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 $Ax \ge b$ 

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## • Optimization under parameter (data) uncertainty

- Ben-Tal and Nemirovsky, El Ghaoui et al
- Bertsimas et al
- Uncertainty is modeled by assuming that data is not known precisely, and will instead lie in known sets.
- Example: a coefficient  $a_i$  is uncertain. We allow  $a_i \in [l_i, u_i]$ .
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## "Tractable"

convex uncertainty models  $\rightarrow$  convex optimization techniques

polynomial-time algorithms

→ sacrifice model richness in favor of theoretical algorithm efficiency

→ in practice, SOCP not quite so "tractable"

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• Parameters:  $0 \le \gamma_1 \le \gamma_2 \le \ldots \le \gamma_K \le 1$ , integers  $0 \le n_i \le N_i$ ,  $1 \le i \le K$ for each asset *j*:  $\bar{\mu}_i =$  expected return

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- $\sum_{j} \mu_{j} \geq \Gamma \sum_{j} \bar{\mu}_{j}; \ \Gamma > 0$  a parameter
- (R. Tütüncü) For  $1 \le h \le H$ ,

• a set ("tier")  $T_h$  of assets, and a parameter  $\Gamma_h > 0$ 

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Benders' decomposition (= cutting-plane algorithm)

Generic problem:  $\min_{x \in X} \max_{d \in D} f(x, d)$ 

## $\rightarrow$ Maintain a finite subset $\tilde{D}$ of $\mathcal{D}$ (a "model")

## GAME

Implementor: solve min<sub>x∈X</sub> max<sub>d∈Ď</sub> f(x, d), with solution x\*

3 Adversary: solve max<sub>d∈D</sub> f(x\*, d), with solution d

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## Add d to D, and go to 1.

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- Decoupling of implementor and adversary yields considerably simpler, and smaller, problems
- Decoupling allows us to use more sophisticated uncertainty models
- If number of iterations is small, implementor's problem is a small "convex" problem

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 Most progress will be achieved in initial iterations – permits "soft" termination criteria

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## Implementor's problem

A convex quadratic program

At iteration *m*, solve

min  $\lambda x^T Q x - r$ 

Subject to:

 $Ax \ge b$ 

$$r \leq \mu_{(i)}^T \mathbf{x}, \ i = 1, \ldots, m$$

Here,  $\mu_{(1)}, \ldots, \mu_{(m)}$  are given return vectors

## Adversarial problem: A mixed-integer program

 $\mathbf{x}^* =$  given asset weights

min  $\sum_{j} \mathbf{x}_{j}^{*} \mu_{j}$ 

Subject to:

$$\begin{split} \bar{\mu}_{j}(1 - \sum_{i} \gamma_{i-1} y_{ij}) &\leq \mu_{j} \leq \bar{\mu}_{j}(1 - \sum_{i} \gamma_{i} y_{ij}) \quad \forall i \geq 1 \\ \sum_{i} y_{ij} \leq 1, \quad \forall j \quad \text{(each asset in at most one segment)} \\ n_{i} &\leq \sum_{j} y_{ij} \leq N_{i}, \quad 1 \leq i \leq K \quad \text{(segment cardinalities)} \\ \sum_{j \in T_{h}} \mu_{j} \geq \Gamma_{h} \sum_{j \in T_{h}} \bar{\mu}_{j}, \quad 1 \leq h \leq H \quad \text{(tier ineqs.)} \\ \mu_{j} \text{ free, } y_{ij} = 0 \text{ or } 1, \text{ all i, j} \end{split}$$

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Why the adversarial problem is "easy"

$$(K = no. of segments, H = no. of tiers)$$

**Theorem.** For every fixed **K** and **H**, and for every  $\epsilon > 0$ , there is an algorithm that finds a solution to the adversarial problem with optimality relative error  $\leq \epsilon$ , in time polynomial in  $\epsilon^{-1}$  and **n** (= no. of assets).
#### The simplest case

max  $\sum_j x_j^* \delta_j$ 

Subject to:

 $\sum_{j} \delta_{j} \leq \Gamma$   $\mathbf{0} \leq \delta_{j} \leq u_{j}y_{j}, y_{j} = \mathbf{0} \text{ or } \mathbf{1}, \text{ all } \mathbf{j}$  $\sum_{j} y_{j} \leq N$ 

