

# Experiments with Robust Optimization

Daniel Bienstock

Columbia University  
New York

ISMP 2006, Rio

# Robust Optimization

- Optimization under parameter (data) uncertainty
- Ben-Tal and Nemirovsky, El Ghaoui et al
- Bertsimas et al
- Uncertainty is modeled by assuming that data is not known precisely, and will instead lie in known sets.
- Example: a coefficient  $a_i$  is uncertain. We allow  $a_i \in [l_i, u_i]$ .
- Typically, a minimization problem becomes a min-max problem.

# Robust Optimization

- Optimization under parameter (data) uncertainty
- Ben-Tal and Nemirovsky, El Ghaoui et al
- Bertsimas et al
- Uncertainty is modeled by assuming that data is not known precisely, and will instead lie in known sets.
- Example: a coefficient  $a_i$  is uncertain. We allow  $a_i \in [l_i, u_i]$ .
- Typically, a minimization problem becomes a min-max problem.

# Robust Optimization

- Optimization under parameter (data) uncertainty
- Ben-Tal and Nemirovsky, El Ghaoui et al
- Bertsimas et al
- Uncertainty is modeled by assuming that data is not known precisely, and will instead lie in known sets.
- Example: a coefficient  $a_i$  is uncertain. We allow  $a_i \in [l_i, u_i]$ .
- Typically, a minimization problem becomes a min-max problem.

# Robust Optimization

- Optimization under parameter (data) uncertainty
- Ben-Tal and Nemirovsky, El Ghaoui et al
- Bertsimas et al
- Uncertainty is modeled by assuming that data is not known precisely, and will instead lie in known sets.
- Example: a coefficient  $a_i$  is uncertain. We allow  $a_i \in [l_i, u_i]$ .
- Typically, a **minimization** problem becomes a **min-max** problem.

# Robust Optimization

- Optimization under parameter (data) uncertainty
- Ben-Tal and Nemirovsky, El Ghaoui et al
- Bertsimas et al
- Uncertainty is modeled by assuming that data is not known precisely, and will instead lie in known sets.
- Example: a coefficient  $\mathbf{a}_i$  is uncertain. We allow  $\mathbf{a}_i \in [l_i, u_i]$ .
- Typically, a minimization problem becomes a min-max problem.

# Robust Optimization

- Optimization under parameter (data) uncertainty
- Ben-Tal and Nemirovsky, El Ghaoui et al
- Bertsimas et al
- Uncertainty is modeled by assuming that data is not known precisely, and will instead lie in known sets.
- Example: a coefficient  $\mathbf{a}_i$  is uncertain. We allow  $\mathbf{a}_i \in [l_i, u_i]$ .
- Typically, a **minimization** problem becomes a **min-max** problem.

## Example: Linear Programs with Row-Wise uncertainty

Ben-Tal and Nemirovsky, 1999 (also: Soyster (1973))

$$\min c^t x$$

Subject to:

$$Ax \geq b \quad \text{for all } A \in \mathcal{U}$$

$$x \in X$$

$\mathcal{U}$  = uncertainty set

→ the  $i^{\text{th}}$  row of  $A$  belongs to an ellipsoidal set  $\mathcal{E}_i$

$$\text{e.g. } \sum_j \alpha_{ij}^2 (\mathbf{a}_{ij} - \bar{\mathbf{a}}_{ij})^2 \leq 1$$

→ can be solved using SOCP techniques

# Other forms of optimization under uncertainty

- Stochastic programming
- Adversarial queueing, online optimization
- “Risk-aware” optimization
- Optimization of utility functions as a substitute for handling infeasibilities

# Other forms of optimization under uncertainty

- Stochastic programming
- Adversarial queueing, online optimization
- “Risk-aware” optimization
- Optimization of utility functions as a substitute for handling infeasibilities

# Other forms of optimization under uncertainty

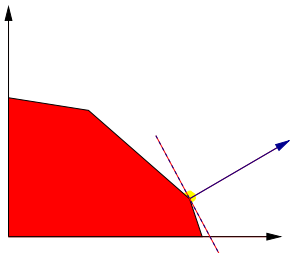
- Stochastic programming
- Adversarial queueing, online optimization
- “Risk-aware” optimization
- Optimization of utility functions as a substitute for handling infeasibilities

# Other forms of optimization under uncertainty

- Stochastic programming
- Adversarial queueing, online optimization
- “Risk-aware” optimization
- Optimization of utility functions as a substitute for handling infeasibilities

## Scenario I: Stability

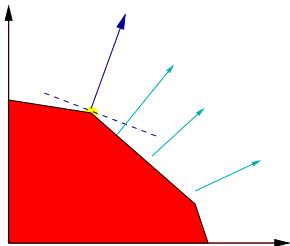
Data is fairly accurate, though possibly noisy – small errors are possible



→ Idiosyncratic decisions and small changes in data could have major impact

## Scenario I: Stability

Data is fairly accurate, though possibly noisy – small errors are possible



→ Idiosyncratic decisions and small changes in data could have major impact

## Scenario II: Hedging

Significant, but within order-of-magnitude, data uncertainty

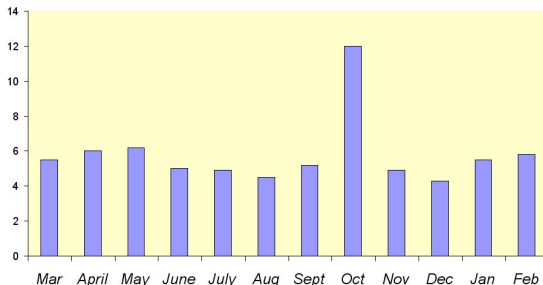
### Example:

A certain parameter,  $\alpha$ , is volatile. Its long-term average is **1.5** but it we could expect changes of the order of **.3**.

- Possibly more than just noise
- Could use deviations to our advantage, especially if there are several uncertain parameters that act “correlated”
- Are we guarding against risk or are we **hedging**?

## Scenario III: Insurance

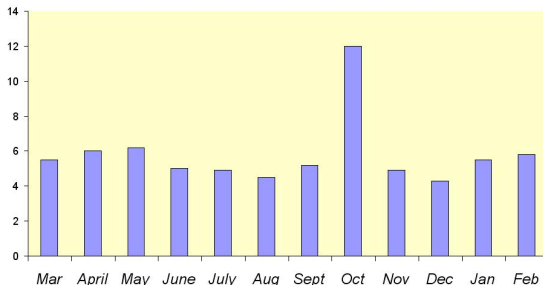
Real world data can exhibit undesirable and unexpected behavior



- Classical goal: how can we protect without becoming too risk averse
- Need to clearly spell out desired tradeoff between risk and performance
- Magnitude and geometry of risk are not the same

## Scenario III: Insurance

Real world data can exhibit undesirable and unexpected behavior



- Classical goal: how can we protect without becoming too risk averse
- Need to clearly spell out desired tradeoff between risk and performance
- **Magnitude** and **geometry** of risk are not the same

## Application: Portfolio Optimization

- $\min \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \boldsymbol{\mu}^T \mathbf{x}$

Subject to:

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

- $\boldsymbol{\mu}$  = vector of “returns”,  $\mathbf{Q}$  = “covariance” matrix
- $\mathbf{x}$  = vector of “asset weights”
- $\mathbf{A} \mathbf{x} \geq \mathbf{b}$ : general linear constraints
- $\lambda \geq 0$  = “risk-aversion” multiplier

## Application: Portfolio Optimization

- $\min \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \boldsymbol{\mu}^T \mathbf{x}$

Subject to:

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

- $\boldsymbol{\mu}$  = vector of “returns”,  $\mathbf{Q}$  = “covariance” matrix
- $\mathbf{x}$  = vector of “asset weights”
- $\mathbf{A} \mathbf{x} \geq \mathbf{b}$ : general linear constraints
- $\lambda \geq 0$  = “risk-aversion” multiplier

## Application: Portfolio Optimization

- $\min \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \boldsymbol{\mu}^T \mathbf{x}$

Subject to:

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

- $\boldsymbol{\mu}$  = vector of “returns”,  $\mathbf{Q}$  = “covariance” matrix
- $\mathbf{x}$  = vector of “asset weights”
- $\mathbf{A} \mathbf{x} \geq \mathbf{b}$ : general linear constraints
- $\lambda \geq 0$  = “risk-aversion” multiplier

## Application: Portfolio Optimization

- $\min \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \boldsymbol{\mu}^T \mathbf{x}$

Subject to:

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

- $\boldsymbol{\mu}$  = vector of “returns”,  $\mathbf{Q}$  = “covariance” matrix
- $\mathbf{x}$  = vector of “asset weights”
- $\mathbf{A} \mathbf{x} \geq \mathbf{b}$ : general linear constraints
- $\lambda \geq 0$  = “risk-aversion” multiplier

## Application: Portfolio Optimization

- $\min \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \boldsymbol{\mu}^T \mathbf{x}$

Subject to:

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

- $\boldsymbol{\mu}$  = vector of “returns”,  $\mathbf{Q}$  = “covariance” matrix
- $\mathbf{x}$  = vector of “asset weights”
- $\mathbf{A} \mathbf{x} \geq \mathbf{b}$ : general linear constraints
- $\lambda \geq \mathbf{0}$  = “risk-aversion” multiplier

