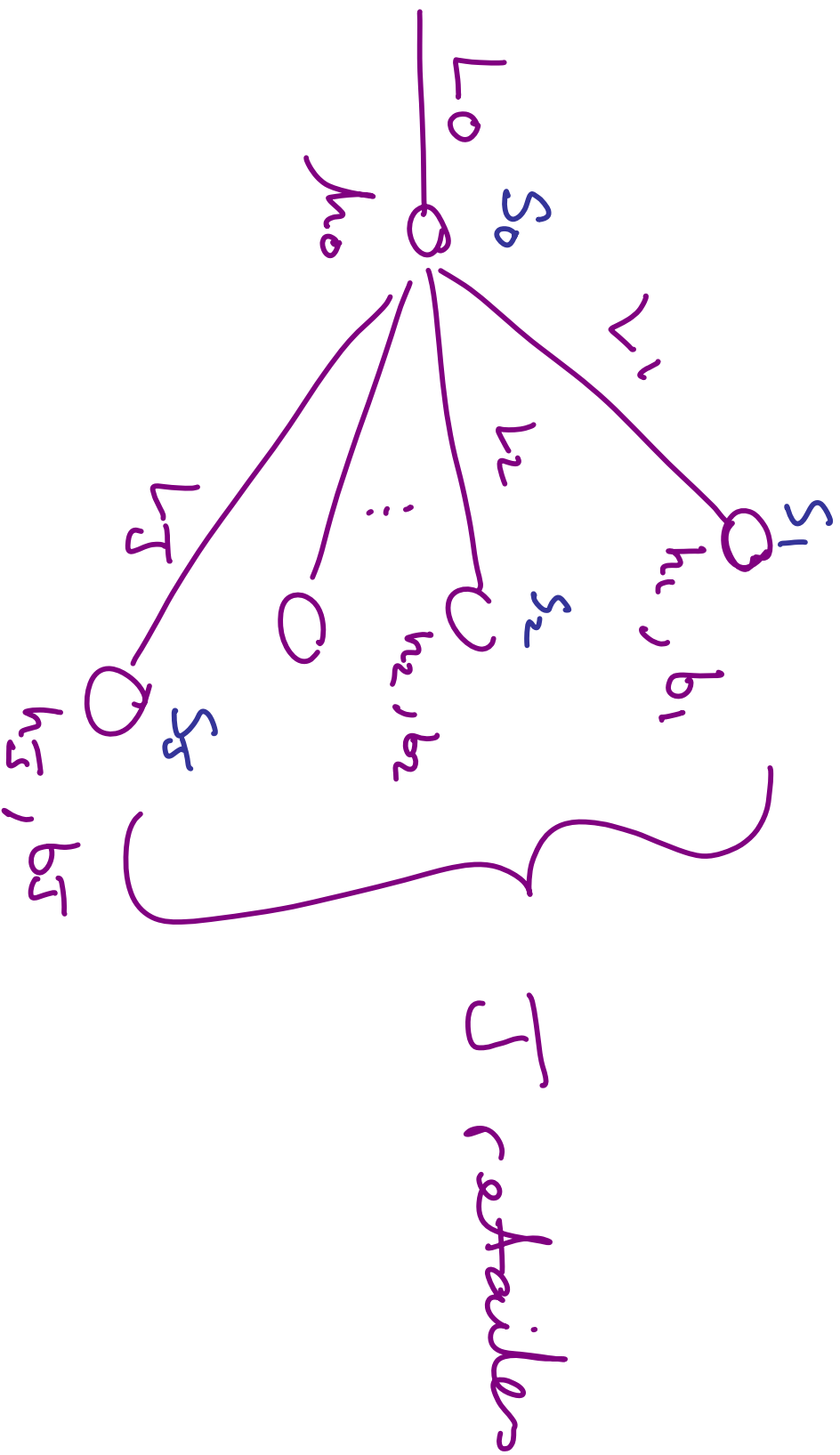


Distribution Algorithms



Local control

- Each store follows base stock policy
- Warehouse use FIFO

Central control

- Warehouse may use smarter allocation rules.

Parsons Demands

λ_j = demand rate at
retailer $j \in \{1, \dots, J\}$

$\lambda_0 = \sum_{j=1}^J \lambda_j$ overall demand
rate

S_j = base stock level
at retailer $j \in \{1, \dots, J\}$

S_0 = base stock at warehouse

$D_j \sim \text{Poisson}(\lambda_j L_j) \quad j \in \{1, \dots, J\}$
 $D_0 \sim \text{Poisson}(\lambda_0 L_0)$ warehouse

Backlogs & inventory at warehouse

$$B_0 = (D_0 - S_0)^+$$

$$I_0 = (S_0 - D_0)^+$$

$B_{0j} =$ backlogs at warehouse
due to retailer j

$$B_0 = \sum_{j=1}^J B_{0j}$$

B_j = backorder at retailer j

$$B_j = (B_{0j} + D_j - S_j)^+$$

$$I_j = (S_j - D_j - B_{0j})^+$$

$$C(S_0, S_1, \dots, S_n) = h_0 E[I_0]$$

$$+ \sum_{j=1}^n c_j c_j(S_0, S_j)$$

$$c_j c_j(S_0, S_j) = h_j E[I_j] + b_j E[B_j].$$

Given s_0 the problem of selecting $s_j(s_0)$ as a Newsreader problem $\mathcal{N}(h_j, b_j, \underbrace{\beta_{0j} + \mathcal{D}_j})$

