

IEOR 4106 lec 9

1) Expected number of visits to transient states

2) Binomial lattice model for risky assets.

Introduction to option pricing theory

(See Lecture Notes 6 on course website)

Computing

$$S_{ij} = E \left[\begin{array}{l} \text{number of visits to} \\ \text{state } j \mid X_0 = i \end{array} \right] \quad i, j \in T$$

$$S_{ij} = E \left[\sum_{n=0}^{\infty} I\{X_n = j \mid X_0 = i\} \right]$$
$$= \sum_{n=0}^{\infty} P_{ij}^{(n)}$$

$$\mathcal{J} \text{ finite} = \{1, 2, \dots, N\}$$
$$|\mathcal{J}| = N$$

($1 \leq b < N$ transient states)

$$T = \{1, \dots, b\}$$

($N - b$ recurrent states
left over)

$$S = (S_{ij})_{i,j \in T} \quad b \times b \text{ matrix}$$

$$P_T = (P_{ij})_{i,j \in T} \quad b \times b \text{ matrix}$$

Not a stochastic matrix; some rows will not sum to 1

$$S_{ii} \geq 1 \quad (\text{via } X_0 = i)$$

$$N_i = \text{Total number of visits to } i \mid X_0 = i, \quad i \in T, \quad P(N_i = k) = f_i^{k-1} (1 - f_i)$$
$$f_i = P(\tau_{ii} < \infty) < 1$$

$$N_i = \sum_{n=0}^{\infty} I\{X_n = i \mid X_0 = i\}$$

$$E(N_i) = E(S_{ii})$$

$$\frac{1}{1-f_i}$$

but at this point,
we do not know
how to compute either f_i
or $E(S_{ii})$

Prop 1.1

Let I denote the $b \times b$
identity matrix

$$I = \begin{bmatrix} 1 & & & 0 \\ & \ddots & & \\ & & \ddots & \\ 0 & & & \ddots \\ & & & & \ddots \\ & & & & & 1 \end{bmatrix}$$

Then

$$S = I + P_T S$$

$$\Rightarrow S - P_T S = I$$

$$\Rightarrow \underbrace{(I - P_T)}_M S = I \Rightarrow S = (I - P_T)^{-1}$$

Recall: $\text{Det}(AB)$

$$= \text{Det}(A) \text{Det}(B)$$

$$\Rightarrow \text{Det}\left(\left(\mathbf{I} - \mathbf{P}_T\right)(S)\right) = \text{Det}(\mathbf{I}) = 1 \neq 0$$

$$\stackrel{=}{=} \text{Det}(\mathbf{I} - \mathbf{P}_T) \text{Det}(S)$$

$$\Rightarrow \begin{array}{c} \uparrow \\ \text{both} \end{array} \quad \nearrow \neq 0$$

$$\Rightarrow \left(\mathbf{I} - \mathbf{P}_T\right)^{-1} \text{ exists.}$$

Case 1: $j=i$

$$S_{ii} = 1 + \sum_{k \in T} P_{ik} S_{ki} \quad \leftarrow \text{Markov Property}$$

$j \neq i$

$$S_{ij} = \sum_{k \in T} P_{ik} S_{kj}$$

Matrix form

$$S = I + P_T S$$



$$S = (I - P_T)^{-1} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{pmatrix} 4 & 3 & 3 & 3 \\ 3 & 3.5 & 2.5 & 3 \\ 3 & 2.5 & 3.5 & 3 \\ 2 & 2 & 2 & 3 \end{pmatrix} \end{matrix}$$

$$S_{14} = 3, \quad S_{11} = 4$$

$$S_{11} + S_{12} + S_{13} + S_{14} = 13$$

$$= E(T_{10})$$

= E(escape time

($X_0 = 1$)

as we computed before
using other methods

Gambler's Ruin problem

$$\mathcal{S} = \{0, 1, 2, \dots, N\}$$

$$\mathcal{T} = \{1, 2, \dots, N-1\}$$

$$S_{ij} = E(\# \text{ gambles leading to } \left. \begin{array}{l} \text{fortune } = j \\ \text{before game ends} \end{array} \right| X_0 = i)$$

$$\text{Sum of row } i = E(\# \text{ gambles } \left. \begin{array}{l} \text{until stopping} \\ \text{until stopping} \end{array} \right| X_0 = i)$$

Suppose $N = 10, i = 3$

$$S_{3,6} = E(\# \text{ times gambler } \left. \begin{array}{l} \text{had } \$6 \\ \text{before stopping} \end{array} \right| X_0 = 3)$$

What about computing

$$f_i = \mathbb{P}(\tau_{ii} < \infty), \quad i \in T \quad ?$$

more
generally

$$f_{ij} = \mathbb{P}(X_n \text{ will even visit } j \mid X_0 = i)$$

$j \neq i$

$$S_{ii} = \frac{1}{1 - f_i}$$

$E(N_i)$



$$f_i = \frac{S_{ii} - 1}{S_{ii}}$$

$j \neq i$

$$S_{ij} = f_{ij} S_{ij} + (1 - f_{ij}) \underline{0}$$

$$= f_{ij} S_{ij}$$

$$f_{ij} = \frac{S_{ij}}{S_{ij}}$$

Rat in open maze

$$f_{2,3} = \frac{S_{23}}{S_{33}} = \frac{2.5}{3.5} = \frac{5}{7}$$

($f_{2,4} = 1$)