## Robust Portfolio Optimization

Goldfarb and Iyengar, 2001

→  $\mathbf{Q}$  and  $\mu$  are uncertain

### Robust Problem

$$\min_{\mathbf{x}} \left\{ \max_{\mathbf{Q} \in \mathcal{Q}} \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \min_{\mu \in \mathcal{E}} \mu^T \mathbf{x} \right\}$$

Subject to:

$$\sum_j \mathbf{x}_j = \mathbf{1}, \quad \mathbf{x} \geq \mathbf{0}$$

→ When  $\mathcal{Q}$  is an ellipsoid and  $\mathcal{E}$  is a product of intervals the robust problem can be solved as an **SOCP**

## Robust Portfolio Optimization

### A different uncertainty model

→ Want to model that deviations of the returns  $\mu_j$  from their nominal values are **rare** but could be **significant**

#### A simple example

- Parameters:  $\mathbf{0} \leq \gamma \leq \mathbf{1}$ , integer  $\mathbf{N} \geq \mathbf{0}$ ,  
for each asset  $j$ :  
 $\bar{\mu}_j =$  expected return,  $\mathbf{0} \leq \delta_j$  small (possibly zero)
- Well-behaved asset  $j$ :  $\bar{\mu}_j - \delta_j \leq \mu_j \leq \bar{\mu}_j + \delta_j$
- Misbehaving asset  $j$ :  $(1 - \gamma)\bar{\mu}_j \leq \mu_j \leq \bar{\mu}_j$
- At most  $N$  assets misbehave

## Robust Portfolio Optimization

### A different uncertainty model

→ Want to model that deviations of the returns  $\mu_j$  from their nominal values are **rare** but could be **significant**

#### A simple example

- Parameters:  $\mathbf{0} \leq \gamma \leq \mathbf{1}$ , integer  $\mathbf{N} \geq \mathbf{0}$ ,  
for each asset  $j$ :  
 $\bar{\mu}_j =$  expected return,  $\mathbf{0} \leq \delta_j$  small (possibly zero)
- **Well-behaved asset  $j$ :**  $\bar{\mu}_j - \delta_j \leq \mu_j \leq \bar{\mu}_j + \delta_j$
- **Misbehaving asset  $j$ :**  $(1 - \gamma)\bar{\mu}_j \leq \mu_j \leq \bar{\mu}_j$
- At most  $N$  assets misbehave

## Robust Portfolio Optimization

### A different uncertainty model

→ Want to model that deviations of the returns  $\mu_j$  from their nominal values are **rare** but could be **significant**

#### A simple example

- Parameters:  $\mathbf{0} \leq \gamma \leq \mathbf{1}$ , integer  $\mathbf{N} \geq \mathbf{0}$ ,  
for each asset  $j$ :  
 $\bar{\mu}_j =$  expected return,  $\mathbf{0} \leq \delta_j$  small (possibly zero)
- Well-behaved asset  $j$ :**  $\bar{\mu}_j - \delta_j \leq \mu_j \leq \bar{\mu}_j + \delta_j$
- Misbehaving asset  $j$ :**  $(1 - \gamma)\bar{\mu}_j \leq \mu_j \leq \bar{\mu}_j$
- At most  $N$  assets misbehave

## Robust Portfolio Optimization

### A different uncertainty model

→ Want to model that deviations of the returns  $\mu_j$  from their nominal values are **rare** but could be **significant**

#### A simple example

- Parameters:  $\mathbf{0} \leq \gamma \leq \mathbf{1}$ , integer  $\mathbf{N} \geq \mathbf{0}$ ,  
for each asset  $j$ :  
 $\bar{\mu}_j =$  expected return,  $\mathbf{0} \leq \delta_j$  small (possibly zero)
- **Well-behaved asset  $j$ :**  $\bar{\mu}_j - \delta_j \leq \mu_j \leq \bar{\mu}_j + \delta_j$
- **Misbehaving asset  $j$ :**  $(1 - \gamma)\bar{\mu}_j \leq \mu_j \leq \bar{\mu}_j$
- **At most  $\mathbf{N}$  assets misbehave**

## A more comprehensive setting

- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$

for each asset  $j$ :  $\bar{\mu}_j =$  expected return

- between  $n_i$  and  $N_i$  assets  $j$  satisfy:

$$(1 - \gamma_i)\bar{\mu}_j \leq \mu_j \leq (1 - \gamma_{i-1})\bar{\mu}_j, \text{ for each } i \geq 1 \quad (\gamma_0 = \mathbf{0})$$

## A more comprehensive setting

- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$

for each asset  $j$ :  $\bar{\mu}_j =$  expected return

- between  $n_i$  and  $N_i$  assets  $j$  satisfy:

$$(\mathbf{1} - \gamma_i)\bar{\mu}_j \leq \mu_j \leq (\mathbf{1} - \gamma_{i-1})\bar{\mu}_j, \text{ for each } i \geq \mathbf{1} \quad (\gamma_0 = \mathbf{0})$$

## A more comprehensive setting

### Alternative

- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$

for each asset  $j$ :  $\bar{\mu}_j =$  expected return,  $\bar{\delta}_j =$  “standard deviation” of return

- between  $n_i$  and  $N_i$  assets  $j$  satisfy:

$$\bar{\mu}_j - \gamma_i \delta_j \leq \mu_j \leq \bar{\mu}_j - \gamma_{i-1} \delta_j$$

## A more comprehensive setting

### Alternative

- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$

for each asset  $j$ :  $\bar{\mu}_j =$  expected return,  $\bar{\delta}_j =$  “standard deviation” of return

- between  $n_i$  and  $N_i$  assets  $j$  satisfy:

$$\bar{\mu}_j - \gamma_i \delta_j \leq \mu_j \leq \bar{\mu}_j - \gamma_{i-1} \delta_j$$

## A more comprehensive setting

- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$   
for each asset  $j$ :  $\bar{\mu}_j =$  expected return
- between  $n_i$  and  $N_i$  assets  $j$  satisfy:  
 $(1 - \gamma_i)\bar{\mu}_j \leq \mu_j \leq (1 - \gamma_{i-1})\bar{\mu}_j$
- $\sum_j \mu_j \geq \Gamma \sum_j \bar{\mu}_j$ ;  $\Gamma > \mathbf{0}$  a parameter
- (R. Tütüncü) For  $\mathbf{1} \leq h \leq H$ ,
  - a set ("tier")  $\mathcal{T}_h$  of assets, and a parameter  $\Gamma_h > 0$
 for each  $h$ ,  $\sum_{j \in \mathcal{T}_h} \mu_j \geq \Gamma_h \sum_{j \in \mathcal{S}_h} \bar{\mu}_j$

Note: only downwards changes are modeled

## A more comprehensive setting

- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$   
for each asset  $j$ :  $\bar{\mu}_j =$  expected return
- between  $n_i$  and  $N_i$  assets  $j$  satisfy:  
 $(\mathbf{1} - \gamma_i)\bar{\mu}_j \leq \mu_j \leq (\mathbf{1} - \gamma_{i-1})\bar{\mu}_j$
- $\sum_j \mu_j \geq \Gamma \sum_j \bar{\mu}_j$ ;  $\Gamma > \mathbf{0}$  a parameter
- (R. Tütüncü) For  $\mathbf{1} \leq h \leq H$ ,
  - a set (“tier”)  $T_h$  of assets, and a parameter  $\Gamma_h > \mathbf{0}$
 for each  $h$ ,  $\sum_{j \in T_h} \mu_j \geq \Gamma_h \sum_{j \in S_h} \bar{\mu}_j$

Note: only downwards changes are modeled

## A more comprehensive setting

- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$   
for each asset  $j$ :  $\bar{\mu}_j =$  expected return
- between  $n_i$  and  $N_i$  assets  $j$  satisfy:  
 $(\mathbf{1} - \gamma_i)\bar{\mu}_j \leq \mu_j \leq (\mathbf{1} - \gamma_{i-1})\bar{\mu}_j$
- $\sum_j \mu_j \geq \Gamma \sum_j \bar{\mu}_j$ ;  $\Gamma > \mathbf{0}$  a parameter
- (R. Tütüncü) For  $\mathbf{1} \leq h \leq H$ ,
  - a set ("tier")  $\mathcal{T}_h$  of assets, and a parameter  $\Gamma_h > 0$
 for each  $h$ ,  $\sum_{j \in \mathcal{T}_h} \mu_j \geq \Gamma_h \sum_{j \in \mathcal{S}_h} \bar{\mu}_j$

Note: only downwards changes are modeled

## A more comprehensive setting

- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$   
for each asset  $j$ :  $\bar{\mu}_j =$  expected return
- between  $n_i$  and  $N_i$  assets  $j$  satisfy:  
 $(\mathbf{1} - \gamma_i)\bar{\mu}_j \leq \mu_j \leq (\mathbf{1} - \gamma_{i-1})\bar{\mu}_j$
- $\sum_j \mu_j \geq \Gamma \sum_j \bar{\mu}_j$ ;  $\Gamma > \mathbf{0}$  a parameter
- (R. Tütüncü) For  $\mathbf{1} \leq h \leq H$ ,
  - a set (“tier”)  $T_h$  of assets, and a parameter  $\Gamma_h > 0$

for each  $h$ ,  $\sum_{j \in T_h} \mu_j \geq \Gamma_h \sum_{j \in S_h} \bar{\mu}_j$

