

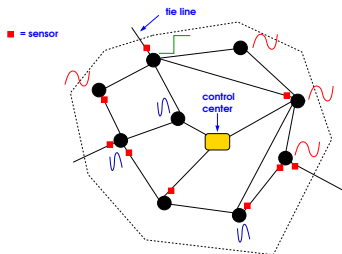
Real-time control of network physical structures to bypass complexity: Optimization, Stochastics and Structure Recognition

D. Bienstock, J. Blanchet, V. Goyal and G. Iyengar

Columbia University, Stanford University

July 2018

Real-time control of networked structures governed by physics



- Today: control enforces separation by time domain
e.g. in power grids: governor reaction (10^{-3} sec), AGC (sec), OPF (mins)
- Opportunity: fast sensors, algorithms
Challenges: “smart” loads, complex noise
- **Research Goals:**
 - ▶ Avoid separation
 - ▶ Quickly recognize system structure. **Time frame: seconds or less**
 - ▶ Quickly detect intrusion. **Time frame: seconds or very few minutes**

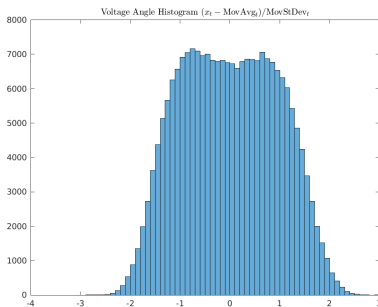
Challenges:

- How do we solve in near-real time hard problems that we cannot solve offline? E.g. nonconvex, polynomial optimization problems
 - To do: warm restart of ADMM-like methods for bilinear optimization.
- How do we handle noise/structure that we do not really understand?
 - Now doing: learning real-time correlation (or covariance) from noisy inputs. (NIPS Time series workshop 2017+)
- How do we combine first-order optimization with poorly understood “noise”?
 - Now doing: **Variance-aware** first-order optimization. (PSCC 2018)

Noise is not just noise

(We have **28 TB** of real data)

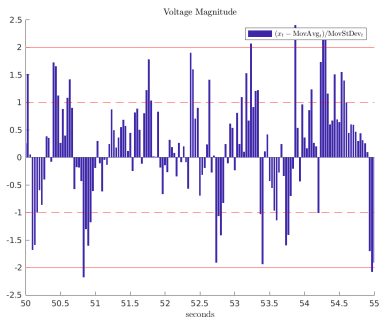
Voltage angle deviation histogram



Kolmogorov-Smirnoff gaussianity test strongly rejected, always

Noise is not just noise

From real time series, voltage magnitude deviations



Strong and nontrivial autocorrelation structure

Concrete problem: learning covariance in real time

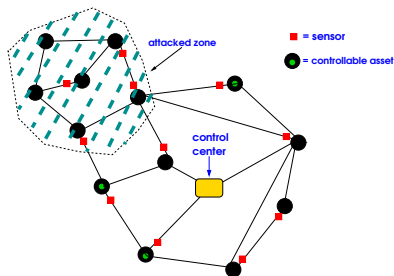
- **PCA:** principal component analysis.
 - ▶ Covariance of real-world data usually has low rank.
 - ▶ Fast PCA methods approximately capture the leading modes.
- **Streaming** PCA:
 - ▶ Old data gets stale
 - ▶ Cannot hold a lot of data
- **Non-stationary** regime for streaming PCA
 - ▶ Goal: detect change
 - ▶ Research question: what are fundamental computational limits?

with PhD student Apurv Shukla plus S. Kim (ex-LANL)

A streaming algorithm to identify PCA structure within time window

- Generative Model
 - ▶ Non-Stationary Spiked Covariance Model
- Sample Complexity:
 - ▶ Lower bound relating recovery error to number of samples
- Algorithm:
 - ▶ Two-phase iterative algorithm
 - ★ Phase-I: Iterative Eigenvector Computation
 - ★ Phase-II: Matrix Sketching
- **Theorems:** See NIPS paper and forthcoming paper

