## Behavioral Economics

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Final Exam

## Friday 18th May

Question 1 (35 pts) We are going to consider preferences over compound lotteries. These are lotteries that give other lotteries as prizes. Let  $\{p_1, q, p_2, r\}$  be the lottery that with probability  $p_1$  gives the lottery q, and with probability  $p_2$  gives the lottery r. For example, consider the lottery in which I flip a coin. If it comes down heads then I roll a die, and if I roll a 1 or 2 (out of 6) I give you \$5 (otherwise nothing). If it comes down heads, I roll a die, and if I get 1,2,3 or 4 I give you \$4 (otherwise nothing). We would write this as

$$\{0.5, q, 0.5, r\}$$

$$q = \frac{1}{3}\$5, \frac{2}{3}\$0$$

$$r = \frac{2}{3}\$4, \frac{1}{3}\$0$$

Call this example a. We will also write  $\{q\}$  for the situation in which the DM receives the lottery  $\{q\}$  for sure. (i.e., in the above example,  $\{q\}$  would be a 100% chance of getting  $\frac{1}{3}$ \$5,  $\frac{2}{3}$ \$0)

Here is one way of calculating the utility of the above lottery (1) calculate the expected utility U(q) and U(r). (2) calculate the utility of the compound lottery as

$$U(\{p_1, q, p_2, r\}) = p_1 U(q) + p_2 U(r)$$

We will call this recursive expected utility approach

1. Assume that the utility of amount u(x) = x. Calculate the recursive expected utility of the lottery of example a

- 2. Show that, for a recursive expected utility maximizer the compound lottery in example *a* is indifferent to receiving the lottery that gives \$5 with probability  $\frac{1}{6}$ , \$4 with probability  $\frac{1}{3}$  and \$0 with probability  $\frac{1}{2}$
- 3. Assume (for simplicity) that the lotteries we consider are over whole dollar amounts between \$0 and \$10. We say that preferences satisfy reduction of compound lotteries if, for ever compound lottery {p<sub>1</sub>, q, p<sub>2</sub>, r}

$$\{p_1, q, p_2, r\}$$
 is indifferent to the lottery  $\{s\}$ 

Where s is the lottery such that, from each  $x \in \{0, 1, ..., 10\}$   $s(x) = p_1q(x) + p_2r(x)$  and s(x), q(x) and r(x) are, respectively, the probability assigned to x by the lotteries s, q and r

Show that the recursive expected utility approach satisfies the reduction of compound lotteries

4. Now consider another way of calculating the utility of a compound lottery. Let  $\pi$  be a cumulative probability weighting function. (1) use  $\pi$  to calculate the non-expected utility  $\bar{U}(q)$  and  $\bar{U}(r)$  of the lotteries q and r (i.e. using the cumulative probability weighting model) (2) calculate the non-expected utility as

$$U(\{p_1, q, p_2, r\}) = \pi(p_1)\bar{U}(q) + (1 - \pi(p_1))\bar{U}(r)$$

if  $\bar{U}(q) \ge \bar{U}(r)$ , or

$$U(\{p_1, q, p_2, r\}) = \pi(p_1)\bar{U}(r) + (1 - \pi(p_1))\bar{U}(q)$$

if  $\bar{U}(r) > \bar{U}(q)$ 

We will call this the recursive non-expected utility approach.

Consider the recursive lottery in example (a). Show that the recursive non-expected utility approach does not necessarily satisfy the reduction of compound lotteries (Make life simple for yourself - assume u(x) = x and remember that you can pick numbers for the probability weighting function, as long as  $\pi(\frac{1}{6}) < \pi(\frac{1}{3}) < \pi(\frac{1}{2}) < \pi(\frac{2}{3})$ )

5. If the probability weighting function is a power function, will the reduction of compound lotteries hold for the recursive lottery in example (a)? (if you get stuck, try it for  $\pi(p) = p^2$ .) Question 2 (45 pts) Consider a decision maker who is choosing over what menu they want to choose from tomorrow. These menus can consist of subsets of three items: apples (a), bourbon (b) and (c) cigarillos. Say that the decision maker has a utility function u such

$$u(a) = 1$$
  
 $u(b) = 2$   
 $u(c) = 3$ 

- Say that the decision maker is standard: i.e. from any menu they will choose the best object in that menu, and so value the menu according to its best option. Let  $\succeq$  represent preferences over menus
  - 1. Calculate the utility of the 7 possible menus that can be constructed from subsets of  $\{a, b, c\}$
  - 2. Notice that we can write the utility of a menu X as

$$U(X) = \max_{x \in X} u(x)$$

Verify that a preference function  $\succeq$  that can be represented by this utility function satisfies the property that, if  $X \succeq Y$ , then  $X \sim X \cup Y$  (If you can show this for general case, at least show its true for the 7 menus you looked at in part 1)