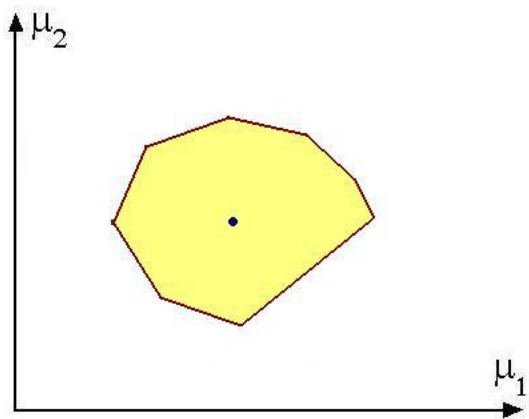
Note: only downwards changes are modeled

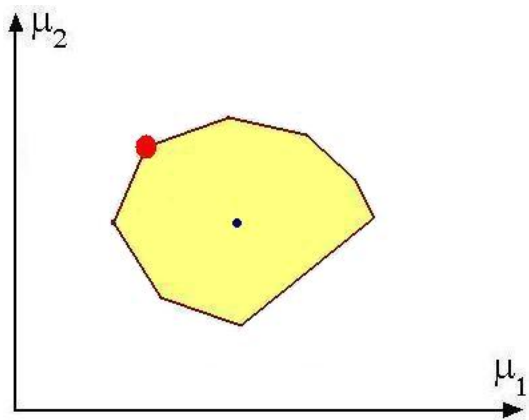
## A more comprehensive setting

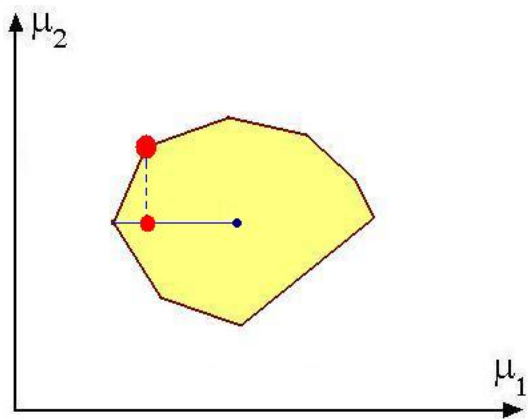
- Parameters:  $\mathbf{0} \leq \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_K \leq \mathbf{1}$ ,  
integers  $\mathbf{0} \leq n_i \leq N_i$ ,  $\mathbf{1} \leq i \leq K$   
for each asset  $j$ :  $\bar{\mu}_j =$  expected return
- between  $n_i$  and  $N_i$  assets  $j$  satisfy:  
 $(\mathbf{1} - \gamma_i)\bar{\mu}_j \leq \mu_j \leq (\mathbf{1} - \gamma_{i-1})\bar{\mu}_j$
- $\sum_j \mu_j \geq \Gamma \sum_j \bar{\mu}_j$ ;  $\Gamma > \mathbf{0}$  a parameter
- (R. Tütüncü) For  $\mathbf{1} \leq h \leq H$ ,
  - a set (“tier”)  $T_h$  of assets, and a parameter  $\Gamma_h > 0$

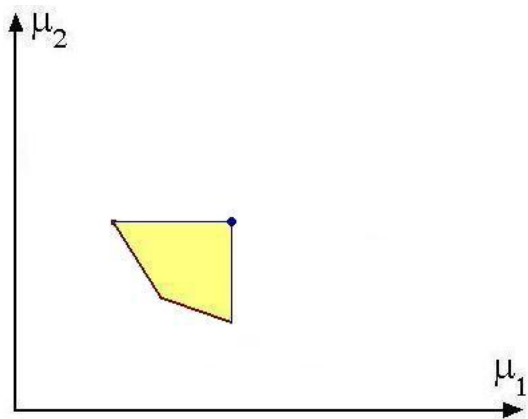
for each  $h$ ,  $\sum_{j \in T_h} \mu_j \geq \Gamma_h \sum_{j \in S_h} \bar{\mu}_j$

**Note:** only downwards changes are modeled









## General methodology:

Benders' decomposition (= cutting-plane algorithm)

Generic problem:  $\min_{x \in X} \max_{d \in \mathcal{D}} f(x, d)$

→ Maintain a **finite subset**  $\tilde{\mathcal{D}}$  of  $\mathcal{D}$  (a “model”)

### GAME

- 1 Implementor: solve  $\min_{x \in X} \max_{d \in \tilde{\mathcal{D}}} f(x, d)$ , with solution  $x^*$
- 2 Adversary: solve  $\max_{d \in \mathcal{D}} f(x^*, d)$ , with solution  $\tilde{d}$
- 3 Add  $\tilde{d}$  to  $\tilde{\mathcal{D}}$ , and go to 1.

## General methodology:

Benders' decomposition (= cutting-plane algorithm)

Generic problem:  $\min_{x \in X} \max_{d \in \mathcal{D}} f(x, d)$

→ Maintain a **finite subset**  $\tilde{\mathcal{D}}$  of  $\mathcal{D}$  (a “model”)

### GAME

- 1 **Implementor:** solve  $\min_{x \in X} \max_{d \in \tilde{\mathcal{D}}} f(x, d)$ , with solution  $x^*$
- 2 **Adversary:** solve  $\max_{d \in \mathcal{D}} f(x^*, d)$ , with solution  $\tilde{d}$
- 3 Add  $\tilde{d}$  to  $\tilde{\mathcal{D}}$ , and go to 1.

## General methodology:

Benders' decomposition (= cutting-plane algorithm)

Generic problem:  $\min_{x \in X} \max_{d \in \mathcal{D}} f(x, d)$

→ Maintain a **finite subset**  $\tilde{\mathcal{D}}$  of  $\mathcal{D}$  (a “model”)

### GAME

1 **Implementor:** solve  $\min_{x \in X} \max_{d \in \tilde{\mathcal{D}}} f(x, d)$ ,  
with solution  $x^*$

2 **Adversary:** solve  $\max_{d \in \mathcal{D}} f(x^*, d)$ , with solution  $\tilde{d}$

3 **Add**  $\tilde{d}$  to  $\tilde{\mathcal{D}}$ , and go to 1.

## General methodology:

Benders' decomposition (= cutting-plane algorithm)

Generic problem:  $\min_{x \in X} \max_{d \in \mathcal{D}} f(x, d)$

→ Maintain a **finite subset**  $\tilde{\mathcal{D}}$  of  $\mathcal{D}$  (a “model”)

### GAME

- 1 **Implementor:** solve  $\min_{x \in X} \max_{d \in \tilde{\mathcal{D}}} f(x, d)$ ,  
with solution  $x^*$
- 2 **Adversary:** solve  $\max_{d \in \mathcal{D}} f(x^*, d)$ , with solution  $\tilde{d}$
- 3 Add  $\tilde{d}$  to  $\tilde{\mathcal{D}}$ , and go to 1.

## General methodology:

Benders' decomposition (= cutting-plane algorithm)

Generic problem:  $\min_{x \in X} \max_{d \in \mathcal{D}} f(x, d)$

→ Maintain a **finite subset**  $\tilde{\mathcal{D}}$  of  $\mathcal{D}$  (a “model”)

### GAME

- 1 **Implementor:** solve  $\min_{x \in X} \max_{d \in \tilde{\mathcal{D}}} f(x, d)$ ,  
with solution  $x^*$
- 2 **Adversary:** solve  $\max_{d \in \mathcal{D}} f(x^*, d)$ , with solution  $\tilde{d}$
- 3 **Add**  $\tilde{d}$  to  $\tilde{\mathcal{D}}$ , and go to 1.