Application! Detecting intrusion through random physics



- Attacked zone is unknown to control center
- Attacker causes physical damage and alters sensor signals
- Defense:
 - ▶ Use controllable assets to alter **covariance** of physics
 - ▶ Changes unpredictable to attacker
 - ▶ Attacker (if aware of defense) can learn variance
 - ▶ But that takes time and sensor stream is continuous
 - ▶ So defender can **learn** the true covariance matrix

Distributionally Robust Optimization

Data-Driven Distributionally Robust Optimization

$$\min_{\theta} \sup_{P: D_c(P, P_n) \leq \delta} E_P(h(\theta, X)) \leftarrow \text{game formulation}$$

$$P_n(dx) = \frac{1}{n} \sum_{k=1}^n \delta_{\{X_k\}}(dx) = \text{Empirical Data}$$

$$E_{P_n}[h(\theta, X)] = \frac{1}{n} \sum_{k=1}^n h(\theta, X_k).$$

In words: Select the best response to model perturbations around the data (need to specify D_c).

Distributionally Robust Optimization (DRO)

- Extensive literature on DRO (Scarf (1958), Dupuis, James, Petersen (2000), Hansen and Sargent (2001), Ben-Tal, El Ghaoui, Nemirovski (2009), Delange & Ye (2010),...).
- Typical choices of $D(P, P_n)$

$$D(P, P_n) = E_P \left(\log \left(\frac{dP}{dP_n} \right) \right).$$

- Problem in data-driven setting: must preserve absolute continuity with respect to P_n .
- Choose $D(\cdot)$ based on optimal transport instead of divergence.
 - Works in practice and recovers exactly many machine learning estimators (e.g. Lasso, SVMs, adaptive ridge etc.)

Optimal Transport Metric and Wasserstein Distances

- Definition of Optimal Transport Discrepancy:

$$D_c(P, Q) = \min\{E_\pi(c(X, Y)) : \pi_X = P, \pi_Y = Q\}.$$

- Algorithmically $D_c(\cdot)$ is obtained by solving an LP

$$\begin{aligned} \min \sum_{x,y} c(x, y) \pi(x, y) \quad \text{subject to} \\ \sum_y \pi(x, y) = P(x), \quad \sum_x \pi(x, y) = Q(y) \\ \pi(x, y) \geq 0 \text{ for all } x, y. \end{aligned}$$

- Formulation includes Wasserstein and earth-mover's distance.

Applications of Data-Driven DRO

- We briefly explain fundamental connections to machine learning.
- Consider linear regression: Estimate $\beta_* \in R^m$ for model

$$Y_i = \beta_* X_i + e_i,$$

where $\{(Y_i, X_i)\}_{i=1}^n$ is a set of data points.

- Optimal Least Squares approach: Estimate β_* via

$$\min_{\beta} MSE(\beta) := \min_{\beta} n^{-1} \sum_{k=1}^n \left(Y_k - \beta^T X_k \right)^2$$

- We now apply the DRO formulation via optimal transport...

Connection to Sqrt-Lasso

Theorem (B., Kang, Murthy (2017)) Suppose that

$$c((x, y), (x', y')) = \begin{cases} \|x - x'\|_q^2 & \text{if } y = y' \\ \infty & \text{if } y \neq y' \end{cases}.$$

Then,

$$\max_{P: D_c(P, P_n) \leq \delta} E_P^{1/2} \left((Y - \beta^T X)^2 \right) = \sqrt{MSE(\beta)} + \sqrt{\delta} \|\beta\|_p.$$

Remark: This form of Lasso is called sqrt-Lasso (Belloni et al. (2011)).

Enhance Out-of-Sample Performance

- More general adversaries \rightarrow better decisions!
- Intuitive choice:

Generalized Mahalanobis: $c(x, y) = (y - x)^T A(x) (y - x)$,
 $A(\cdot)$ positive definite.

