

Claim
 If $p(v) = d(v, v)$ then

$$\bar{c}(v, w) = c(v, w) + p(v) - p(w) \geq 0$$

PF
 Because $d(\cdot)$ are sp. distances
 $d(w) \leq d(v) + c(v, w) \quad \forall (v, w)$

$$\Rightarrow c(v, w) + d(v) - d(w) \geq 0$$

$$\Rightarrow c(v, w) + p(v) - p(w) = \bar{c}(v, w) \geq 0$$

Feb 14-11:00 AM

Bidirectional Algorithm: Pitfalls

The algorithm is not as simple as it looks.

The searches meet at x , but x is not on the shortest path.

SOFSEM 07 8

Feb 14-11:16 AM

mmc.pdf - Adobe Reader

File Edit View Window Help

1 / 4 100%

Tools Comment

$$\mu(X) = \frac{\sum_{v,w \in X} c(v,w)}{|X|}$$

The minimum mean cycle is the cycle with smallest mean value.

$$\mu^* = \min_{\text{cycles } X} \mu(X)$$

11.00 x 8.50 in

Feb 14-11:45 AM

mmc.pdf - Adobe Reader

File Edit View Window Help

2 / 4 100%

Tools Comment

Computing a minimum mean cycle

Notes

- Computing the smallest value cycle is NP-hard
- The minimum mean cycle “approximates” the smallest value cycle.
- Choose v_1 as a “source”

Definition Let $d^k(v)$ be the length of a shortest directed walk from v_1 to v containing exactly k edges. (∞ if no such walk exists).

We can compute d^k for all k and via the recurrence:

$$d^k(w) = \min_{(v,w) \in E} d^{k-1}(v) + c(v,w)$$

- We initialize d^0 to ∞ for all vertices other than v_1 .

11.00 x 8.50 in

Feb 14-11:52 AM

10

$d^{10}(v_2) = 16$

| | d^0 | d^1 | d^2 | d^3 | d^4 | $\max(d^n(v) - d^k(v))/(n - k)$ |
|-------|----------|----------|----------|----------|----------|---------------------------------|
| v_1 | 0 | ∞ | ∞ | 20 | 14 | $7/2$ |
| v_2 | ∞ | 1 | ∞ | 12 | 6 | $5/3$ |
| v_3 | ∞ | 10 | 4 | ∞ | 15 | $11/2$ |
| v_4 | ∞ | ∞ | 12 | 6 | ∞ | ∞ |

11.00 x 8.50 in

Feb 14-11:55 AM

Pf let $\mu^* = \text{val of m.m.r.}$
 Assume $\mu^* = 0$ (case 1)
 \Rightarrow no neg. cycles
 \Rightarrow At least 1 0-cost cycle
 Compute shortest paths from v_1 , call these dist $d(v)$.
 Def. $c^d(v,w) = c(v,w) + d(v) - d(w)$.

Feb 14-12:06 PM

Think about graph w/ edges costs c^d

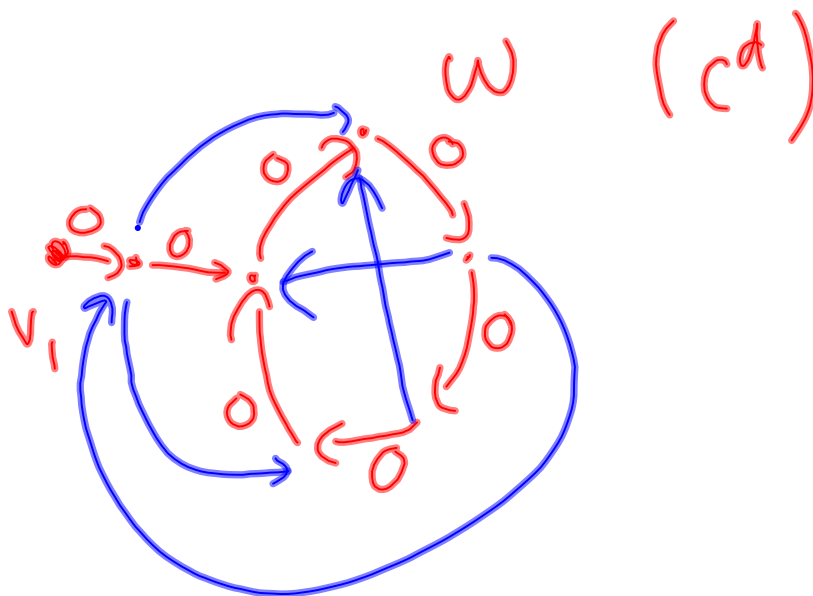
$$c^d(v,w) \geq 0 \quad \forall (v,w)$$

- let W be a 0-cost cycle w/r/t c .

$$c^d(v,w) = 0 \quad \forall (v,w) \in W$$

\neg If (v,w) is on a s.p. from v_i to some vertex, then $c^d(v,w) = 0$.

Feb 14-12:12 PM



Feb 14-12:16 PM