## Why this approach

- Decoupling of implementor and adversary yields considerably simpler, and smaller, problems
- Decoupling allows us to use more sophisticated uncertainty models
- If number of iterations is small, implementor's problem is a small "convex" problem
- Most progress will be achieved in initial iterations – permits "soft" termination criteria

## Why this approach

- Decoupling of implementor and adversary yields considerably simpler, and smaller, problems
- Decoupling allows us to use more sophisticated uncertainty models
- If number of iterations is small, implementor's problem is a small "convex" problem
- Most progress will be achieved in initial iterations – permits "soft" termination criteria

## Why this approach

- Decoupling of implementor and adversary yields considerably simpler, and smaller, problems
- Decoupling allows us to use more sophisticated uncertainty models
- If number of iterations is small, implementor's problem is a small "convex" problem
- Most progress will be achieved in initial iterations – permits "soft" termination criteria

## Why this approach

- Decoupling of implementor and adversary yields considerably simpler, and smaller, problems
- Decoupling allows us to use more sophisticated uncertainty models
- If number of iterations is small, implementor's problem is a small "convex" problem
- Most progress will be achieved in initial iterations – permits "soft" termination criteria

## Why this approach

- Decoupling of implementor and adversary yields considerably simpler, and smaller, problems
- Decoupling allows us to use more sophisticated uncertainty models
- If number of iterations is small, implementor's problem is a small "convex" problem
- Most progress will be achieved in initial iterations – permits "soft" termination criteria

## Implementor's problem

A convex quadratic program

At iteration  $m$ , solve

$$\min \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - r$$

Subject to:

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

$$r \leq \mu_{(i)}^T \mathbf{x}, \quad i = 1, \dots, m$$

Here,  $\mu_{(1)}, \dots, \mu_{(m)}$  are given return vectors

## Adversarial problem: A mixed-integer program

$\mathbf{x}^*$  = given asset weights

$$\min \sum_j \mathbf{x}_j^* \mu_j$$

Subject to:

$$\bar{\mu}_j(\mathbf{1} - \sum_i \gamma_{i-1} \mathbf{y}_{ij}) \leq \mu_j \leq \bar{\mu}_j(\mathbf{1} - \sum_i \gamma_i \mathbf{y}_{ij}) \quad \forall i \geq 1$$

$$\sum_i \mathbf{y}_{ij} \leq \mathbf{1}, \quad \forall j \quad (\text{each asset in at most one segment})$$

$$\mathbf{n}_i \leq \sum_j \mathbf{y}_{ij} \leq \mathbf{N}_i, \quad \mathbf{1} \leq i \leq \mathbf{K} \quad (\text{segment cardinalities})$$

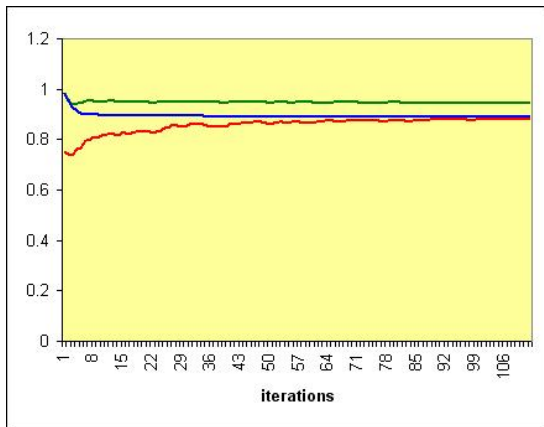
$$\sum_{j \in T_h} \mu_j \geq \Gamma_h \sum_{j \in T_h} \bar{\mu}_j, \quad \mathbf{1} \leq h \leq \mathbf{H} \quad (\text{tier ineqs.})$$

$$\mu_j \text{ free, } \mathbf{y}_{ij} = \mathbf{0} \text{ or } \mathbf{1}, \text{ all } i, j$$

2464 assets, 152 factors; total CPU time = 93.27 sec.

4 segments: (400, 0.01), (50, 0.05), (30, 0.2), (10, 0.3)

3 tiers: the three top deciles lose at most **10%** each



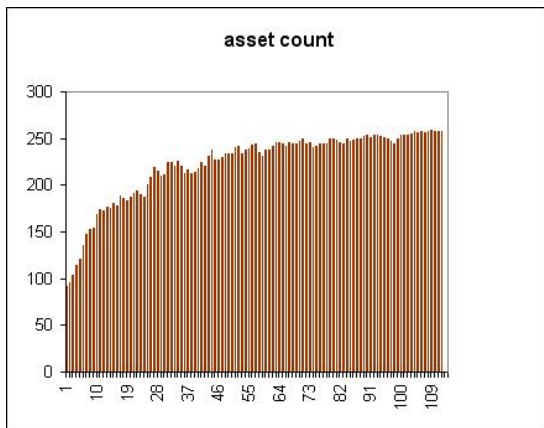
green = nominal return, blue = estimate, red = adversarial

## Same run

2464 assets, 152 factors; total CPU time = 93.27 sec.

4 segments: (400, 0.01), (50, 0.05), (30, 0.20), (10, 0.30)

3 tiers: the three top deciles lose at most **10%** each



## Summary

	columns	rows	iterations	time (sec.)	imp. time	adv. time
1	500	20	47	1.85	1.34	0.46
2	500	20	3	0.09	0.01	0.03
3	703	108	1	0.29	0.13	0.04
4	499	140	3	3.12	2.65	0.05
5	499	20	19	0.42	0.21	0.17
6	1338	81	7	0.45	0.17	0.08
7	2019	140	8	41.53	39.6	0.36
8	2443	153	2	12.32	9.91	0.07
9	2464	153	111	100.81	60.93	36.78

## Why the adversarial problem is “easy”

(  $\mathbf{x}^*$  = given asset weights)

$$\min \sum_j \mathbf{x}_j^* \mu_j$$

Subject to:

$$\bar{\mu}_j(\mathbf{1} - \sum_i \gamma_{i-1} \mathbf{y}_{ij}) \leq \mu_j \leq \bar{\mu}_j(\mathbf{1} - \sum_i \gamma_i \mathbf{y}_{ij})$$

$$\sum_i \mathbf{y}_{ij} \leq \mathbf{1}, \quad \forall j \quad (\text{each asset in at most one segment})$$

$$\mathbf{n}_i \leq \sum_j \mathbf{y}_{ij} \leq \mathbf{N}_i, \quad \forall i \quad (\text{segment cardinalities})$$

$$\sum_{j \in \mathcal{T}_h} \mu_j \geq \Gamma_h (\sum_{j \in \mathcal{T}_h} \bar{\mu}_j), \quad \forall h \quad (\text{tier inequalities})$$

$$\mu_j \text{ free, } \mathbf{y}_{ij} = 0 \text{ or } 1, \text{ all } i, j$$

## Why the adversarial problem is “easy”

(  $K$  = no. of segments,  $H$  = no. of tiers)

**Theorem.** For every fixed  $K$  and  $H$ , and for every  $\epsilon > 0$ , there is an algorithm that finds a solution to the adversarial problem with optimality relative error  $\leq \epsilon$ , in time polynomial in  $\epsilon^{-1}$  and  $n$  (= no. of assets).

## The simplest case

$$\max \sum_j x_j^* \delta_j$$

Subject to:

$$\sum_j \delta_j \leq \Gamma$$

$$0 \leq \delta_j \leq u_j y_j, \quad y_j = 0 \text{ or } 1, \text{ all } j$$

$$\sum_j y_j \leq N$$

... a *cardinality constrained knapsack problem*

B. (1995), DeFarias and Nemhauser (2004)

## What is the impact of the uncertainty model

All runs on the same data set with 1338 columns and 81 rows

- 1 segment: (200, 0.5)  
robust random return = **4.57**, **157** assets
- 2 segments: (200, 0.25), (100, 0.5)  
robust random return = **4.57**, **186** assets
- 2 segments: (200, 0.2), (100, 0.6)  
robust random return = **3.25**, **213** assets
- 2 segments: (200, 0.1), (100, 0.8)  
robust random return = **1.50**, **256** assets
- 1 segment: (100, 1.0)  
robust random return = **1.24**, **281** assets

## Ambiguous chance-constrained models

- 1 The implementor chooses a vector  $x^*$  of assets
- 2 The adversary chooses a *probability distribution*  $P$  for the returns vector
- 3 A random returns vector  $\mu$  is drawn from  $P$

→ Implementor wants to choose  $x^*$  so as to minimize **value-at-risk** (conditional value at risk, etc.)

Erdogan and Iyengar (2004), Calafiore and Campi (2004)

→ We want to model *correlated* errors in the returns

## Ambiguous chance-constrained models

- 1 The implementor chooses a vector  $x^*$  of assets
- 2 The adversary chooses a *probability distribution*  $P$  for the returns vector
- 3 A random returns vector  $\mu$  is drawn from  $P$

→ Implementor wants to choose  $x^*$  so as to minimize **value-at-risk** (conditional value at risk, etc.)

Erdogan and Iyengar (2004), Calafiore and Campi (2004)

→ We want to model *correlated* errors in the returns

## Ambiguous chance-constrained models

- 1 The implementor chooses a vector  $x^*$  of assets
- 2 The adversary chooses a *probability distribution*  $P$  for the returns vector
- 3 A random returns vector  $\mu$  is drawn from  $P$

→ Implementor wants to choose  $x^*$  so as to minimize **value-at-risk** (conditional value at risk, etc.)

Erdogan and Iyengar (2004), Calafiore and Campi (2004)

→ We want to model *correlated* errors in the returns

## Ambiguous chance-constrained models

- 1 The implementor chooses a vector  $x^*$  of assets
- 2 The adversary chooses a *probability distribution*  $P$  for the returns vector
- 3 A random returns vector  $\mu$  is drawn from  $P$

→ Implementor wants to choose  $x^*$  so as to minimize **value-at-risk** (conditional value at risk, etc.)

Erdogan and Iyengar (2004), Calafiore and Campi (2004)

→ We want to model *correlated* errors in the returns

## Ambiguous chance-constrained models

- 1 The implementor chooses a vector  $x^*$  of assets
- 2 The adversary chooses a *probability distribution*  $P$  for the returns vector
- 3 A random returns vector  $\mu$  is drawn from  $P$

→ Implementor wants to choose  $x^*$  so as to minimize **value-at-risk** (conditional value at risk, etc.)