- We assume $\ell(\cdot)$ is a convex loss function to be minimized.
- Affine decision rules $\beta^T x$ (so empirical risk minimization involves minimizing $E_{P_n} [\ell(\beta^T X)]$ over β).
- We explore distributionally robust version of this problem.

A Class of DRO Problems with Affine Policies

Theorem (B., Fan, Murthy '18)

Under natural assumptions

$$\inf_{\beta \in \mathcal{R}^d} \sup_{P: D_c(P, P_n) \leq \delta} E_{P_n} [\ell(\beta^T X)] = \inf_{\beta, \lambda \geq 0} E_{P_n} [f(\beta, \lambda, X)],$$

for a tractable $f(\cdot)$ which is strongly convex in β, λ for $\delta \in [0, \delta_0]$ and some $\delta_0 > 0$. So, robust problem is not harder to solve than non-robust counterpart. Also, the worst case adversary can be computed.

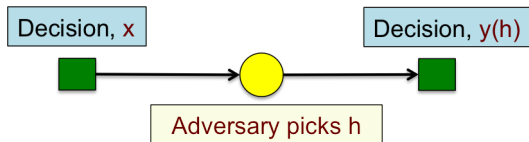
Consequence: Should robustify. Challenge: Calibrate the function $c(x, y)$.

Two-stage Adjustable Robust problem

$$z_{AR}(\mathcal{U}) = \min \mathbf{c}^T \mathbf{x} + \max_{\mathbf{h} \in \mathcal{U}} \min_{\mathbf{y}(\mathbf{h})} \mathbf{d}^T \mathbf{y}(\mathbf{h})$$

$$\mathbf{A}\mathbf{x} + \mathbf{B}\mathbf{y}(\mathbf{h}) \geq \mathbf{h}$$

$$\mathbf{x}, \mathbf{y}(\mathbf{h}) \in \mathbb{R}_+^n$$



- Demand uncertainty in unit commitment, facility location ...
- Hard to approximate within a factor better than $O(\log n / \log \log n)$ (Feige et al. 2007).

Approximate Solution Policies

- **Static robust solution:**

- ▶ Single solution (\mathbf{x}, \mathbf{y}) feasible for all scenarios.
- ▶ Easy to compute.
- ▶ Good approximation for symmetric sets (Bertsimas, G and Sun (2011)).
- ▶ Worst case performance bound is $\Omega(m)$.

- **Piecewise static policies:**

- ▶ Also known as K-adaptability policies.
- ▶ Divide uncertainty set into pieces and a static solution for each piece.
- ▶ Optimal pieces may be exponentially many (El Housni and G (2017))
- ▶ Even designing small number of optimal pieces is NP-hard. (Bertsimas and Caramanis (2012)).

Approximate Solution Policies

- **Affine policy:**

- ▶ $y(\mathbf{h}) = \mathbf{P}\mathbf{h} + \mathbf{q}$.
- ▶ Introduced by Ben-Tal et al. (2004)
- ▶ Can be computed efficiently and have good empirical performance.
- ▶ Optimal for simplex uncertainty sets
- ▶ Tight $O(\sqrt{m})$ -approximation for general sets (Bertsimas and G (2010)).

- **Piecewise affine policies**

- ▶ Chen and Zhang (2009), Bertsimas and Georghiou (2014), Bertsimas and Dunning (2014), Postek and Den Hertog (2016), Ben-Tal, El Housni and G (2016).
- ▶ Optimal for convex uncertainty sets (Zhen et al. (2016))
- ▶ Hard to compute the optimal pieces that may be exponentially many.

Current Work

- Characterization of performance of affine policies for important classes of uncertainty sets
- **Budget uncertainty sets:**

$$\mathcal{U} = \left\{ \mathbf{h} \in [0, 1]^m \mid \sum_{i=1}^m w_i h_i \leq \Gamma \right\}$$

- ▶ Very commonly used class of uncertainty sets.
 - ▶ Captures confidence interval sets and CLT sets.
 - ▶ Adjustable problem is $\Omega(\log n / \log \log n)$ -hard even for these sets
-
- **Intersection of Budget uncertainty sets**