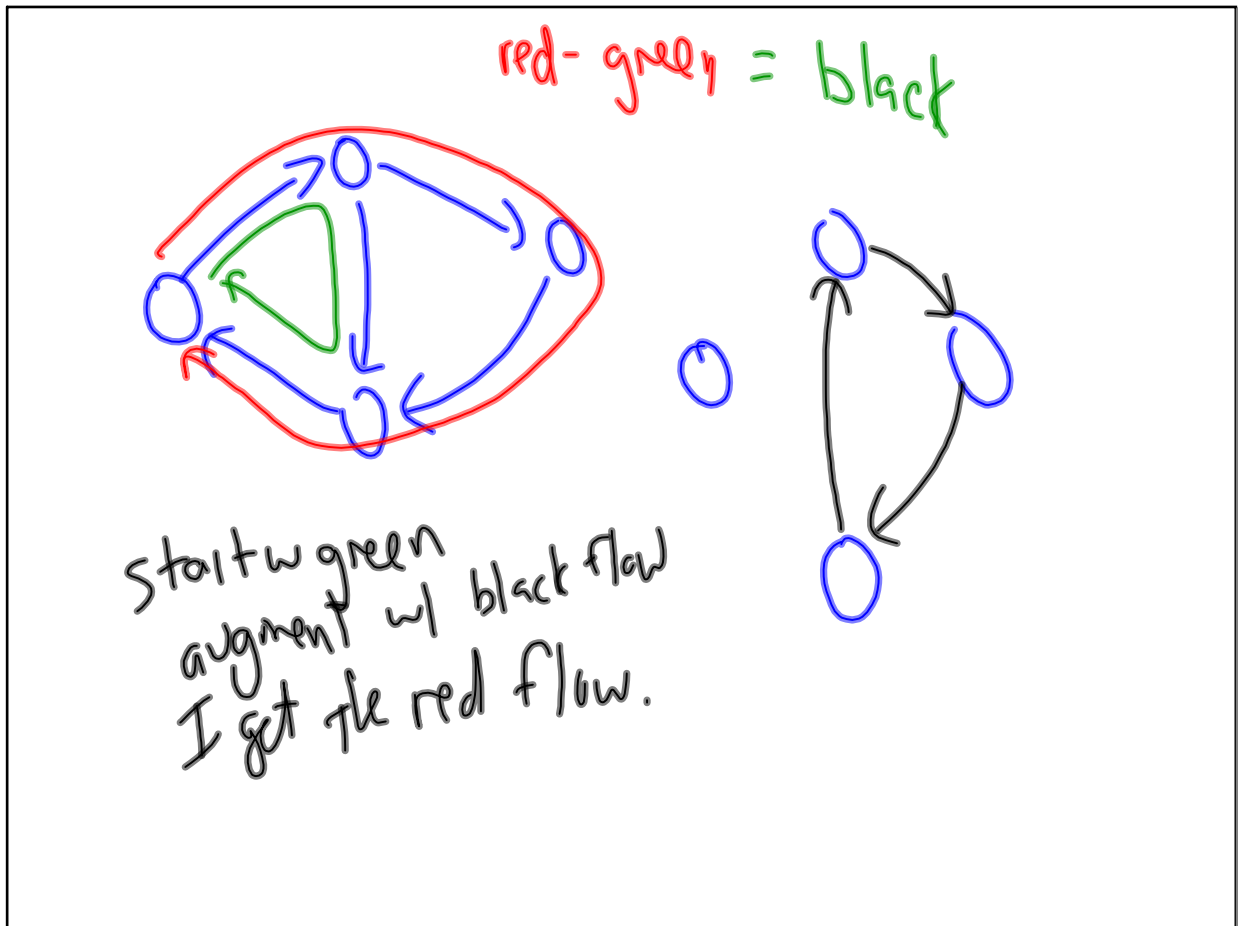
If G_f has no neg cycles, then
 $f - f^*$ is a collection of non-neg.
cycles $\rightarrow c(f) \leq c(f^*)$
 $\Rightarrow f$ is optimal.