... a cardinality constrained knapsack problem B. (1995), DeFarias and Nemhauser (2004)

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# The LP relaxation $x^* =$ given asset weights should (?) be tight

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#### **Robust problem:**

 $V \doteq \min \lambda x^T Q x - r$ Subject to:  $Ax \ge b$  $r \le \mu^T x$ ,  $\forall \mu$  achievable by adversary

Robust problem for relaxed adversary:

$$V \doteq \min \lambda x^{T} Q x - r$$
  
Subject to:  $A x \ge b$   
 $r \le \mu^{T} x$ ,  $\forall \mu$  achievable by **relaxed** adversary

 $V^* \ge V$ , perhaps  $V^* \approx V$ ,

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or,  $V \doteq \min \lambda x^T Q x - r$ Subject to:  $Ax \ge b$  $r \le \min \operatorname{minimum return}(x)$ 

**but**, minimum return(x) = min  $\sum_{i} x_{i}^{*} \mu_{i}$ 

Subject to:  $M_1\mu + M_2y \ge \Psi$ 

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duality: minimum return(x) = max  $\Psi^T \alpha$ 

Subject to:  $M_1^T \alpha = \mathbf{x}, \ M_2^T \alpha = \mathbf{0}, \ \alpha \ge \mathbf{0}$ 

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$$r - \Psi^T \alpha \leq 0$$

$$M_1^T \alpha - \mathbf{x} = \mathbf{0}, \ M_2^T \alpha = \mathbf{0}, \ \alpha \ge \mathbf{0}$$

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Let  $\hat{\mu} =$  optimal duals for (\*\*)

$$m{V}^* = m{min} \ \lambda m{x}^T m{Q} m{x} - m{r} \ m{S}$$
ubject to:  $m{A} m{x} \geq m{b} \ m{r} - \hat{\mu}^T m{x} \leq m{0}$ 

 $(\mathbf{r} - \boldsymbol{\mu}^{\mathsf{T}} \mathbf{x} \leq \mathbf{0}, \forall \boldsymbol{\mu}$  available to strict adversary)

**Problem:** Find  $\mu$  available to strict adversary, and with  $\mu \approx \hat{\mu}$ 

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$$V^* \doteq \min \lambda x^T Q x - r$$
  
Subject to:  $Ax \ge b$   
 $r - \Psi^T \alpha \le 0$   
(\*\*)  $M_1^T \alpha - x = 0$   
 $M_2^T \alpha = 0, \ \alpha \ge 0$ 

Let  $\hat{\mu}$  = optimal duals for (\*\*)

$$m{V}^* = \min \ \lambda m{x}^T m{Q} m{x} - m{r}$$
  
Subject to:  $m{A} m{x} \ge m{b}$   
 $m{r} - \hat{\mu}^T m{x} \le m{0}$ 

 $(\mathbf{r} - \boldsymbol{\mu}^T \mathbf{x} \leq \mathbf{0}, \forall \boldsymbol{\mu} \text{ available to strict adversary})$ 

**Problem:** Find  $\mu$  available to strict adversary, and with  $\mu \approx \hat{\mu}$ 

# Benders' algorithm with strengthening

**Step 1.** Solve relaxed robust problem; answer =  $\hat{\mu}$ 

**Step 2.** Solve MIP to obtain vector  $\vec{\mu}$  which is legal for histogram model, and with  $\vec{\mu} \approx \hat{\mu}$ 

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**Step 3.** Run Benders beginning with the cut  $\mathbf{r} - \boldsymbol{\mu}^T \mathbf{x} \leq \mathbf{0}$ 

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# Alternate algorithm?

**Step 1.** Solve relaxed robust problem, let  $\hat{\mu}$  be the min-max return vector

**Step 2.** Find a cut  $\alpha^T \mu \le \alpha_0$ , that separates  $\hat{\mu}$  from the convex hull of vectors available to the strict adversary

**Step 3.** Add  $\alpha^T \hat{\mu} \leq \alpha_0$  to the definition of the adversarial problem, and go to 1.

# Example: 2464 assets, 152-factor model. CPU time: 300 seconds No Strengthening – straight Benders

10 segments (a: "heavy tail") 6 tiers: the top five deciles lose at most 10% each, total loss < 5%



