

$(S, r) \leftarrow (S, S)$ Policies

- Algorithms for (S, r) policies
- Exact implementors
- Stochastic Dynamic Programming
- Next topic:

Compound Poisson

$X =$ random demand size

Special case $P(X=1) = 1$
reduces to Poisson case

$$e(G, r) = \frac{KX \left(E_{\text{mean}}(G, X) \right) + \sum_{j=1}^{r+Q} G(j)}{Q}$$

$\underbrace{\hspace{10em}}_{= 1 \text{ in Poisson case}}$

$r + j$; $j \sim \cup \{1, \dots, Q\}$

$r + j - X \leq r$ iff $X \geq j$

$\frac{1}{Q} \sum_{j=1}^Q P(X \geq j) =$ probability of placing an order

$\frac{1}{Q} E_{\text{num}}(Q, X)$

• General result

by X discrete, non-negative r.v.
+ takes integer values

$$EX = 1 \cdot p_1 + 2p_2 + 3p_3 + \dots$$

$$= p_1 + p_2 + p_3 + \dots$$

$$+ p_2 + p_3 + \dots$$

$$p_3 + \dots$$

$$p_4 + \dots$$

$$EX = P(X \geq 1) + P(X \geq 2) + \dots$$

$$= \sum_{j=1}^{\infty} P(X \geq j)$$

$$X^{\infty} = \min(Q, X)$$

$$E X^{\infty} = P(X^{\infty} \geq 1) + P(X^{\infty} \geq 2) + \dots$$

$$P(X^{\infty} \geq j) = \begin{cases} P(X \geq j) & j \leq Q \\ 0 & j > Q \end{cases}$$

$$\begin{aligned} \Rightarrow E X^{\infty} &= E \min(Q, X) \\ &= \sum_{j=1}^Q P(X \geq j) \end{aligned}$$

Ordering cost (average)

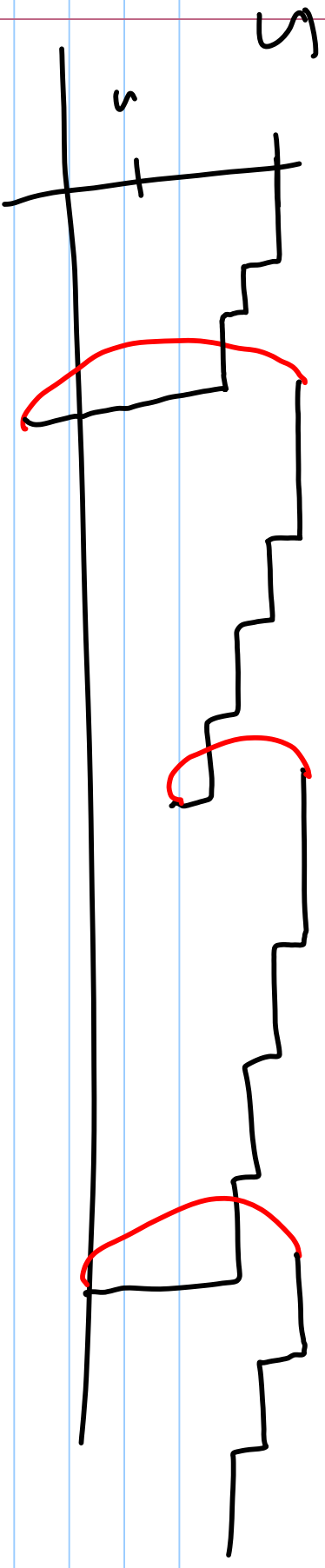
$$KX \frac{1}{Q} \sum_{j=1}^Q P(X \geq j)$$

$$= KX \frac{E_{\text{num}}(Q, X)}{Q}$$

□

(s, S) policies

- Place an order to bring IP to S whenever it falls to or below s.



IP has a distribution

concentrated over $\{s+1, s+2, \dots, S\}$

but it is not necessarily uniform.

$$C(s, S) = K \cdot \text{prob of placing an order} \\ + \sum_{j=1}^{S-s} G(s+j) F_j$$

$$F_j = P_r (\text{IP} = s+j) \quad j=1, \dots, S-s$$

depends on s, S .

Q: How do we find the distribution of the IP?

