### Preference for Commitment

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- In order to discuss preference for commitment we need to be able to discuss preferences over menus
- Interpretation: choosing a set of alternatives from which you will make a choice at a later date.
- What would be the standard way of assessing a menu of options  $A = \{a_1, a_2, a_3, ...\}$ ?
- Assume that you will choose the best option from the menu at the later date
- Then a menu A is preferred to menu B if the best option in A is better than the best option in B
- i.e.

$$A \succeq B$$
 if and only if  $\max_{a \in A} u(a) \geq \max_{b \in B} u(b)$ 

- For a 'standard' decision maker, more options to choose from is always (weakly) better
- Add alternative a to a choice set A
  - Either a is preferred to all the options already in A
    - a will be chosen from the expanded choice set
    - $\{a\} \cup A$  is better than A
  - Or there is some b in A which is preferred to a
    - a will not be chosen from the expanded choice set
    - $\{a\} \cup A$  is no better, and no worse than A
- DM will always prefer to have a bigger menu to choose from

$$B \subset A$$
$$\Rightarrow A \succ B$$

- This may not be the case if the DM suffers from problems of temptation:
- Classic example: A dieter might prefer to a restaurant with the menu

fish salad

rather than one with the menu

fish burger salad

- Why?
- (At least) two possible reasons
  - Would prefer to not eat the burger, but worries they will succumb to temptation if the burger is available
  - 2 Thinks they will be able to overcome the temptation to eat the burger, but it will be costly to do so

- We are going to discuss a model of menu preferences and choice that captures both these forces
- Based on the classic work of Gul and Pesendorfer [2001]
- Updated (and better explained) by Lipman and Pesendorfer [2013]

- Let C be a compact metric space
- $\Delta(C)$  set of all measures on the Borel  $\sigma$ -algebra of C (i.e. all lotteries)
  - Use lotteries because it means set of choice objects is convex
- ullet Endow  $\Delta(\mathcal{C})$  with topology of weak convergence
- Z all non empty compact subsets of  $\Delta(\mathcal{C})$  (Hausdorff topology)
- Let  $\succeq$  be a preference relation on Z
  - Interpretation: preference over menus from which you will later get to choose
- Let  $\trianglerighteq$  be a preference relation on  $\Delta(C)$ 
  - Interpretation: preferences when asked to choose from a menu

• For  $x, y \in Z$  and  $\alpha \in (0, 1)$  define

$$\alpha x + (1 - \alpha)y$$

$$= \{ p = \alpha q + (1 - \alpha)r | q \in x, r \in y, \}$$

• E.g. if  $x=\{\delta_a\}$ ,  $y=\{\delta_b,\delta_c\}$  the

$$= \left\{ \begin{array}{l} \alpha x + (1 - \alpha)y \\ \alpha a + (1 - \alpha)b \\ \alpha a + (1 - \alpha)c \end{array} \right\}$$

• Mixture of all elements in menu x with all elements in menu y

## Modelling Preference over Menus

- First, we will consider conditions which are necessary and sufficient for the standard model
  - Single utility function
  - Represents ≥ (choice from menus)
  - <u>></u> (choice between menus) represented using largest utility in the set
- Next, consider how to alter these axioms in order to generate the 'Gul Pesendorfer' model
  - Allows for both 'temptation' and 'self control' to be expressed in menu preferences

Axiom 1 (Preference Relations)  $\succeq$ ,  $\trianglerighteq$  are complete preference relations

Axiom 2 (Independence) 
$$x \succeq y$$
 implies  $\alpha x + (1 - \alpha)z \succeq \alpha y + (1 - \alpha)z \ \forall \ x, y, z \in Z$ ,  $\alpha \in (0, 1)$ 

- Notice that this is not the same as 'standard' independence
- Mixing operation is different
- Need to think a bit about how to interpret it

- Interpretation of independence: Standard Independence + Indifference to Timing of Uncertainty
  - Imagine we extended 

    to preferences over lotteries over menus
  - Independence would now say that, if we prefer choosing from x to choosing from y then we prefer choosing from x  $\alpha\%$  of the time (and z  $(1-\alpha)\%$  of the time) to choosing from y  $\alpha\%$  of the time (and z  $(1-\alpha)\%$  of the time)
  - Randomization occurs before choosing at second stage
- Claim: choosing contingent plans in this set up gives rise to the same probability distribution over outcomes as come about from 'Gul Pesendorfer' mixing