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### Same run

# 2464 assets, 152 factors; 10 segments, 6 tiers



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# Summary of average problems with 3-4 segments, 2-3 tiers

	columns	rows	iterations	time	imp. time	adv. time
				(sec.)		
1	500	20	47	1.85	1.34	0.46
2	500	20	3	0.09	0.01	0.03
3	703	108	1	0.29	0.13	0.04
4	499	140	3	3.12	2.65	0.05
5	499	20	19	0.42	0.21	0.17
6	1338	81	7	0.45	0.17	0.08
7	2019	140	8	41.53	39.6	0.36
8	2443	153	2	12.32	9.91	0.07
9	2464	153	111	100.81	60.93	36.78

	time	bigQP	bigMIP	iters	impT	advT	01vars
Α	327.04	2.52	211.72	135	12.27	100.24	5000
С	29.32	3.01	9.35	27	1.02	15.76	4990
F	74.06	13.57	15.96	27	2.47	41.42	13380
<b>G</b> *	681.12	_	_	19	64.7	615.54	20190
I	124.82	93.38	22.58	1	4.17	2.46	24640

Table: Heavy-tailed instances, 10 segments, 6 tiers, tol. =  $1.0e^{-03}$ 

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	500	500	499	<b>499</b> <sup>b</sup>	703 *	1338	2443
error	× 20	× 20	× 20	× 140	× 108	× 81	× 153
5.0e <sup>-2</sup>	214.53	14.81	144.86	122.53	11.77	274.64	140.29
1.0e <sup>-2</sup>	223.21	15.49	144.86	122.53	14.66	356.98	140.29
5.0e <sup>-3</sup>	254.73	16.03	162.41	126.63	34.16	363.84	140.29
1.0e <sup>-3</sup>	300.88	35.23	183.12	157.49	64.61	469.75	140.29
5.0e <sup>-4</sup>	361.20	37.92	216.52	167.40	73.87	598.94	140.29

Table: Convergence time on heavy-tailed instances, 10 segments, 6 tiers

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# What is the impact of the uncertainty model

All runs on the same data set with 1338 columns and 81 rows

- 1 segment: (200, 0.5) robust random return = 4.57, 157 assets
- 2 segments: (200, 0.25), (100, 0.5) robust random return = 4.57, 186 assets
- 2 segments: (200, 0.2), (100, 0.6) robust random return = 3.25, 213 assets
- 2 segments: (200, 0.1), (100, 0.8) robust random return = 1.50, 256 assets
- 1 segment: (100, 1.0) robust random return = 1.24, 281 assets

# The implementor chooses a vector x\* of assets

- 2 The adversary chooses a probability distribution P for the returns vector
- 3 A random returns vector  $\mu$  is drawn from P

 $\rightarrow$  Implementor wants to choose  $x^*$  so as to minimize **value-at-risk** (conditional value at risk, etc.)

Erdogan and Iyengar (2004), Calafiore and Campi (2004)  $\rightarrow$  We want to model *correlated* errors in the returns

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 $\rightarrow$  We want to model *correlated* errors in the returns

# **Uncertainty set**

### Given a vector **x**\* of assets, the adversary

- Chooses a vector  $w \in \mathbb{R}^n$  (n = no. of assets) with  $0 \le w_j \le 1$  for all j.
- ② Chooses a random variable  $0 \le \delta \le 1$
- ightarrow Random return:  $\mu_j = ar{\mu}_j$  (1  $-\delta w_j$ ) ( $ar{\mu} =$  nominal returns).

**Definition (Rockafellar and Uryasev):** Given reals  $\nu$  and  $0 \le \theta \le 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \ge 0$  such that

 $Prob(\nu - \mu^T \mathbf{x}^* \geq \rho) \geq \theta$ 

#### $\rightarrow$ The adversary wants to maximize VaR

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VaR Definition

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VaR Definition

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#### **Uncertainty set**

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- $\rightarrow$  Random return:  $\mu_j = \bar{\mu}_j (1 \delta w_j)$  ( $\bar{\mu} =$  nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \le \theta \le 1$  the *conditional value-at-risk* of  $x^*$  equals

 $E(\nu - \mu^T \mathbf{x}^* \mid \nu - \mu^T \mathbf{x}^* \geq \rho)$  where  $\rho = VaR$ 

 $\rightarrow$  The adversary wants to maximize CVaR

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# $\rightarrow$ Random return<sub>j</sub> = $\bar{\mu}_j (1 - \delta w_j)$ where $0 \le w_j \le 1 \forall j$ , and $0 \le \delta \le 1$ is a random variable.