Renewal Theory

$$y \in \{s+1, \dots, S\}$$

$K(s, y) =$ total expected cost until
the next order is placed

starting from $IP = s$

$M(i) =$ expected time until

an order is placed starting

from $IP = s+i$,
 $j=1, \dots, S-s$.

$$c(s, S) = \frac{K(s, S)}{M(s, S)}$$

$$K(s, y) = G(y) + K P_r(\mathcal{D} \geq y-s) \\ + \sum_{j=0}^{y-s-1} K(s, y-j) P_r(\mathcal{D} = j)$$

$$K(s, y) = G(y) + K \sum_{j \geq y-s} P_j \\ + \sum_{j=0}^{y-s-1} K(s, y-j) P_j .$$

$$M(j) = 0 \quad \forall \quad j = 0$$

$j > 1$

$$M(j) = 1 + O_P(D \geq j) + \sum_{x=1}^{j-1} M(j-x) P(D=x)$$

Solution is

$$M(j) = M(j-1) + m(j-1)$$

$$m(0) = \frac{1}{1-\rho_0}$$

$$m(j) = \sum_{k=0}^j p_k m(j-k) \quad j = 1, 2, \dots$$

$$k(s, s) = k + \sum_{j=0}^{s-1} m(j) G(s-j)$$

$$c(s, s) = \frac{k + \sum_{j=0}^{s-1} m(j) G(s-j)}{M(s-s)}$$

$$= \frac{k + \sum_{j=1}^{s-1} G(s+j) m(s-s-j)}{M(s-s)}$$

$$\frac{1}{M(S-s)} = \text{prob of placing an order}$$

$$\frac{m(S-s-j)}{M(S-s)} = P_r(\text{IP} = s+j) = q_j$$

o Algorithm

Key idea:

$e(s-1, S)$ in terms of $e(s, S)$ and $G(s)$

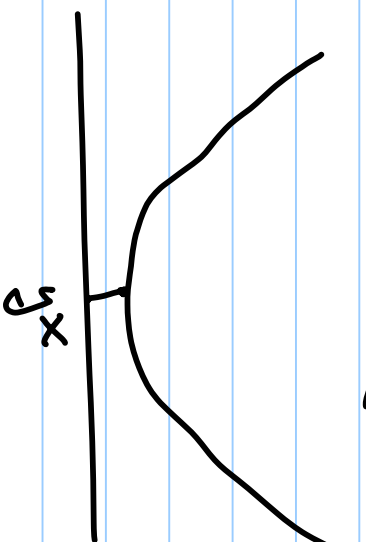
$$e(s-1, S) = d_n e(s, S) + (1-d_n) G(s)$$

$$d_n = \frac{M(n)}{M(n+1)} \quad ; n = S-1.$$

At position

Let y^x be a number of

$$G_1(y).$$



Step 0

$$S^0 = y^x, \quad S = y^{x-1}$$

DO WHILE $e(S, S^0) > G(y)$

$$S = S^{-1};$$

ENDDO;

$c^0 = c(s, S^0)$, $S = S^0 + 1$

Step 1

DO UNTIL $G(S) > c^0$;

IF $c(s, S) < c^0$

$S^0 = S$

DO WHILE $c(s, S^0) \leq G(s+1)$

$s = s + 1$;

ENDDO;