Mar 20-11:20 AM

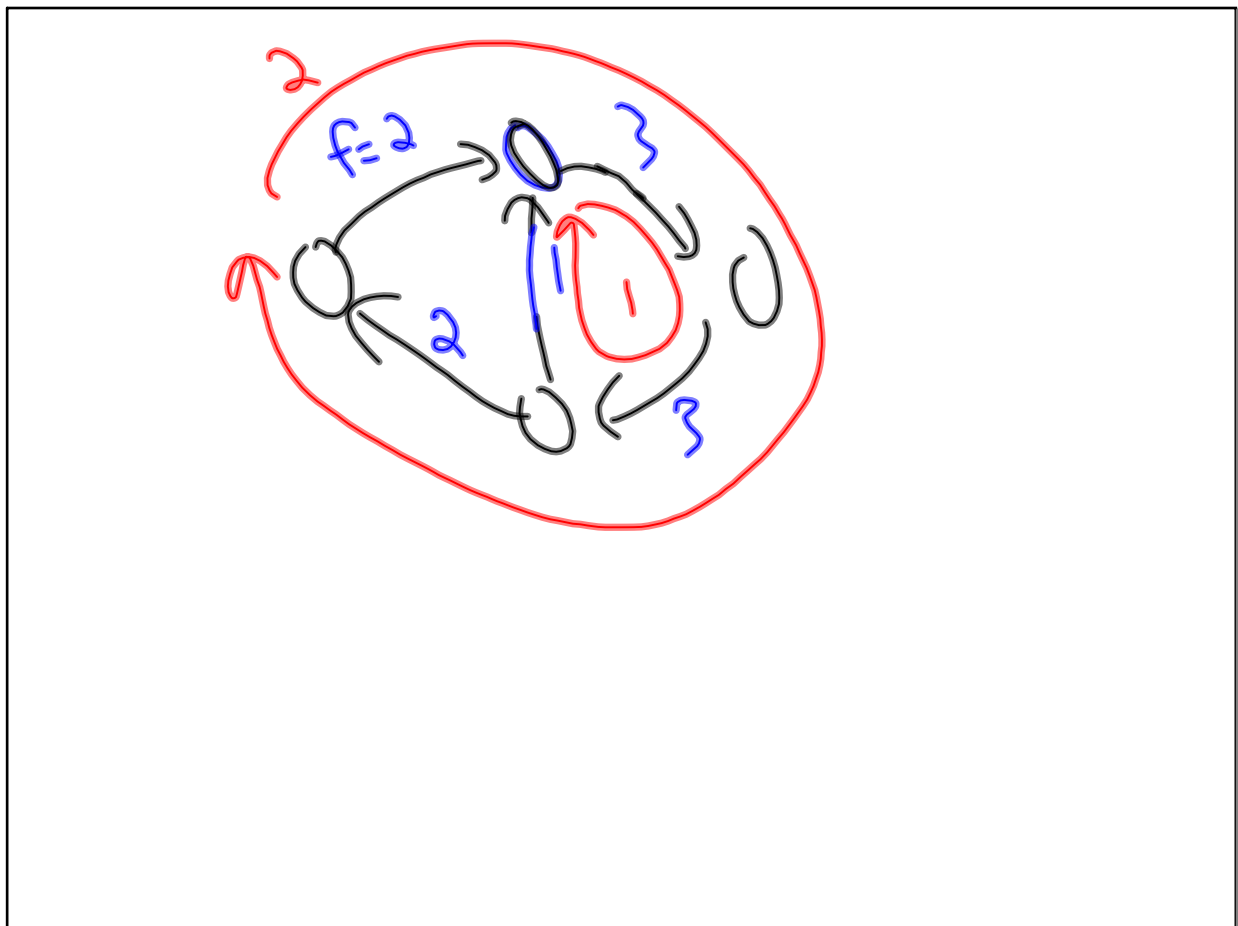
difference between two feas. mcf is a
collection of cycles



Mar 20-11:22 AM



Mar 20-11:27 AM



Mar 20-11:36 AM

Optimality 2: Reduced Cost Optimality

Reduced Cost Optimality: A feasible flow f is optimal iff there exists potentials π such that

$$c^\pi(v, w) \geq 0 \quad \forall (v, w) \in G_f$$

Pf Show G_f has no neg cost cycles \Leftrightarrow

$$\Leftrightarrow c^\pi(v, w) \geq 0 \quad \forall (v, w) \in G_f$$

$$\Rightarrow \text{for any cycle } X, \sum_{(v, w) \in X} c^\pi(v, w) \geq 0$$

$$\Rightarrow \sum_{(v, w) \in X} c(v, w) \geq 0.$$

Mar 20-11:46 AM

Optimality 2: Reduced Cost Optimality

G_f no neg. cost cycles \Rightarrow

Reduced Cost Optimality: A feasible flow f is optimal iff there exists potentials π such that