A discrete distribution:

- We are given **fixed** values  $0 = \delta_0 \le \delta_2 \le ... \le \delta_K = 1$ example:  $\delta_i = \frac{i}{K}$
- Adversary chooses  $\pi_i = \operatorname{Prob}(\delta = \delta_i), 0 \le i \le K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^I \leq \pi_i \leq \pi_i^u$  (and possibly other constraints)
- Tier constraints: for sets ("tiers")  $T_h$  of assets,  $1 \le h \le H$ , we require:

 $E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h$  (given)

or,  $(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$ 

Definition

 $\rightarrow$  Random return<sub>i</sub> =  $\bar{\mu}_i (1 - \delta w_i)$  where  $0 \le w_i \le 1 \forall j$ , and  $0 < \delta < 1$  is a random variable.

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VaR Defi

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 *E*(δ ∑<sub>j∈T<sub>h</sub></sub> *w<sub>j</sub>*) ≤ Γ<sub>h</sub> (given)

or,  $(\sum_i \delta_i \pi_i) \sum_{j \in \mathcal{T}_h} w_j \leq \Gamma_h$ 

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VaR Defin

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VaR Defi

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# Robust optimization problem (VaR case): Given $0 < \epsilon$ , min V

Subject to:

 $\lambda \mathbf{x}^{\mathsf{T}} \mathbf{Q} \mathbf{x} - \mu^{\mathsf{T}} \mathbf{x} \leq \mathbf{v} * + \epsilon$ 

 $Ax \ge b$ 

 $V \geq VaR^{max}(x)$ 

Here,  $\mathbf{v}^* \doteq \min \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \mu^T \mathbf{x}$ 

Subject to:

 $Ax \ge b$ 

# Robust optimization problem (VaR case): Given $0 < \epsilon$ , min V

Subject to:

 $\lambda \mathbf{x}^{\mathsf{T}} \mathbf{Q} \mathbf{x} - \mu^{\mathsf{T}} \mathbf{x} \leq \mathbf{v} * + \epsilon$ 

 $\mathbf{A}\mathbf{x} \geq \mathbf{b}$ 

 $V \geq VaR^{max}(x)$ 

**Theorem:** The problem can be reduced to K SOCPs. K = number of points in discrete distribution

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## Adversarial problem – a nonlinear MIP

Recall: random return<sub>j</sub>  $\mu_j = \bar{\mu}_j (1 - \delta w_j)$ where  $\delta = \delta_i$  (given) with probability  $\pi_i$  (chosen by adversary),  $0 \le \delta_0 \le \delta_1 \le \ldots \le \delta_K = 1$  and  $0 \le w$ 

$$\begin{split} \min_{\pi,w,V} \min_{1 \le i \le k} V_i \\ \text{Subject to} \\ \mathbf{0} \le w_j \le 1, \text{ all j, } \pi_i^I \le \pi_i \le \pi_i^u, \text{ all i,} \\ \sum_i \pi_i = 1, \\ V_i = \sum_j \bar{\mu}_j (1 - \delta_i w_j) \mathbf{x}_j^*, \text{ if } \pi_i + \pi_{i+1} + \dots + \pi_K \ge \theta \\ V_i = M \text{ (large), otherwise} \\ (\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \le \Gamma_h, \text{ for each tier h} \end{split}$$

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 $\min_{\pi, w, V} \min_{1 \le i \le k} V_i$ 

Subject to

 $\begin{array}{l} \mathbf{0} \leq \mathbf{w}_{j} \leq \mathbf{1}, \ \text{all j}, \ \pi_{i}^{I} \leq \pi_{i} \leq \pi_{i}^{u}, \ \text{all i}, \\ \sum_{i} \pi_{i} = \mathbf{1}, \\ V_{i} = \sum_{j} \bar{\mu}_{j} (\mathbf{1} - \delta_{i} \mathbf{w}_{j}) \mathbf{x}_{j}^{*}, \ \text{ if } \ \pi_{i} + \pi_{i+1} + \ldots + \pi_{K} \geq \theta \\ V_{i} = M \ (large), \ \text{ otherwise} \\ (\sum_{i} \delta_{i} \pi_{i}) \sum_{i \in T_{h}} \mathbf{w}_{i} \leq \Gamma_{h}, \ \text{ for each tier h} \end{array}$ 