$c^0 = c(s, S^0)$

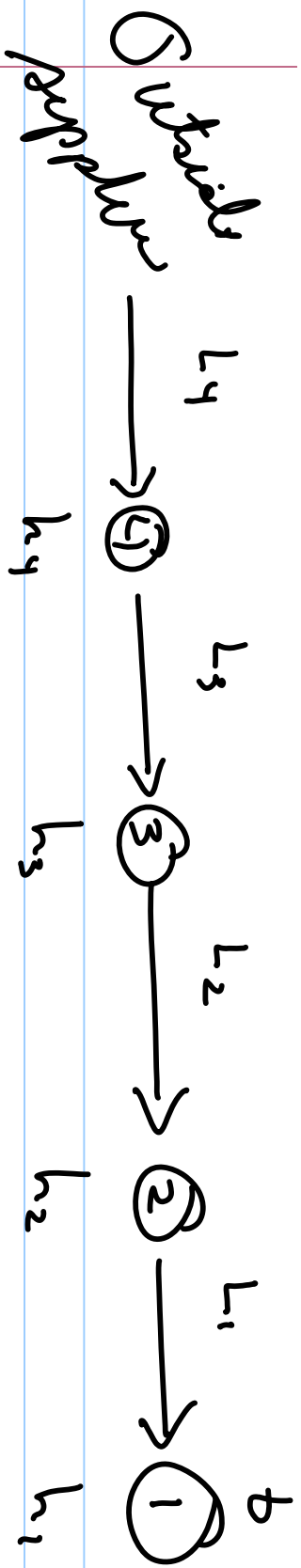
ENDIF
S = S + 1
END0;
END

Questions?

Single Location Models

- Newsvendor
 - Finite horizon
 - Infinite horizon
- Address Fixed costs
 - (Q, r) models
 - (s, S) "

Multi-echelon Models

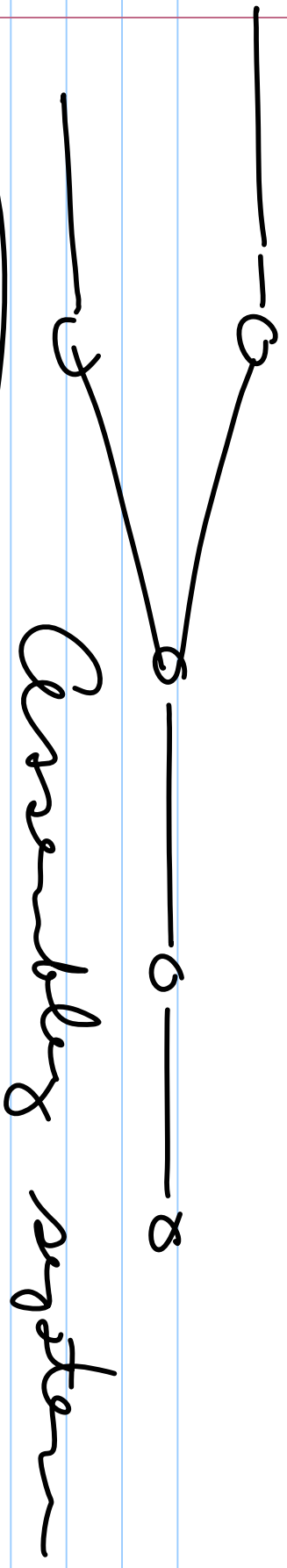


S E R I A L S Y S T E M

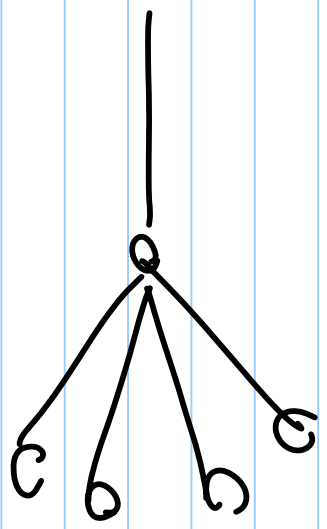
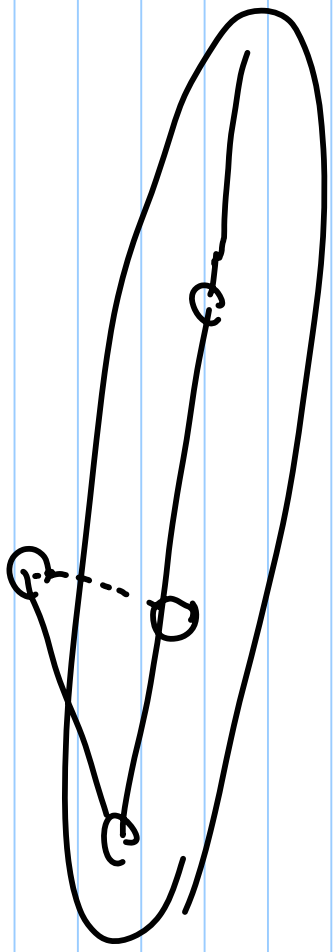
Base stock policies are optimal when setup costs are zero.

S_4 S_3 S_2 S_1

Server Allocation



Assembly system



Distribution system

• Approx / Heuristics