Example

$$\frac{1}{2}x + \frac{1}{2}z$$

$$x = \{x_1, x_2\}, \ z = \{z_1, z_2\}$$

• Gul-Pesendorfer mixing: a menu of

$$\left\{ \begin{array}{l} \frac{1}{2}x_1 + \frac{1}{2}z_1 \\ \frac{1}{2}x_2 + \frac{1}{2}z_1 \\ \frac{1}{2}x_1 + \frac{1}{2}z_2 \\ \frac{1}{2}x_2 + \frac{1}{2}z_2 \end{array} \right\}$$

- 'Standard' Mixing: 50% chance of menu x, 50% chance of menu y
  - Contingent plan: choose either  $x_1$  or  $x_2$  from x and either  $y_1$  or  $y_2$  from y
  - Uncertainty decided before second stage choice
  - Set of contingent plans gives rise to same menu of lotteries over outcomes as does GP mixing

- If timing of resolution of uncertainty is not important there is an equivalence between
  - Choosing a contingent plan for a lottery over menus
  - Choosing from a menu of lotteries generated by 'Gul Pesendorfer' mixing
- Thus, 'standard' independence and indifference to timing of uncertainty give rise to GP independence

#### Axiom 3 (Sophistication) $x \cup \{p\} \succ x \Leftrightarrow p \rhd q \ \forall \ q \in x$

- This is the axiom that links together first and second stage choice.
- Whether or not people are sophisticated is going to be an important empirical question
  - Do they understand the choices they will make from a given menu?
  - If not, may underestimate their degree of self control
  - e.g. sign up for gym memberships they do not use
  - or make costly commitments which they subsequently do not stick to.

#### Axiom 4 (Continuity) Three continuity conditions:

- ① (Upper Semi Continuity): The sets  $\{z \in Z | z \succeq x\}$  and  $\{p \in \Delta(C) | p \trianglerighteq q\}$  are closed for all x and q
- **2** (Lower vNM Continuity):  $x \succ y \succ z$  implies  $\alpha x + (1-a)z \succ y$  for some  $\alpha \in (0,1)$
- (Lower Singleton Continuity): The sets  $\{p: \{q\} \succeq \{p\}\}$  are closed for every q

## Standard Model

• The Standard Model of preference over menus

$$U(z) = \max_{p \in z} u(p)$$

for some linear, continuous utility  $u:\Delta(C)\to\mathbb{R}$  such that

- *U* represents ≥
- u represents ⊵

### Standard Model

Equivalent to axioms 1-4 and

$$x \succeq y \Rightarrow x \cup y \sim x$$

- x ≥ y implies that the best alternative in x is weakly better than the best alternative in y
- The best alternative in  $x \cup y$  is the same as the best alternative in x
- Thus  $x \cup y \sim x$
- Note that this implies

$$x \supset y \Rightarrow x \succeq y$$

- Say y ≻ x
  - either  $x/y \succeq y$  in which case

$$x = x/y \cup y \sim x/y \succeq y \succ x$$

• or  $y \succeq x/y$ 

$$x = x/y \cup y \sim y \succ x$$

#### The Gul Pesendorfer Model

Preference over menus given by

$$U(x) = \max_{p \in x} [u(p) + v(p)] - \max_{q \in x} v(q)$$

- *u* : 'long run' utility
- v : 'temptation' utility
- Interpretation:
  - Choose p to maximize u(p) + v(p)
  - Suffer temptation cost v(p) v(q)
- Unlike the standard model, the Gul Pesendorfer model can lead to strict preference for smaller choice sets

$$x \supset y$$
 but  $x \prec y$ 

Case 1: Commitment

Object	и	V
Salad	4	0
Fish	2	1
Burger	1	4

• Which menu would the DM prefer?  $\{s\}$  or  $\{s, b\}$ ?

$$U(\{s\}) = \max_{x \in \{s\}} (u(x) + v(x)) - \max_{y \in \{s\}} v(y)$$
  
= 4 + 0 - 0  
= 4

$$U(\{s,b\}) = \max_{x \in \{s,b\}} (u(x) + v(x)) - \max_{y \in \{s,b\}} v(y)$$
  
= 1 + 4 - 4

Case 1: Commitment

Object	и	V
Salad	4	0
Fish	2	1
Burger	1	4

- Menu  $\{s\}$  preferred to  $\{s.b\}$
- Interpretation: b would be chosen from the latter menu

• 
$$u(b) + v(b) > u(s) + v(s)$$

- But s has higher long run utility
  - u(s) > u(b)
- The DM would rather not have b in their menu, because if it is available they will choose it.