Erdogan and Iyengar (2004), Calafiore and Campi (2004)

→ We want to model *correlated* errors in the returns

## Ambiguous chance-constrained models

- 1 The implementor chooses a vector  $x^*$  of assets
- 2 The adversary chooses a *probability distribution*  $P$  for the returns vector
- 3 A random returns vector  $\mu$  is drawn from  $P$

→ Implementor wants to choose  $x^*$  so as to minimize **value-at-risk** (conditional value at risk, etc.)

Erdogan and Iyengar (2004), Calafiore and Campi (2004)

→ We want to model *correlated* errors in the returns

## Uncertainty set

Given a vector  $x^*$  of assets, the adversary

- 1 Chooses a vector  $w \in R^n$  ( $n = \text{no. of assets}$ ) with  $0 \leq w_j \leq 1$  for all  $j$ .
- 2 Chooses a random variable  $0 \leq \delta \leq 1$

→ Random return:  $\mu_j = \bar{\mu}_j (1 - \delta w_j)$  ( $\bar{\mu}$  = nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \leq \theta \leq 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \geq 0$  such that

$$\text{Prob}(\nu - \mu^T x^* \geq \rho) \geq \theta$$

→ The adversary wants to maximize VAR

## Uncertainty set

Given a vector  $x^*$  of assets, the adversary

- 1 Chooses a vector  $w \in R^n$  ( $n = \text{no. of assets}$ ) with  $0 \leq w_j \leq 1$  for all  $j$ .
- 2 Chooses a random variable  $0 \leq \delta \leq 1$

→ Random return:  $\mu_j = \bar{\mu}_j (1 - \delta w_j)$  ( $\bar{\mu}$  = nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \leq \theta \leq 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \geq 0$  such that

$$\text{Prob}(\nu - \mu^T x^* \geq \rho) \geq \theta$$

→ The adversary wants to maximize VAR

## Uncertainty set

Given a vector  $x^*$  of assets, the adversary

- 1 Chooses a vector  $w \in R^n$  ( $n = \text{no. of assets}$ ) with  $0 \leq w_j \leq 1$  for all  $j$ .
- 2 Chooses a random variable  $0 \leq \delta \leq 1$

→ Random return:  $\mu_j = \bar{\mu}_j (1 - \delta w_j)$  ( $\bar{\mu}$  = nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \leq \theta \leq 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \geq 0$  such that

$$\text{Prob}(\nu - \mu^T x^* \geq \rho) \geq \theta$$

→ The adversary wants to maximize VAR

## Uncertainty set

Given a vector  $x^*$  of assets, the adversary

- 1 Chooses a vector  $w \in R^n$  ( $n = \text{no. of assets}$ ) with  $0 \leq w_j \leq 1$  for all  $j$ .
- 2 Chooses a random variable  $0 \leq \delta \leq 1$

→ Random return:  $\mu_j = \bar{\mu}_j (1 - \delta w_j)$  ( $\bar{\mu}$  = nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \leq \theta \leq 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \geq 0$  such that

$$\text{Prob}(\nu - \mu^T x^* \geq \rho) \geq \theta$$

→ The adversary wants to maximize VAR

## Uncertainty set

Given a vector  $x^*$  of assets, the adversary

- 1 Chooses a vector  $w \in R^n$  ( $n = \text{no. of assets}$ ) with  $0 \leq w_j \leq 1$  for all  $j$ .
- 2 Chooses a random variable  $0 \leq \delta \leq 1$

→ Random return:  $\mu_j = \bar{\mu}_j (1 - \delta w_j)$  ( $\bar{\mu}$  = nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \leq \theta \leq 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \geq 0$  such that

$$\text{Prob}(\nu - \mu^T x^* \geq \rho) \geq \theta$$

→ The adversary wants to maximize VAR

## Uncertainty set

Given a vector  $x^*$  of assets, the adversary

- 1 Chooses a vector  $w \in R^n$  ( $n = \text{no. of assets}$ ) with  $0 \leq w_j \leq 1$  for all  $j$ .
- 2 Chooses a random variable  $0 \leq \delta \leq 1$

→ Random return:  $\mu_j = \bar{\mu}_j (1 - \delta w_j)$  ( $\bar{\mu}$  = nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \leq \theta \leq 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \geq 0$  such that

$$\text{Prob}(\nu - \mu^T x^* \geq \rho) \geq \theta$$

→ The adversary wants to maximize VAR

## OR ...

Given a vector  $x^*$  of assets, the adversary

- 1 Chooses a vector  $w \in R^n$  ( $n = \text{no. of assets}$ ) with  $0 \leq w_j \leq W$  for all  $j$ .
- 2 Chooses a random variable  $0 \leq \delta \leq 1$

→ Random return:  $\mu_j = \bar{\mu}_j - \delta w_j$  ( $\bar{\mu}$  = nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \leq \theta \leq 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \geq 0$  such that

$$\text{Prob}(\nu - \mu^T x^* \geq \rho) \geq \theta$$

→ The adversary wants to maximize VAR

OR ...

Given a vector  $x^*$  of assets, the adversary

- 1 Chooses a vector  $w \in R^n$  ( $n = \text{no. of assets}$ ) with  $0 \leq w_j \leq W$  for all  $j$ .
- 2 Chooses a random variable  $0 \leq \delta \leq 1$

→ Random return:  $\mu_j = \bar{\mu}_j - \delta w_j$  ( $\bar{\mu}$  = nominal returns).

**Definition:** Given reals  $\nu$  and  $0 \leq \theta \leq 1$  the *value-at-risk* of  $x^*$  is the real  $\rho \geq 0$  such that

$$\text{Prob}(\nu - \mu^T x^* \geq \rho) \geq \theta$$

→ The adversary wants to maximize VAR

## The classical factor model for returns

$$\mu = \bar{\mu} + V^T f + \epsilon$$

where

- $\bar{\mu}$  = expected return,
- $V$  = “factor exposure matrix”,
- $f$  = a bounded random variable,
- $\epsilon$  = residual errors

$V$  is  $r \times n$  with  $r \ll n$ .

→ Random return  $r_j = \bar{\mu}_j(1 - \delta w_j)$  where  $0 \leq w_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

A discrete distribution:

- We are given **fixed** values  $0 = \delta_0 \leq \delta_2 \leq \dots \leq \delta_K = 1$   
example:  $\delta_i = \frac{i}{K}$
- Adversary chooses  $\pi_i = \text{Prob}(\delta = \delta_i)$ ,  $0 \leq i \leq K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^l \leq \pi_i \leq \pi_i^u$   
(and possibly other constraints)
- Tier constraints: for sets (“tiers”)  $T_h$  of assets,  $1 \leq h \leq H$ , we require:  
 $E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h$  (given)  
or,  $(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$
- Cardinality constraint:  $w_j > 0$  for at most  $N$  indices  $j$

→ Random return  $r_j = \bar{\mu}_j(1 - \delta w_j)$  where  $0 \leq w_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

A *discrete distribution*:

- We are given **fixed** values  $0 = \delta_0 \leq \delta_2 \leq \dots \leq \delta_K = 1$

example:  $\delta_i = \frac{i}{K}$

- Adversary chooses  $\pi_i = \text{Prob}(\delta = \delta_i)$ ,  $0 \leq i \leq K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^l \leq \pi_i \leq \pi_i^u$  (and possibly other constraints)
- Tier constraints: for sets (“tiers”)  $T_h$  of assets,  $1 \leq h \leq H$ , we require:

$$E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h \quad (\text{given})$$

$$\text{or, } (\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$$

- Cardinality constraint:  $w_j > 0$  for at most  $N$  indices  $j$

→ Random return  $r_j = \bar{\mu}_j(1 - \delta w_j)$  where  $0 \leq w_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

A *discrete distribution*:

- We are given **fixed** values  $0 = \delta_0 \leq \delta_2 \leq \dots \leq \delta_K = 1$   
example:  $\delta_i = \frac{i}{K}$
- Adversary chooses  $\pi_i = \text{Prob}(\delta = \delta_i)$ ,  $0 \leq i \leq K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^l \leq \pi_i \leq \pi_i^u$   
(and possibly other constraints)
- Tier constraints: for sets (“tiers”)  $T_h$  of assets,  $1 \leq h \leq H$ , we require:  
 $E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h$  (given)  
or,  $(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$
- Cardinality constraint:  $w_j > 0$  for at most  $N$  indices  $j$

→ Random return  $r_j = \bar{\mu}_j(1 - \delta w_j)$  where  $0 \leq w_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

A discrete distribution:

- We are given **fixed** values  $0 = \delta_0 \leq \delta_2 \leq \dots \leq \delta_K = 1$   
example:  $\delta_i = \frac{i}{K}$
- Adversary chooses  $\pi_i = \text{Prob}(\delta = \delta_i)$ ,  $0 \leq i \leq K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^l \leq \pi_i \leq \pi_i^u$   
(and possibly other constraints)
- Tier constraints: for sets (“tiers”)  $T_h$  of assets,  $1 \leq h \leq H$ , we require:  
 $E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h$  (given)  
or,  $(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$
- Cardinality constraint:  $w_j > 0$  for at most  $N$  indices  $j$