↳ require calculations

$$s_j(s_0) = \arg \min_{\{s_j \in \mathcal{H} : P_r(\mathcal{D}_j + \beta_{0j} \leq s_j)\}}$$

$$\succ \left. \frac{b_j}{b_j + h_j} \right\}$$

$$\beta_{0j} \sim \text{bin}(\theta_j, \beta_0) \quad \theta_j \equiv \lambda_j / \lambda_0$$

$$\max_{s_0} [f(s_0)] + \underbrace{\sum_{j=1}^n c_j (s_0, s_j(s_0))}_{\text{function of } s_0 \text{ alone}}$$

- Not convex in s_0
- Search over $LB \leq s_0 \leq UB$ and find optimal solution.

Derivatives?

CD Derivatives

$$S_0 = 0$$

$$\begin{aligned} B_0 &= (D_0 - S_0)^+ \\ &= D_0 = \sum_{j=1}^n D_{0j} \end{aligned}$$

$D_{0j} \sim \text{Poisson}(\lambda_j t_0)$

$$B_{0j} = D_{0j}$$

$$B_{0j} + D_j \sim \text{Poisson}(\gamma_j L_0 + \gamma_j L_j) \\ \sim \text{Poisson}(\gamma_j (L_0 + L_j))$$

Simply solve \mathcal{I} numerically

problems

$$\mathcal{N}(L_j, b_j \text{Poisson}(\gamma_j (L_0 + L_j)))$$

Stack - pooling heuristic

$$B_j = (B_{0j} + D_j - S_j)^+ \\ \leq B_{0j} + (D_j - S_j)^+$$

Cost reduces to $T+1$ warehouse

problems: Retailer is $N(h_j, b_j, D_j)$
"as if the warehouse had single supply!"

Weakstar solves $\approx b_0$

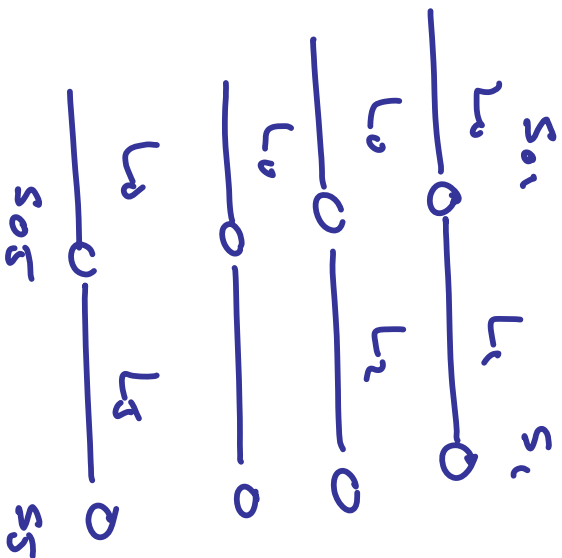
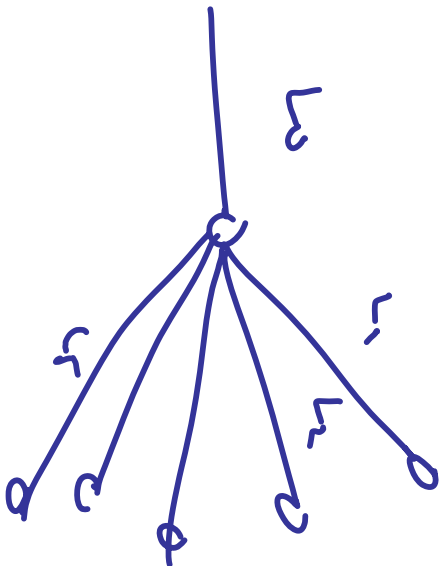
$$\mathcal{N}(A_0, \sum_{j=1}^J b_j \theta_j, D_0)$$

$$S_j = \min \left\{ s \in \mathbb{Z} : P_c(D_j \leq s) \geq \frac{b_j}{b_j + u_j} \right\}$$

$$S_0 = \min \left\{ s \in \mathbb{Z} : P_c(D_0 \leq s) \geq \frac{b_0}{b_0 + u_0} \right\}$$

reducing $J+1$ weakstar problems.

Projectus Horrentis



Solve T
two stage
Gravel capture

Find s_0 such that

$$F(D_0 - s_0)^+ = \sum_{j=1}^n (D_{0j} - s_{0j})^+$$

For each $j \in \{1, \dots, n\}$ find

$$s_j(s_0) = \min\{s \in \mathbb{R} : (b_{0j} + D_j \leq s) \geq \frac{d_j}{b_j + h_j}\}$$

Repeat for all iterations.

$$h_0(F_0) + \sum_{j=1}^n c_j(s_0, s_j(s_0))$$