• More generally, consider p, q, such that

$$u(p) > u(q)$$
  
 $u(q) + v(q) > u(p) + v(p)$ 

Then

$$U(\{p\}) = u(p)$$
  
 $U(\{p,q\}) = u(q) + v(q) - v(q) = u(q)$   
 $U(\{q\}\} = u(q)$ 

- Interpretation: give in to temptation and choose q
- 'Weak set betweenness'

$$\{p\} \succ \{p,q\} \sim \{q\}$$

Case 2: Avoid 'Willpower Costs'

Object	и	V
Salad	4	0
Fish	2	1
Burger	1	4

• Which menu would the DM prefer?  $\{s\}$  or  $\{s, f\}$ ?

$$U(\{s\}) = \max_{x \in \{s\}} (u(x) + v(x)) - \max_{y \in \{s\}} v(y)$$

$$= 4 + 0 - 0$$

$$= 4$$

$$U(\{s, f\}) = \max_{x \in \{s, f\}} (u(x) + v(x)) - \max_{y \in \{s, f\}} v(y)$$

$$= 4 + 0 - 1$$

$$= 3$$

Case 2: Avoid 'Willpower Costs'

Object	и	V
Salad	4	0
Fish	2	1
Burger	1	4

- Menu  $\{s\}$  is preferred to menu  $\{s, f\}$
- However, this time, s would be chosen from both menus, as

$$u(s) + v(s) > u(f) + v(f)$$

- The DM still prefers to have f removed from the menu because it is more tempting: v(f) > v(s)
- The DM is able to exert self control if both options are on the menu, but it is costly to do so

Case 2: Avoid 'Willpower Costs'

More generally, consider p, q, such that

$$u(p) > u(q)$$
  
 $v(q) > v(p)$   
 $u(p) + v(p) > u(q) + v(q)$ 

Then

$$U(\{p\}) = u(p)$$
  
 $U(\{p,q\}) = u(p) + v(p) - v(q)$   
 $U(\{q\}\} = u(q)$ 

- Interpretation: fight temptation, but this is costly
- 'Strict set betweenness'

$$\{p\} \succ \{p,q\} \succ \{q\}$$

### Temptation and Self Control

- We say that q tempts p if  $\{p\} \succ \{p, q\}$
- We say that a decision maker exhibits self control at y if there exists x, z such that  $x \cup z = y$  and

$$\{x\} \succ \{y\} \succ \{z\}$$

- {x} ≻ {y} implies there exists something in z which is tempting relative to items in x
- $\{y\} \succ \{z\}$  implies tempting item not chosen
- if it were then

$$\max_{p \in y} u(p) + v(p) = \max_{p \in z} u(p) + v(p) \Rightarrow$$

$$U(y) = \max_{p \in y} (u(p) + v(p)) - \max_{q \in y} v(q)$$

$$\leq \max_{p \in z} (u(p) + v(p)) - \max_{q \in z} v(q)$$

$$= U(z)$$

# Why 'Long Run' and 'Temptation' Utilities?

- So far we have described u as 'long run' utility and v as 'temptation' utility
- Why is this a behaviorally appropriate description?
- u describes choices over singleton menus:

$$U(\{p\}) = u(p) + v(p) - v(p) = u(p)$$

and so describes preferences when the DM is not tempted

# Why 'Long Run' and 'Temptation' Utilities?

- v leads to temptation: q tempts p only if v(q) > v(p)
  - Case 1: u(p) + v(p) > u(q) + v(q)

$$U(\lbrace p\rbrace) > u(\lbrace p, q\rbrace)$$

$$\Rightarrow u(p) > u(p) + v(p) - \max_{r \in \lbrace p, q\rbrace} v(r)$$

$$\Rightarrow \max_{r \in \lbrace p, q\rbrace} v(r) > v(p)$$

$$\Rightarrow v(q) = \max_{r \in \lbrace p, q\rbrace} v(r) > v(p)$$

# Why 'Long Run' and 'Temptation' Utilities?

- v leads to temptation: q tempts p only if v(q) > v(p)
  - Case 2: u(q) + v(q) > u(p) + v(p)

$$\begin{array}{ll} U(\{p\}) &>& u(\{p,q\}) \\ &\Rightarrow & u(p) > u(q) + v(q) - \max_{r \in \{p,q\}} v(r) \\ &\Rightarrow & u(p) + \max_{r \in \{p,q\}} v(r) > u(q) + v(q) \\ &\Rightarrow & \max_{r \in \{p,q\}} v(r) = v(q) > v(p) \end{array}$$