→ Random return  $r_j = \bar{\mu}_j(1 - \delta w_j)$  where  $0 \leq w_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

A discrete distribution:

- We are given **fixed** values  $0 = \delta_0 \leq \delta_2 \leq \dots \leq \delta_K = 1$   
example:  $\delta_i = \frac{i}{K}$
- Adversary chooses  $\pi_i = \text{Prob}(\delta = \delta_i)$ ,  $0 \leq i \leq K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^l \leq \pi_i \leq \pi_i^u$   
(and possibly other constraints)
- Tier constraints: for sets (“tiers”)  $T_h$  of assets,  $1 \leq h \leq H$ , we require:  
 $E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h$  (given)  
or,  $(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$
- Cardinality constraint:  $w_j > 0$  for at most  $N$  indices  $j$

→ Random return  $r_j = \bar{\mu}_j(1 - \delta w_j)$  where  $0 \leq w_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

A discrete distribution:

- We are given **fixed** values  $0 = \delta_0 \leq \delta_2 \leq \dots \leq \delta_K = 1$   
example:  $\delta_i = \frac{i}{K}$
- Adversary chooses  $\pi_i = \text{Prob}(\delta = \delta_i)$ ,  $0 \leq i \leq K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^l \leq \pi_i \leq \pi_i^u$   
(and possibly other constraints)
- **Tier** constraints: for sets (“tiers”)  $T_h$  of assets,  $1 \leq h \leq H$ , we require:

$$E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h \quad (\text{given})$$

$$\text{or, } (\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$$

- Cardinality constraint:  $w_j > 0$  for at most  $N$  indices  $j$

→ Random return  $r_j = \bar{\mu}_j(1 - \delta w_j)$  where  $0 \leq w_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

A *discrete distribution*:

- We are given **fixed** values  $0 = \delta_0 \leq \delta_2 \leq \dots \leq \delta_K = 1$   
example:  $\delta_i = \frac{i}{K}$
- Adversary chooses  $\pi_i = \text{Prob}(\delta = \delta_i)$ ,  $0 \leq i \leq K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^l \leq \pi_i \leq \pi_i^u$   
(and possibly other constraints)
- **Tier** constraints: for sets (“tiers”)  $T_h$  of assets,  $1 \leq h \leq H$ , we require:

$$E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h \quad (\text{given})$$

$$\text{or, } (\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$$

- Cardinality constraint:  $w_j > 0$  for at most  $N$  indices  $j$

→ Random return  $r_j = \bar{\mu}_j(1 - \delta w_j)$  where  $0 \leq w_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

A *discrete distribution*:

- We are given **fixed** values  $0 = \delta_0 \leq \delta_2 \leq \dots \leq \delta_K = 1$   
example:  $\delta_i = \frac{i}{K}$
- Adversary chooses  $\pi_i = \text{Prob}(\delta = \delta_i)$ ,  $0 \leq i \leq K$
- The  $\pi_i$  are *constrained*: we have fixed bounds,  $\pi_i^l \leq \pi_i \leq \pi_i^u$   
(and possibly other constraints)
- **Tier** constraints: for sets (“tiers”)  $T_h$  of assets,  $1 \leq h \leq H$ , we require:

$$E(\delta \sum_{j \in T_h} w_j) \leq \Gamma_h \quad (\text{given})$$

$$\text{or, } (\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h$$

- Cardinality constraint:  $w_j > 0$  for at most  $N$  indices  $j$

## The adversarial problem is “easy”

$K$  = no. of points in discrete distribution,  $H$  = no. of tiers

### Theorem

- Without the cardinality constraint, for each fixed  $K$  and  $H$  the adversarial problem can be solved as a polynomial number of linear programs.
- With the cardinality constraint, for each fixed  $K$  and  $H$  the adversarial problem can be solved as a polynomial number of knapsack problems.

## The adversarial problem is “easy”

$K$  = no. of points in discrete distribution,  $H$  = no. of tiers

### Theorem

- Without the cardinality constraint, for each fixed  $K$  and  $H$  the adversarial problem can be solved as a polynomial number of linear programs.
- With the cardinality constraint, for each fixed  $K$  and  $H$  the adversarial problem can be solved as a polynomial number of knapsack problems.

## The adversarial problem is “easy”

$K$  = no. of points in discrete distribution,  $H$  = no. of tiers

### Theorem

- Without the cardinality constraint, for each fixed  $K$  and  $H$  the adversarial problem can be solved as a polynomial number of linear programs.
- With the cardinality constraint, for each fixed  $K$  and  $H$  the adversarial problem can be solved as a polynomial number of knapsack problems.

## Adversarial problem as an MIP

Recall: random return  $\mu_j = \bar{\mu}_j(1 - \delta w_j)$

where  $\delta = \delta_i$  (given) with probability  $\pi_i$  (chosen by adversary),

$0 \leq \delta_0 \leq \delta_1 \leq \dots \leq \delta_K = 1$  and  $0 \leq w$

$$\min_{\pi, w, v} \min_{1 \leq i \leq k} V_i$$

Subject to

$$0 \leq w_j \leq 1, \text{ all } j, \pi_i^l \leq \pi_i \leq \pi_i^u, \text{ all } i,$$

$$\sum_i \pi_i = 1,$$

$$V_i = \sum_j \bar{\mu}_j(1 - \delta_i w_j) x_j^*, \text{ if } \pi_i + \pi_{i+1} + \dots + \pi_K \geq 1 - \theta$$

$$V_i = M \text{ (large)}, \text{ otherwise}$$

$$(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h, \text{ for each tier } h$$

## Adversarial problem as an MIP

Recall: random return  $\mu_j = \bar{\mu}_j(1 - \delta w_j)$

where  $\delta = \delta_i$  (given) with probability  $\pi_i$  (chosen by adversary),

$0 \leq \delta_0 \leq \delta_1 \leq \dots \leq \delta_K = 1$  and  $0 \leq w$

$$\min_{\pi, w, v} \min_{1 \leq i \leq k} V_i$$

Subject to

$$0 \leq w_j \leq 1, \text{ all } j, \pi_i^l \leq \pi_i \leq \pi_i^u, \text{ all } i,$$

$$\sum_i \pi_i = 1,$$

$$V_i = \sum_j \bar{\mu}_j(1 - \delta_i w_j) x_j^*, \text{ if } \pi_i + \pi_{i+1} + \dots + \pi_K \geq 1 - \theta$$

$$V_i = M \text{ (large)}, \text{ otherwise}$$

$$(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h, \text{ for each tier } h$$

## Adversarial problem as an MIP

Recall: random return  $\mu_j = \bar{\mu}_j(1 - \delta w_j)$

where  $\delta = \delta_i$  (given) with probability  $\pi_i$  (chosen by adversary),

$0 \leq \delta_0 \leq \delta_1 \leq \dots \leq \delta_K = 1$  and  $0 \leq w$

$$\min_{\pi, w, v} \min_{1 \leq i \leq k} V_i$$

Subject to

$$0 \leq w_j \leq 1, \text{ all } j, \pi_i^l \leq \pi_i \leq \pi_i^u, \text{ all } i,$$

$$\sum_i \pi_i = 1,$$

$$V_i = \sum_j \bar{\mu}_j(1 - \delta_i w_j) x_j^*, \text{ if } \pi_i + \pi_{i+1} + \dots + \pi_K \geq 1 - \theta$$

$$V_i = M \text{ (large)}, \text{ otherwise}$$

$$(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h, \text{ for each tier } h$$

## Adversarial problem as an MIP

Recall: random return  $\mu_j = \bar{\mu}_j(1 - \delta w_j)$

where  $\delta = \delta_i$  (given) with probability  $\pi_i$  (chosen by adversary),

$0 \leq \delta_0 \leq \delta_1 \leq \dots \leq \delta_K = 1$  and  $0 \leq w$

$$\min_{\pi, w, v} \min_{1 \leq i \leq k} V_i$$

Subject to

$$0 \leq w_j \leq 1, \text{ all } j, \pi_i^l \leq \pi_i \leq \pi_i^u, \text{ all } i,$$

$$\sum_i \pi_i = 1,$$

$$V_i = \sum_j \bar{\mu}_j(1 - \delta_i w_j) x_j^*, \text{ if } \pi_i + \pi_{i+1} + \dots + \pi_K \geq 1 - \theta$$

$$V_i = M \text{ (large)}, \text{ otherwise}$$

$$(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h, \text{ for each tier } h$$

## Approximation

$$\left(\sum_i \delta_i \pi_i\right) \sum_{j \in T_h} w_j \leq \Gamma_h, \quad \text{for each tier } h \quad (*)$$

Let  $N > 0$  be an integer. For  $1 \leq k \leq N$ , write

$$\frac{k}{N} \sum_{j \in T_h} w_j \leq \Gamma_h + M(1 - z_{hk}), \quad \text{where}$$

$$z_{hk} = 1 \text{ if } \frac{k-1}{N} < \sum_i \delta_i \pi_i \leq \frac{k}{N}$$

$$z_{hk} = 0 \text{ otherwise}$$

$$\sum_k z_{hk} = 1$$

and  $M$  is large

**Lemma.** Under reasonable conditions, replacing  $(*)$  with this system changes the value of the problem by at most a factor of  $(1 + \frac{1}{N})$