$$c^\pi(v, w) \geq 0 \quad \forall (v, w) \in G_f$$

\Rightarrow Pick a source s , compute s.p. dists $d(v)$ from s

Set $\pi(v) = -d(v)$

$$(d(w) \leq d(v) + c(v, w) \quad \forall (v, w) \in E)$$

$$c^\pi(v, w) = c(v, w) - \pi(v) + \pi(w)$$

$$= c(v, w) + d(v) - d(w) \geq 0 \quad \square$$

Mar 20-11:52 AM

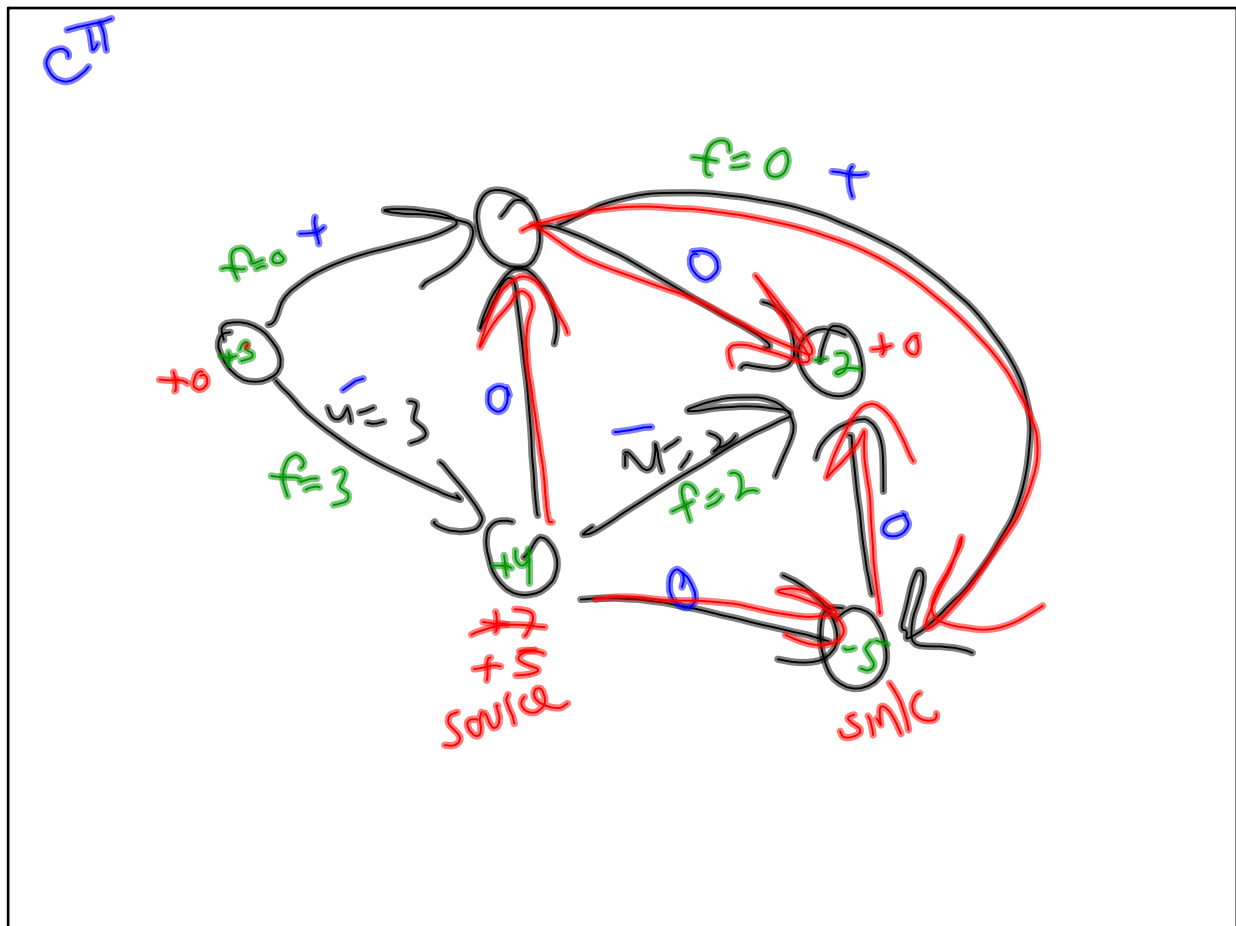
Optimality 3: Complimentary Slackness

A feasible flow f is optimal iff there exists potentials π such that for all edges $(v, w) \in G$

- if $c^\pi(v, w) > 0$ then $f(v, w) = 0$
- if $0 < f(v, w) < u(v, w)$ then $c^\pi(v, w) = 0$
- if $c^\pi(v, w) < 0$ then $f(v, w) = u(v, w)$.

$c(v, w) = (v, v) = 0$

Mar 20-12:07 PM



Mar 20-12:07 PM