• Last line follows from assumption u(q) + v(q) > u(p) + v(p)

## Limiting Case: No Willpower

- Imagine that differences in v are large relative to differences in u
- In the limit, model reduces to

$$U(x) = \max_{p \in x} u(p)$$
 s.t.  $v(p) \ge v(q) \ \forall \ q \in x$ 

- This is the 'Strolz' model
- Implies no strict set betweenness, and not self control
- $\beta \delta$  model is of this class

• Set Betweenness: for any x, y s.t  $x \succeq y$ 

$$x \succeq x \cup y \succeq y$$

Notice the difference to the 'standard' model

$$x \succeq y \Rightarrow x \cup y \sim x$$

Smaller sets can be strictly preferred

• Set Betweenness: for any x, y s.t  $x \succeq y$ 

$$x \succeq x \cup y \succeq y$$

- Necessity:
  - $x \succeq y$  implies that

$$u(p^{x}) + v(p^{x}) - v(q^{x}) \ge u(p^{y}) + v(p^{y}) - v(q^{y})$$

where

$$p^i = \arg\max_{p \in i} u(p) + v(p)$$

and

$$q^i = rg \max_{q \in i} v(q)$$

• NTS  $x \succeq x \cup y$ 

- Two cases:
- Case 1:  $u(p^x) + v(p^x) \ge u(p^y) + v(p^y)$

$$u(p^{x}) + v(p^{x}) \geq u(p^{y}) + v(p^{y}) \Rightarrow$$

$$u(p^{x}) + v(p^{x}) = u(p^{x \cup y}) + v(p^{x \cup y}) \Rightarrow$$

$$u(p^{x}) + v(p^{x}) - v(q^{x}) \geq u(p^{x \cup y}) + v(p^{x \cup y}) - v(q^{x \cup y})$$

- Case 2:  $u(p^x) + v(p^x) < u(p^y) + v(p^y)$ 
  - implies  $v(q^x) \le v(q^y)$  as x is preferred to y

$$u(p^{y}) + v(p^{y}) = u(p^{x \cup y}) + v(p^{x \cup y}) v(q^{x \cup y}) = v(q^{y}) \Rightarrow u(p^{x \cup y}) + v(p^{x \cup y}) - v(q^{x \cup y}) = u(p^{y}) + v(p^{y}) - v(q^{y}) \leq u(p^{x}) + v(p^{x}) - v(q^{x})$$

#### **Theorem**

 $\succeq$  satisfies Axioms 1, 2, 4 and set betweenness if and only if it has a Strolz representation or a G-P representation

#### **Theorem**

The proper relation  $\succeq$  and  $\trianglerighteq$  satisfy Axioms 1-4 and set betweenness if and only if

- $\succeq$  has a Stroltz representation and  $p \trianglerighteq q$  if and only if v(p) > v(q) or v(p) = v(q) and  $u(p) \trianglerighteq u(q)$
- or  $\succeq$  has a G-P representation and u(p) + v(p) represents  $\trianglerighteq$

## Sketch of Proof that Axioms Imply Representation

- **Lemma 1:** Axioms 1, 2, 4 imply a linear  $U: Z \to \mathbb{R}$  that represents  $\succeq$  and is continuous on singleton sets
  - This is standard, and makes use of the mixture space axioms

## Sketch of Proof that Axioms Imply Representation

• Lemma 2: Show that

$$\begin{array}{rcl} U(x) & = & \displaystyle \max_{p \in x} \min_{q \in x} U(\{p, q\}) \\ & = & \displaystyle \min_{q \in x} \max_{p \in x} U(\{p, q\}) \end{array}$$