## Approximation

$$(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h, \quad \text{for each tier } h \quad (*)$$

Let  $N > 0$  be an integer. For  $1 \leq k \leq N$ , write

$$\frac{k}{N} \sum_{j \in T_h} w_j \leq \Gamma_h + M(1 - z_{hk}), \quad \text{where}$$

$$z_{hk} = 1 \text{ if } \frac{k-1}{N} < \sum_i \delta_i \pi_i \leq \frac{k}{N}$$

$$z_{hk} = 0 \text{ otherwise}$$

$$\sum_k z_{hk} = 1$$

and  $M$  is large

**Lemma.** Under reasonable conditions, replacing  $(*)$  with this system changes the value of the problem by at most a factor of  $(1 + \frac{1}{N})$

## Approximation

$$(\sum_i \delta_i \pi_i) \sum_{j \in T_h} w_j \leq \Gamma_h, \quad \text{for each tier } h \quad (*)$$

Let  $N > 0$  be an integer. For  $1 \leq k \leq N$ , write

$$\frac{k}{N} \sum_{j \in T_h} w_j \leq \Gamma_h + M(1 - z_{hk}), \quad \text{where}$$

$$z_{hk} = 1 \text{ if } \frac{k-1}{N} < \sum_i \delta_i \pi_i \leq \frac{k}{N}$$

$$z_{hk} = 0 \text{ otherwise}$$

$$\sum_k z_{hk} = 1$$

and  $M$  is large

**Lemma.** Under reasonable conditions, replacing  $(*)$  with this system changes the value of the problem by at most a factor of  $(1 + \frac{1}{N})$

## Implementor's problem

Find a near-optimal solution with minimum value-at-risk

Nominal problem:

$$v^* = \min_x \lambda x^T Q x - \mu^T x$$

Subject to:

$$Ax \geq b$$

## Implementor's problem

Find a near-optimal solution with minimum value-at-risk

**Nominal problem:**

$$v^* = \min_x \lambda x^T Q x - \mu^T x$$

Subject to:

$$Ax \geq b$$

## Implementor's problem

Find a near-optimal solution with minimum value-at-risk

Given asset weights  $\mathbf{x}$ , we have:

- value-at-risk  $\geq \rho$ , if the adversary can produce a return vector  $\mu$  with

$$\text{Prob}(\nu - \mu^T \mathbf{x} \geq \rho) \geq \theta$$

where  $\nu$  is a fixed reference value.

## Implementor's problem

Find a near-optimal solution with minimum value-at-risk

Implementor's problem at iteration  $r$ :

$$\min V$$

Subject to:

$$\lambda x^T Q x - \mu^T x \leq (1 + \epsilon) v^*$$

$$A x \geq b$$

$$V \geq \nu - \sum_j \bar{\mu}_j \left(1 - \delta_{i(t)} w_j^{(t)}\right) x_j, \quad t = 1, 2, \dots, r-1$$

Here,  $\delta_{i(t)}$  and  $w^{(t)}$  are the adversary's output at iteration  $t < r$ .

## Implementor's problem

Find a near-optimal solution with minimum value-at-risk

### Implementor's problem at iteration $r$ :

$$\min V$$

Subject to:

$$\lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \boldsymbol{\mu}^T \mathbf{x} \leq (1 + \epsilon) \mathbf{v}^*$$

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

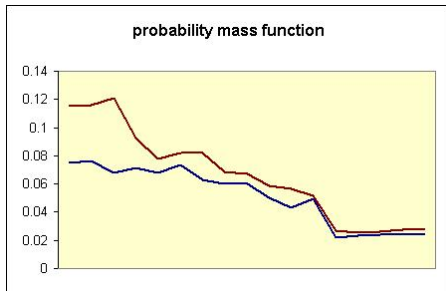
$$V \geq \nu - \sum_j \bar{\mu}_j \left( 1 - \delta_{i(t)} \mathbf{w}_j^{(t)} \right) \mathbf{x}_j, \quad t = 1, 2, \dots, r - 1$$

Here,  $\delta_{i(t)}$  and  $\mathbf{w}^{(t)}$  are the adversary's output at iteration  $t < r$ .

## First set of experiments

1338 assets, 41 factors, 81 rows

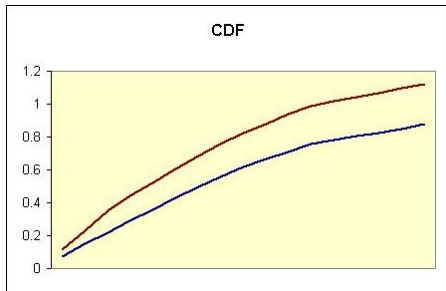
- problem: find a “near optimal” solution with minimum value-at-risk, for a given threshold probability  $\theta$
- experiment: investigate different values of  $\theta$
- “near optimal”: want solutions that are at most 1% more expensive than optimal
- random variable  $\delta$ :



## First set of experiments

1338 assets, 41 factors, 81 rows

- problem: find a “near optimal” solution with minimum value-at-risk, for a given threshold probability  $\theta$
- experiment: investigate different values of  $\theta$
- “near optimal”: want solutions that are at most 1% more expensive than optimal
- random variable  $\delta$ :

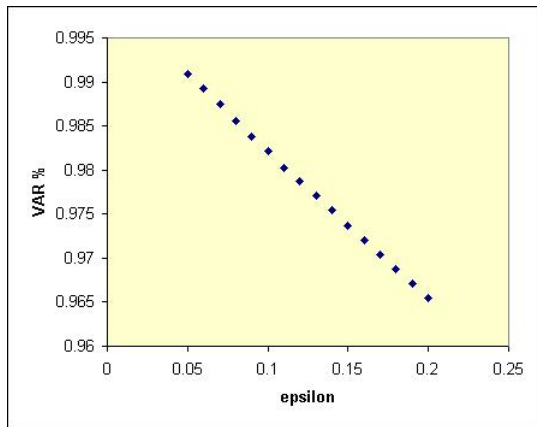


1338 assets, 41 factors, 81 rows,  $\leq 1\%$  suboptimality

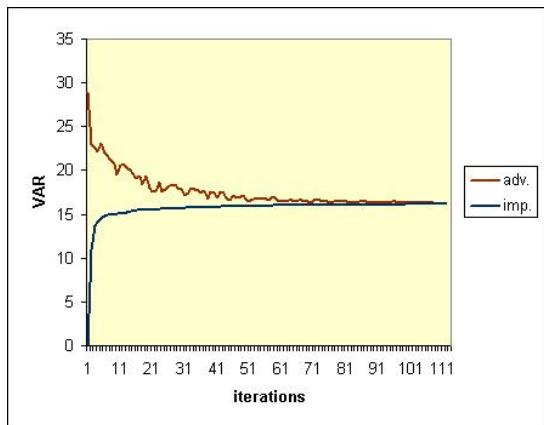
$\theta$	time (sec.)	iters.	VAR	VAR as %
.89	1.18	2	3.43131	71.50
.90	1.42	3	3.74498	78.04
.91	1.42	3	3.74498	78.04
.92	3.47	11	4.05669	84.53
.93	6.35	29	4.05721	84.54
.94	6.37	20	4.05721	84.54
.95	26.59	51	4.35481	90.74
.96	26.25	51	4.35481	90.74
.97	26.20	51	4.35481	90.74
.98	33.07	58	4.63938	96.67
.99	33.11	58	4.63938	96.67

## Second set of experiments

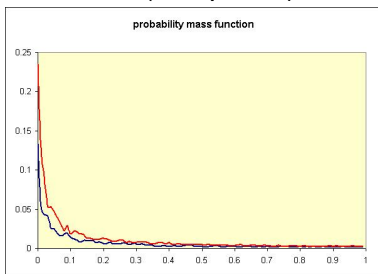
Fix  $\theta = 0.90$  but vary suboptimality criterion



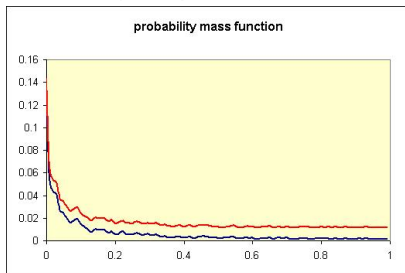
## Typical convergence behavior



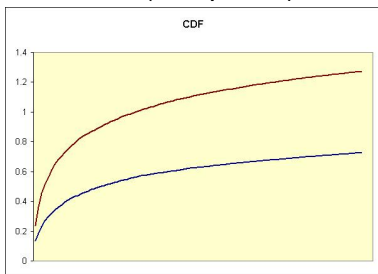
- Heavy tail, proportional error (100 points):



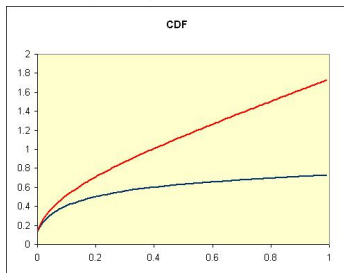
- Heavy tail, constant error (100 points):



- Heavy tail, proportional error (100 points):