- 3. Does  $\succeq$  satisfy set betweenness (again, do the general case if you can, or if not, then show its true for the 7 menus in part 1)
- 4. Now consider a decision maker who does not know what sort of mood they will be in tomorrow. With a 50% chance they think that they will want to be unhealthy, in which case they will have the utility function u (from section 1 above). with a 50% chance they think that they will wake up wanting to be healthy, in which case they will have the utility function v

$$v(a) = 3$$
  
 $v(b) = 2$   
 $v(c) = 1$ 

They calculate the utility of a menu by calculating the expected utility of that menu: i.e., for a menu containing  $\{a, b\}$ , there is a 50% chance that they will wake up with utility function u. In this case b is better than a, and so they will choose b and get utility u(b). With 50% chance they will wake up with utility function v, in which case ais better than b and, they will choose a and get utility v(a). Thus the utility of this set is  $\{a, b\}$  is given by  $U(\{a, b\}) = 0.5u(b) + 0.5v(a)$ 

Calculate the utility for this decision maker of the 7 possible menus that can be constructed from  $\{a, b, c\}$ 

- Do the preferences over menus of this decision maker satisfy the condition described in
   (2) above?
- 6. Do they satisfy set betweenness?
- 7. Now consider a general description of this type of preferences (sometimes called a preference for flexibility): Let  $\Omega$  be a set of alternatives, and assume that the decision maker has a set of moods M. Each mood occurs with probability p(m), and each mood gives rise to a utility function  $u_m$  over the objects in  $\Omega$ . For any subset X of  $\Omega$ , the utility of that subset is calculated as

$$U(X) = \sum_{m \in M} p(m) \max_{x \in X} u_m(x)$$

where  $\max_{x \in X} u_m(x)$  is the highest utility obtainable in X according to the utility function  $u_m$ 

Show that a decision maker who assesses menus in this way will satisfy the following condition:

$$\begin{array}{rccc} X & \supseteq & Y \\ & \Rightarrow & X & \trianglerighteq \end{array}$$

Y

8. Show that they will also satisfy the following condition

$$\begin{array}{rcl} X & \sim & X \cup Y \\ \end{array}$$
 implies that, for any  $Z & \subset & \Omega \\ & X \cup Z & \sim & X \cup Y \cup Z \end{array}$ 

9. We sometimes describe a decision maker as sophisticated if  $X \cup \{x\} \triangleright X$  if and only if x will be chosen from the menu  $X \cup \{x\}$ . Will the preferences described at the start of the question satisfy this description of sophistication?

- 10. Show that the preferences in section 7 will not satisfy sophistication (i.e., there is a chance that  $X \cup \{x\} \triangleright X$ , but x would not be chosen from the second stage menu). Can you think of a new definition of sophistication that would be satisfied by these preferences?
- Question 3 (20 pts) Consider the following game (sometimes called the Nash Bargaining game). Two players have to share \$10. Each player makes a bid  $b_1$  and  $b_2$ , which can be any number between 0 and 10. If  $b_1 + b_2 \leq 10$ , then each player receives their bid. If  $b_1 + b_2 > 10$  then each player receives zero. These bids are made simultaneously. Assume that utility is linear in money.
  - 1. Show that assuming standard preferences, a pair of strategies  $\{b_1, b_2\}$  is a Nash Equilibrium if  $b_1 + b_2 = 10$ . Are these the only Nash Equilibria of this game? (Remember, a Nash Equilibrium is a pair of strategies  $\{b_1, b_2\}$  such that  $b_1$  is the best that player 1 can do, given  $b_2$ , and  $b_2$  is the best that player 2 can do given  $b_1$ )
  - 2. Imagine that player 1 has standard preferences, and player 2 has inequality averse preferences with  $\alpha > 0$ . Show that there is a threshold for  $\bar{b}$  such that, if  $b_2 < \bar{b}$ , then  $\{b_1, b_2\}$ such that  $b_1 + b_2 = 10$  is not a Nash Equilibrium. Calculate  $\bar{b}$  as a function of  $\alpha$
  - 3. Again imagine that player 1 has standard preferences, and player 2 has inequality averse preferences. Is it always the case that, if  $b_2 > \overline{b}$ , then  $\{b_1, b_2\}$  such that  $b_1 + b_2 = 10$  is a Nash Equilibrium of the game? What if  $\beta > 0.5$ ?