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$$\min_{\pi, w, V} \min_{1 \le i \le k} V_i$$

Subject to

$$\begin{array}{l} \mathbf{0} \leq \mathbf{w}_{j} \leq \mathbf{1}, \ \text{all j, } \ \pi_{i}^{I} \leq \pi_{i} \leq \pi_{i}^{u}, \ \text{all i,} \\ \sum_{i} \pi_{i} = \mathbf{1}, \end{array}$$

$$\begin{array}{l} \mathbf{V}_{i} = \sum_{j} \bar{\mu}_{j} (\mathbf{1} - \delta_{i} \mathbf{w}_{j}) \mathbf{x}_{j}^{*}, \ \text{if } \ \pi_{i} + \pi_{i+1} + \ldots + \pi_{K} \geq \theta \\ \mathbf{V}_{i} = \mathbf{M} \ (\textbf{large}), \ \text{otherwise} \end{array}$$

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$$\min_{\pi, w, V} \min_{1 \le i \le k} V_i$$

Subject to

$$\begin{array}{l} \mathbf{0} \leq \mathbf{w}_{j} \leq \mathbf{1}, \ \text{all j, } \ \pi_{i}^{I} \leq \pi_{i} \leq \pi_{i}^{U}, \ \text{all i,} \\ \sum_{i} \pi_{i} = \mathbf{1}, \end{array}$$

$$\begin{array}{l} \mathbf{V}_{i} = \sum_{j} \bar{\mu}_{j} (\mathbf{1} - \delta_{i} \mathbf{w}_{j}) \mathbf{x}_{j}^{*}, \ \text{if } \pi_{i} + \pi_{i+1} + \ldots + \pi_{K} \geq \theta \\ \mathbf{V}_{i} = \mathbf{M} \ (large), \ \text{otherwise} \end{array}$$

$$(\sum_{i} \delta_{i} \pi_{i}) \sum_{j \in \mathcal{T}_{h}} \mathbf{w}_{j} \leq \Gamma_{h}, \ \text{for each tier h} \end{array}$$

#### Definition

# Approximation

 $(\sum_{i} \delta_{i} \pi_{i}) \sum_{j \in T_{h}} w_{j} \leq \Gamma_{h}, \quad \text{for each tier h} \quad (*)$ 

Let  $\mathit{N} > \mathit{0}$  be an integer. For  $1 \leq \mathit{k} \leq \mathit{N}$ , write

- $rac{k}{N}\sum_{j\in T_h}w_j\leq \Gamma_h + M(1-z_{hk}),$  where
- $\mathbf{z}_{hk} = \mathbf{1}$  if  $rac{k-1}{N} < \sum_i \delta_i \pi_i \leq rac{k}{N}$
- *z<sub>hk</sub>* = 0 otherwise
- $\sum_{k} z_{hk} = 1$

and *M* is large

**Lemma.** Under reasonable conditions, replacing (\*) with this system creates an error of order  $O(\frac{1}{N})$ 

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#### Definition

# Approximation

 $(\sum_{i} \delta_{i} \pi_{i}) \sum_{j \in T_{h}} w_{j} \leq \Gamma_{h}, \text{ for each tier h} \quad (*)$ Let N > 0 be an integer. For  $1 \leq k \leq N$ , write  $\frac{k}{N} \sum_{j \in T_{h}} w_{j} \leq \Gamma_{h} + M(1 - z_{hk}), \text{ where}$  $z_{hk} = 1 \text{ if } \frac{k-1}{N} < \sum_{i} \delta_{i} \pi_{i} \leq \frac{k}{N}$  $z_{hk} = 0 \text{ otherwise}$  $\sum_{k} z_{hk} = 1$ 

and *M* is large

**Lemma.** Under reasonable conditions, replacing (\*) with this system creates an error of order  $O(\frac{1}{N})$ 

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#### Implementor's problem for Benders approach, at iteration r:

Here,  $\delta_{i(t)}$  and  $w^{(t)}$  are the adversary's output at iteration t < r.