- Heavy tail, constant error (100 points):



$\theta$	Proportional		Constant	
	Time	Its	Time	Its
.85	4.79	8	11.84	11
.86	2.32	3	8.27	8
.87	6.40	10	9.55	11
.88	4.34	4	18.10	24
.89	8.00	14	5.85	6
.90	2.58	4	13.54	20
.91	4.79	9	16.31	23
.92	7.99	15	13.13	22
.93	13.43	27	22.47	40
.94	10.04	15	21.99	40
.95	9.59	16	11.90	25
.96	6.63	17	29.89	54
.97	48.43	110	16.45	35
.98	20.25	53	20.25	45
.99	22.02	52	21.89	47

## A difficult case

- 2464 columns, 152 factors, 3 tiers
- time = 6191 seconds
- 258 iterations
- implementor time = 6123 seconds,  
adversarial time = 20 seconds

## A difficult case

- 2464 columns, 152 factors, 3 tiers
- time = 6191 seconds
- 258 iterations
- implementor time = 6123 seconds,  
adversarial time = 20 seconds

## A difficult case

- 2464 columns, 152 factors, 3 tiers
- time = 6191 seconds
- 258 iterations
- implementor time = 6123 seconds,  
adversarial time = 20 seconds

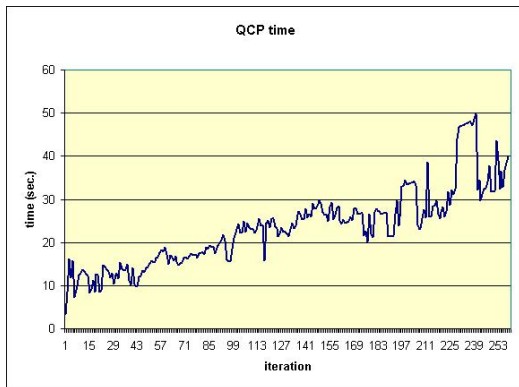
## A difficult case

- 2464 columns, 152 factors, 3 tiers
- time = 6191 seconds
- 258 iterations
- implementor time = 6123 seconds,  
adversarial time = 20 seconds

## A difficult case

- 2464 columns, 152 factors, 3 tiers
- time = 6191 seconds
- 258 iterations
- implementor time = 6123 seconds,  
adversarial time = 20 seconds

## Implementor runtime



## Implementor's problem at iteration $r$

**min**  $V$

Subject to:

$$\lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \mu^T \mathbf{x} \leq (1 + \epsilon) \mathbf{v}^*$$

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

$$V \geq \nu - \sum_j \bar{\mu}_j \left( 1 - \delta_{i(t)} \mathbf{w}_j^{(t)} \right) \mathbf{x}_j, \quad t = 1, 2, \dots, r - 1$$

Here,  $\delta_{i(t)}$  and  $\mathbf{w}^{(t)}$  are the adversary's output at iteration  $t < r$ .

## Implementor's problem at iteration $r$

Approximate version

**min**  $V$

Subject to:

$$2\lambda \mathbf{x}_{(k)}^T \mathbf{Q} \mathbf{x} - \lambda \mathbf{x}_{(k)}^T \mathbf{Q} \mathbf{x}_{(k)} - \boldsymbol{\mu}^T \mathbf{x} \leq (1 + \epsilon) \mathbf{v}^*, \quad \forall k < r$$

$$\mathbf{A} \mathbf{x} \geq \mathbf{b}$$

$$V \geq \nu - \sum_j \bar{\mu}_j \left( 1 - \delta_{i(k)} \mathbf{w}_j^{(k)} \right) \mathbf{x}_j, \quad \forall k < r$$

Here,  $\delta_{i(k)}$  and  $\mathbf{w}^{(k)}$  are the adversary's output at iteration  $k < r$ , and  $\mathbf{x}_{(k)}$  is the implementor's output at iteration  $k$ .

## Does it work?

- Before: 258 iterations, **6191** seconds
- Linearized: 1776 iterations, **3969** seconds

## Does it work?

- Before: 258 iterations, **6191** seconds
- Linearized: 1776 iterations, **3969** seconds

## Averaging

- $\mathbf{x}_{(k)}$  is the implementor's output at iteration  $k$ .
- Define  $\mathbf{y}_{(1)} = \mathbf{x}_{(1)}$
- For  $k > 1$ ,  $\mathbf{y}_{(k)} = \lambda \mathbf{x}_{(k)} + (1 - \lambda) \mathbf{y}_{(k-1)}$ ,  
 $0 \leq \lambda \leq 1$
- Input  $\mathbf{y}_{(k)}$  to the adversary
- Old ideas, also Nesterov, Nemirovsky (2003)

## Averaging

- $\mathbf{x}_{(k)}$  is the implementor's output at iteration  $k$ .
- Define  $\mathbf{y}_{(1)} = \mathbf{x}_{(1)}$
- For  $k > 1$ ,  $\mathbf{y}_{(k)} = \lambda \mathbf{x}_{(k)} + (1 - \lambda) \mathbf{y}_{(k-1)}$ ,  
 $0 \leq \lambda \leq 1$
- Input  $\mathbf{y}_{(k)}$  to the adversary
- Old ideas, also Nesterov, Nemirovsky (2003)

## Averaging

- $\mathbf{x}_{(k)}$  is the implementor's output at iteration  $k$ .
- Define  $\mathbf{y}_{(1)} = \mathbf{x}_{(1)}$
- For  $k > 1$ ,  $\mathbf{y}_{(k)} = \lambda \mathbf{x}_{(k)} + (1 - \lambda) \mathbf{y}_{(k-1)}$ ,  
 $0 \leq \lambda \leq 1$
- Input  $\mathbf{y}_{(k)}$  to the adversary
- Old ideas, also Nesterov, Nemirovsky (2003)

## Averaging

- $\mathbf{x}_{(k)}$  is the implementor's output at iteration  $k$ .
- Define  $\mathbf{y}_{(1)} = \mathbf{x}_{(1)}$
- For  $k > 1$ ,  $\mathbf{y}_{(k)} = \lambda \mathbf{x}_{(k)} + (1 - \lambda) \mathbf{y}_{(k-1)}$ ,  
 $0 \leq \lambda \leq 1$
- Input  $\mathbf{y}_{(k)}$  to the adversary
- Old ideas, also Nesterov, Nemirovsky (2003)

## Averaging

- $\mathbf{x}_{(k)}$  is the implementor's output at iteration  $k$ .
- Define  $\mathbf{y}_{(1)} = \mathbf{x}_{(1)}$
- For  $k > 1$ ,  $\mathbf{y}_{(k)} = \lambda \mathbf{x}_{(k)} + (1 - \lambda) \mathbf{y}_{(k-1)}$ ,  
 $0 \leq \lambda \leq 1$
- Input  $\mathbf{y}_{(k)}$  to the adversary
- Old ideas, also Nesterov, Nemirovsky (2003)

## Does it work?

- Default: 258 iterations, **6191** seconds
- Linearized: 1776 iterations, **3969** seconds
- Averaging plus Linearized: 860 iterations, **530** seconds

## Does it work?

- Default: 258 iterations, **6191** seconds
- Linearized: 1776 iterations, **3969** seconds
- Averaging plus Linearized: 860 iterations, **530** seconds

## Does it work?

- Default: 258 iterations, **6191** seconds
- Linearized: 1776 iterations, **3969** seconds
- Averaging plus Linearized: 860 iterations, **530** seconds

## Other robust models

- Min-max expected loss with orthogonal missing factor

Random return =  $\bar{\mu} \bullet (1 - \delta \mathbf{w})$  where  $-1 \leq \mathbf{w}_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

normalization constraints, e.g.  $\sum_j \mathbf{w}_j = 0$

- errors in covariance matrix  $\mathbf{Q}$

robust problem:  $\rightarrow \min_x \max_{\mathbf{Q} \in \mathcal{Q}} \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \mu^T \mathbf{x}$

## Other robust models

- Min-max expected loss with orthogonal missing factor

Random return =  $\bar{\mu} \bullet (1 - \delta \mathbf{w})$  where  $-1 \leq \mathbf{w}_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

normalization constraints, e.g.  $\sum_j \mathbf{w}_j = \mathbf{0}$

- errors in covariance matrix  $\mathbf{Q}$

robust problem:  $\rightarrow \min_x \max_{\mathbf{Q} \in \mathcal{Q}} \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \mu^T \mathbf{x}$

## Other robust models

- Min-max expected *loss* with orthogonal missing factor

Random return =  $\bar{\mu} \bullet (1 - \delta \mathbf{w})$  where  $-1 \leq \mathbf{w}_j \leq 1 \quad \forall j$ , and  $0 \leq \delta \leq 1$  is a random variable.

normalization constraints, e.g.  $\sum_j \mathbf{w}_j = \mathbf{0}$

- errors in covariance matrix  $\mathbf{Q}$

robust problem:  $\rightarrow \min_{\mathbf{x}} \max_{\mathbf{Q} \in \mathcal{Q}} \lambda \mathbf{x}^T \mathbf{Q} \mathbf{x} - \mu^T \mathbf{x}$

On-going work: a provably good version of Benders' algorithm