

# 4106 lec 4

- 1) More on Simulation of Markov chains
- 2)  $n$ -Step transition probabilities  
(Chapman-Kolmogorov Equations)

$$X_{n+1} = f(X_n, U_n), n \geq 0$$

$$f = f(i, u)$$

↑ ↑  
iid Sequence

( $X_0$  is chosen independent of  $\{U_n\}$ )

always forms a MC ; converge

also true :

Every MC  $(X_n)$

can be expressed as a recursion: there exists

a function  $f = f(x, u)$

and an iid sequence  $(U_n)$

such that

$$X_{n+1} = f(X_n, U_n), \quad n \geq 0$$

In fact,  $(U_n)$  can be chosen as  
iid  $\text{Unif}(0,1)$  rvs

$$\text{CDF } P(U \leq x) = x, \quad x \in (0, 1)$$

(cont inuous  
Uniform  
distribution  
over (0,1))

$$\text{density } f(x) = \begin{cases} 1, & x \in (0, 1) \\ 0, & x \notin (0, 1) \end{cases}$$

$$0 \leq a < b \leq 1$$

$$P(a < U \leq b) = b - a$$

Consider any

$$\text{CDF } F(x) = P(X \leq x), \quad x \in \mathbb{R} \text{ of a rv.}$$

If  $F^{-1}(y)$  exists then  $X = F^{-1}(U)$   
has CDF  $F$

Example

$$F(x) = 1 - e^{-\lambda x}, \quad x \geq 0$$

exponential dist.  
at "rate"  $\lambda$

find  $F^{-1}(y)$ ,  $y \in (0, 1)$  ( $X \geq 0$ )

Solve  $y = F(x)$  for  $x$  in terms  
of  $y$

$$y = 1 - e^{-\lambda x}, \quad 1 - y = e^{-\lambda x}, \quad \ln(1 - y) = -\lambda x$$
$$x = -\frac{1}{\lambda} \ln(1 - y) \quad \checkmark$$

$$x = F^{-1}(x) = -\frac{1}{\lambda} \ln(1-x)$$

$$X = \boxed{-\frac{1}{\lambda} \ln(1-U)} \stackrel{\text{dist.}}{=} \boxed{-\frac{1}{\lambda} \ln(U)}$$

because  $1-U \stackrel{\text{dist.}}{=} U \sim \text{unif}(0,1)$

$$P(1-U \leq x) = P(U > 1-x) = 1 - (1-x) = x \quad \checkmark$$

$X = -\frac{1}{\lambda} \ln(U)$  has the  $\exp(\lambda)$  dist.:

$$P\left(-\frac{1}{\lambda} \ln(U) \leq x\right)$$

$$= P\left(\ln(U) > -\lambda x\right)$$

$$= P\left(U > e^{-\lambda x}\right) = 1 - e^{-\lambda x} \quad \checkmark$$

"Inverse transform method"  
for simulating copies  $X$  distributed  
as  $F$

More generally define  
"generalized inverse function"  
of  $F(x)$

$$F^{-1}(y) = \min \{x : F(x) \geq y\}$$

again  $X = F^{-1}(U)$  yields  
a rv with CDF  $F$  (even if  
 $X$  is discrete)



Suppose  $X$  is discrete  
 on  $\{0, 1, 2, \dots\}$

$$p(k) = P(X=k) \quad k \geq 0$$

Generate  $U$

$$\text{Set } X = F^{-1}(U) = \begin{cases} 0 & \text{if } U \leq p(0) \\ 1 & \text{if } p(0) < U \leq p(0) + p(1) \\ \vdots & \vdots \\ k & \text{if } \sum_{j=0}^{k-1} p(j) < U \leq \sum_{j=0}^k p(j) \\ \vdots & \vdots \end{cases}$$

$$p(k) = P(X=k) \checkmark$$

define  $f(i, u) \stackrel{\text{def}}{=} F_i^{-1}(u)$   $U_j, U_i$   
 $i \in \mathcal{I}$ .

$F_i$  = CDF of  $p$   $i$ th row  
of  $p$   $i \in \mathcal{I}$

$F_i^{-1}$  = generalized inverse  
function of  $F_i$

if  $X_n = i_j$  then

$$X_{n+1} = f(i_j, U_n) = F_{i_j}^{-1}(U_n) \checkmark$$

When  $F^{-1}$  does not have  
a closed form,  
then we can't use  
this  $X = F^{-1}(U)$   
method.