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## Implementor's problem for Benders approach, at iteration r:

## min V

Subject to:

$$\lambda \mathbf{x}^{T} \mathbf{Q} \mathbf{x} - \mu^{T} \mathbf{x} \leq (1 + \epsilon) \mathbf{v}^{*}$$
  

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$
  

$$\mathbf{V} \geq \nu - \sum_{j} \bar{\mu}_{j} \left( 1 - \delta_{i(t)} \mathbf{w}_{j}^{(t)} \right) \mathbf{x}_{j}, \quad t = 1, 2, \dots, r - 1$$

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#### But we can do better:

At iteration t, the adversary produces a probability distribution  $\pi^{(t)}$ and a vector  $\mathbf{w}^{(t)}$ 

and the cut is: 
$$\mathbf{V} \ge \nu - \sum_{j} \bar{\mu}_{j} \left( \mathbf{1} - \delta_{i(t)} \mathbf{w}_{j}^{(t)} \right) \mathbf{x}_{j}$$
  
here,  $i(t)$  is smallest s.t.  $\sum_{i \ge i(t)} \pi_{i}^{(t)} \ge \theta$ 

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#### But we can do better:

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$$m{V} \geq 
u \, - \, \sum_j ar{\mu}_j \left(m{1} - \delta_{i(t)} m{w}_j^{(t)}
ight) m{x}_j$$

How about a cut that is valid for *every* w s.t.  $(\pi^{(t)}, w)$  is feasible for the adversary?

## We want an expression for

```
\begin{split} \min \sum_{j} \bar{\mu}_{j} (\mathbf{1} - \delta_{i(t)} \mathbf{w}_{j}) \mathbf{x}_{j} \\ \text{Subject to} \\ (\sum_{i} \delta_{i} \pi_{i}^{(t)}) \sum_{j \in \mathcal{T}_{h}} \mathbf{w}_{j} \leq \Gamma_{h}, \quad \text{for each tier h} \end{split}
```

→ Use LP duality

 $\rightarrow$  The implementor's problem will gain new variables and rows at each iteration

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### Typical convergence behavior – simple Benders



## Heavy-tailed instances, $\theta = .05$

Heavy tail, proportional error (100 points):



## Heavy-tailed instances, $\theta = .05$

Heavy tail, proportional error (100 points):



Heavy-tailed instances,  $\theta = .05$ , K = 100

VaR	A	D	E	F	G	I
time	1.98	5.02	2.47	2.03	26.51	38.32
iters	2	2	2	2	2	2
impt	0.25	2.25	0.54	1.07	14.09	19.90
advt	1.26	1.14	1.32	0.24	2.17	1.47
adj $ au$	2.8e <sup>-04</sup>	2.4e <sup>-04</sup>	3.0e <sup>-04</sup>	2.5e <sup>-04</sup>	4.7e <sup>-05</sup>	2.1e <sup>-04</sup>

CVaR	A	D	E	F	G	
time	7.10	14.11	6.23	11.45	33.13	88.43
iters	2	2	2	2	2	3
impt	0.16	1.72	1.18	0.66	9.56	52.13
advt	6.72	10.67	4.74	10.33	12.2	23.85
gap	9.8e <sup>-04</sup>	2.2e <sup>-05</sup>	7.3e <sup>-05</sup>	5.1e <sup>-05</sup>	3.2e <sup>-05</sup>	1.3e <sup>-04</sup>
apperr	2.3e <sup>-04</sup>	2.2e <sup>-05</sup>	2.4e <sup>-04</sup>	1.6e <sup>-05</sup>	1.0e <sup>-04</sup>	2.2e <sup>-04</sup>

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#### Impact of tail probability

"confidence level" =  $1 - \theta$ 



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# Impact of suboptimality target

#### Fix $\theta = 0.95$ but vary suboptimality criterion



#### Experiment: sensitivity of model to parameters

Adversary chooses 
$$\pi_i = P(\delta = \delta_i), \quad \pi_i^l \leq \pi_i \leq \pi_i^u$$

Experiment: choose  $\Delta \ge 0$ , and solve robust problems for

 $\pi_i \leftarrow \max\{\pi_i^l - \Delta, \mathbf{0}\}, \ \pi_i^l \leftarrow \pi_i^u + \Delta$ 

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## VaR and CVaR as a function of data errors:



## $(N = 10^4 \text{ for VaR case})$

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