Example:  $X$  has a  $N(0,1)$  dist.

$$F(x) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^x e^{-y^2/2} dy. \quad (\text{Normal}) \quad F^{-1}(y) = ?$$

$$R_{n+1} = R_n + \Delta_{n+1}$$

$$U_n \stackrel{\text{def}}{=} \Delta_{n+1}$$

$$f(i, u) \approx i + u$$



# n-step transition probabilities

$$P_{ij}^{(n)} = P(X_n = j \mid X_0 = i) \quad i, j \in \mathcal{S}$$

$$P^{(n)} = \left( P_{ij}^{(n)} \right)_{i, j \in \mathcal{S}}$$

$n$ -step transition matrix

$$P^{(1)} = P$$

$$n \geq 1$$

$$P^{(0)} \stackrel{\text{def}}{=} I$$

identity matrix

Prop:

$$P^{(n)} = P^n = P \times P \times \dots \times P$$

$\underbrace{\hspace{10em}}_{n \text{ times}}$

Proof based on Chapman Kolmogorov Equations:

$$P_{ij}^{(n+m)} = \sum_{k \in \mathcal{A}} P_{ik}^{(n)} \cdot P_{kj}^{(m)}$$

$$\begin{aligned} n &\geq 0 \\ m &\geq 0 \\ i, j &\in \mathcal{A} \end{aligned}$$

Via Markov Property

$$n = m = 1$$

$$P^{(1)} = P$$

$$P^{(2)}_{ij} = P^{(1+1)}_{ij}$$

$$\sum_{k \in \mathcal{A}} P_{ik} P_{kj}$$

Matrix multiplication!

for all  $i, j \in \mathcal{A}$

$$\Rightarrow P^{(2)} = P \times P = P^2$$

$$P^{(3)}_{ij} = P^{(2+1)}_{ij} = \sum_{k \in \mathcal{A}} P^{(2)}_{ik} P_{kj} = \sum_{k \in \mathcal{A}} P^2_{ik} P_{kj} \Rightarrow$$

$$P^{(3)} = P^2 \times P = P^3$$

Complete by induction on  $n$

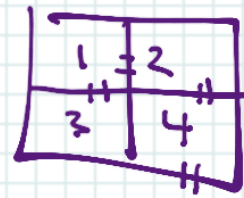
$$P^{(n+1)} = P^{(n)} \cdot P = P^n \cdot P = P^{n+1}$$





# Examples

Rat in open maze



$$S = \{0, 1, 2, 3, 4\}$$

$$P_{1,1}^{(2)} = P(X_2=1 | X_0=1) = \left(\frac{1}{2}\right)^2 + \left(\frac{1}{2}\right)^2 = \frac{1}{2}$$

$P^2$

	0	1	2	3	4
0	0	0	0	0	0
1	0	$\frac{1}{2}$	0	0	0
2	0	0	$\frac{1}{2}$	0	0
3	0	0	0	$\frac{1}{2}$	0
4	0	0	0	0	$\frac{1}{2}$

$$P(X_2=2 | X_0=2)$$

$$P(\text{Rat escapes them no later than } n=7 \mid X_0=1)$$

$$= P_{1,0}^{(7)}$$

$$P_{00} = 1$$

$$P(X_{n+1}=0 \mid X_n=0)$$

$$= P(\text{Rat escapes by time } 7 \mid X_0=1)$$

Gambler's Ruin MC

$$\mathcal{S} = \{0, 1, \dots, N\}$$

$$P_{00} = 1 = P_{NN}$$

$\mathbb{P}(\text{Gambler stops by time } n)$

$$(X_0 = i)$$
$$1 \leq i \leq N-1$$

$$= \boxed{P_{i0}^{(n)} + P_{iN}^{(n)}}$$

Rat in closed maze

$$S = \{1, 2, 3, 4\}$$

$$P_{i,j}^{(n)} = P(X_n = j | X_0 = i)$$

$$1 \leq i, j \leq 4$$

# Weather modeling

$$W_n = \begin{cases} 1 & \text{if it rains on } n^{\text{th}} \text{ day} \\ 0 & \text{if No rain on } n^{\text{th}} \text{ day} \end{cases}$$

$$\mathcal{S} = \{0, 1\}$$

Assume

a MC

$$P = \begin{matrix} & \begin{matrix} 0 & 1 \end{matrix} \\ \begin{matrix} 0 \\ 1 \end{matrix} & \begin{bmatrix} .7 & .3 \\ .4 & .6 \end{bmatrix} \end{matrix}$$

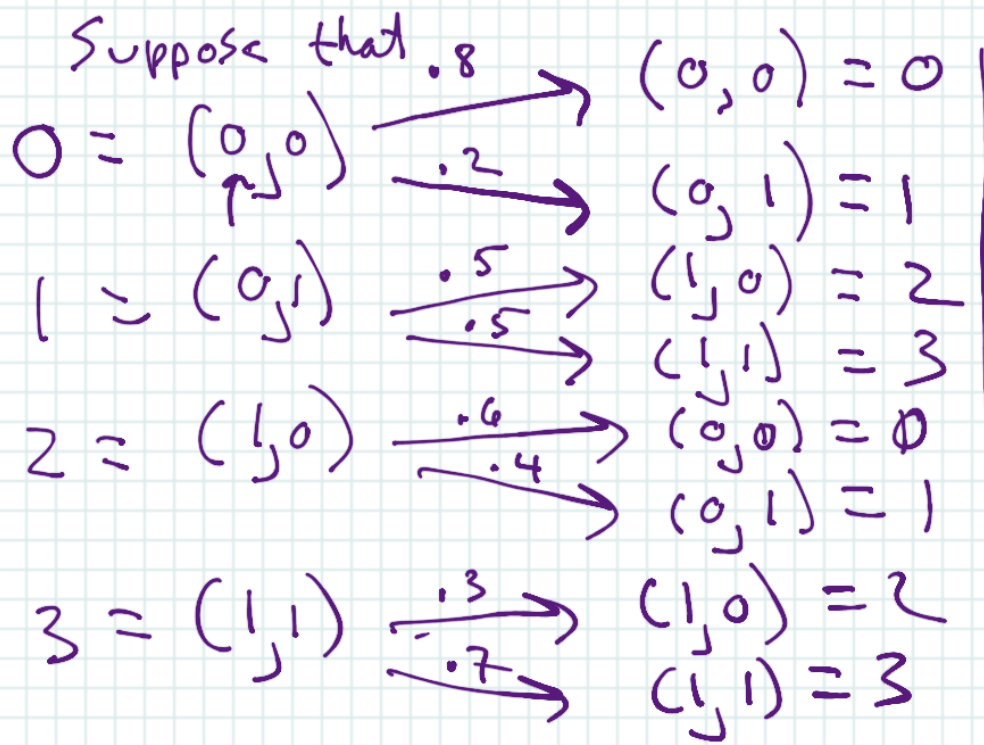
(clearly  
not a  
realistic  
model)

Instead, assume that

$X_n = (W_{n-1}, W_n)$  is a MC

$$\mathcal{F} = \left\{ \begin{array}{c} (0, 0) \\ \downarrow \\ 0 \\ \parallel \\ 0 \end{array} \right\}, \left\{ \begin{array}{c} (0, 1) \\ \parallel \\ 1 \end{array} \right\}, \left\{ \begin{array}{c} (1, 0) \\ \parallel \\ 2 \end{array} \right\}, \left\{ \begin{array}{c} (1, 1) \\ \parallel \\ 3 \end{array} \right\} \right\}$$

more realistic.



$$P = \begin{bmatrix} 0 & 1 & 2 & 3 \\ 0 & .8 & .2 & 0 & 0 \\ 1 & 0 & 0 & .5 & .5 \\ 2 & .6 & .4 & 0 & 0 \\ 3 & 0 & 0 & .3 & .7 \end{bmatrix}$$

$$P(\text{Rain on wed} \mid \text{No rain on Sunday and no rain on Monday}) = ?$$

$$P(X_{n+2} = \begin{matrix} (0,1) = 1 \\ \text{or} \\ (1,1) = 3 \end{matrix} \mid X_n = (0,0)) = P_{0,1}^{(2)} + P_{0,3}^{(2)}$$

$$P^2 = \begin{matrix} & \begin{matrix} 0 & 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} .64 & .16 & .10 & .10 \\ .30 & .20 & .15 & .35 \\ .48 & .12 & .20 & .20 \\ .18 & .12 & .21 & .49 \end{bmatrix} \end{matrix}$$

$$P_{0,1}^{(2)} + P_{0,3}^{(2)} = .16 + .10 = .26$$



Even Better :

$$X_n = (W_{n-2}, W_{n-1}, W_n)$$

$\mathcal{S}$  has 8 states  
||

etc.

# other MC examples

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## Simple Inventory Model

$$\mathcal{S} = \{0, 1, \dots, 10\}$$

headphones (pairs)  
for sale

$$X_0 = 10$$

$X_n$  = inventory at end of  $n^{\text{th}}$  week

$D_n$  = demand on the  $n^{\text{th}}$  week

$(D_n)$  iid poisson ( $\lambda$ )

$$E(D) = \lambda = \text{Var}(D) \quad P(D=j) = e^{-\lambda} \frac{\lambda^j}{j!}, \quad j \geq 0$$

If  $X_n > 0$ , then

$$\underline{X_{n+1}} = (X_n - \textcircled{D_{n+1}})^+$$

$\uparrow = \max\{X_n - D_{n+1}, 0\}$

If  $X_n = 0$ , then restock back to 10

$$X_{n+1} = \left( \underset{\uparrow}{10} - \textcircled{D_{n+1}} \right)^+$$

MC Since it is a recursion

$$P = (P_{ij}) \text{ for}$$

$$S = \{0, 1, 2, 3\}$$

$$X_0 = 3$$

$P =$

	0	1	2	3
0	$P(D \geq 3)$	$P(D=2)$	$P(D=1)$	$P(D=0)$
1	$P(D \geq 1)$	$P(D=0)$	0	0
2	$P(D \geq 2)$	$P(D=1)$	$P(D=0)$	0
3	$P(D=3)$	$P(D=2)$	$P(D=1)$	$P(D=0)$