- Utility depends only on 'chosen element', and 'most tempting element
- Proof: Let  $\bar{u} = \max_{p \in x} \min_{q \in x} U(\{p, q\}) = U(\{p^*, q^*\})$
- Note that  $U(\{p^*, q\}) \ge U(\{p^*, q^*\}) = \bar{u} \ \forall \ q \in A$
- Set betweenness implies  $\bar{u} \leq U(\cup_{q \in x} \{p^*, q\}) = U(x)$
- Also, for every  $p \in A$ ,  $\exists q_p \in A$  such that  $U(\{p, q_p\}) \leq \bar{u}$
- By set betweenness  $\bar{u} \geq U(\cup_{p \in A} \{p, q_p\}) = U(x)$

# Sketch of Proof that Axioms Imply Representation

• **Lemma 3:** Show that

$$\begin{array}{lcl} U(\{x\}) & > & U(\{x,y\}) > U(\{y\}) \\ U(\{a\}) & > & U(\{a,b\}) > U(\{b\}) \end{array}$$

implies

$$U(\alpha \{x, y\} + (1 - \alpha) \{a, b\})$$
=  $U(\{\alpha x + (1 - \alpha)a\}, \alpha y + (1 - \alpha)b\})$ 

• This comes straight from super independence and the fact that  $\alpha x + (1-\alpha)a$  is the best and  $\alpha y + (1-\alpha)b$  the most tempting element

Define

$$\begin{array}{rcl} u(p) & = & U(\{p\}) \\ v(s;p,q,\delta) & = & \frac{U(\{p,q\}) - U(\{p,(1-\delta)q + \delta s\})}{\delta} \end{array}$$

- u is the long run utility
- v is a measure of how tempting s is relative to p and q (under the assumption p is chosen)

• Lemma 4: Show that, if

$$U(\{\mathit{p}\}) > U(\{\mathit{p}, (1-\delta)\mathit{r} + \delta\mathit{s}\}) > U(\{(1-\delta)\mathit{r} + \delta\mathit{s}\})$$

for all  $s \in \Delta(C)$ , then

- 1  $U(\{p\}) > U(\{p,s\}) > U(s) \Rightarrow v(s;p,q,\delta) = U(\{p,q\}) U(\{p,s\})$
- 2  $v(p; p, q, \delta) = U(\{p, q\}) U(\{p\})$
- Follows from Lemma 3

• Lemma 5: Show that, if

$$U(\{p\}) \geq U(\{p,q\}) \geq U(\{q\})$$

and for some r and  $\delta$ 

$$U(\{\mathit{p}\}) > U(\{\mathit{p}, (1-\delta)\mathit{r} + \delta\mathit{s}\}) > U(\{(1-\delta)\mathit{r} + \delta\mathit{s}\})$$

for all  $s \in \Delta(C)$ , then

$$U(\lbrace p,q\rbrace) = \max_{w \in \lbrace p,q\rbrace} [u(w) + v(w;p,r,\delta)] - \max_{z \in \lbrace p,q\rbrace} [v(z;p,r,\delta)]$$

Proof (assuming)

$$U({p}) > U({p,q}) > U({q})$$

By previous lemma

$$v(q; p, r, \delta) = U(\lbrace p, r \rbrace) - U(\lbrace p, q \rbrace)$$

$$\geq U(\lbrace p, r \rbrace) - U(\lbrace p \rbrace)$$

$$= v(p; p, r, \delta)$$

and so

$$\max_{z \in \{p,q\}} [v(z; p, r, \delta)] = v(q; p, r, \delta)$$

Also

• Also  

$$u(p) + v(p; p, r, \delta) = U(\{p\}) + U(\{p, r\}) - U(\{p\}) = U(\{p, r\})$$
  
 $u(q) + v(q; p, r, \delta) = U(\{q\}) + U(\{p, r\}) - U(\{p, q\})$ 

and so 
$$\max_{w \in \{p,q\}} \left[ u(w) + v(w;p,r,\delta) \right] = u(p) + v(p;p,r,\delta)$$

• This then implies

$$\max_{w \in \{p,q\}} [u(w) + v(w; p, r, \delta)] - \max_{z \in \{p,q\}} [v(z; p, r, \delta)]$$

$$= u(p) + v(p; p, r, \delta) - v(q; p, r, \delta)$$
(1)
$$= U(\{p\}) + U(\{p,r\}) - U(\{p\}) - U(\{p,r\}) + U(\{p,q\})$$

$$= U(\{p,q\})$$
(2)

Finally, pick p, q such that

$$U({p}) > U({p,q}) > U({q})$$

(if such exists) and pick  $\delta$  such that

$$U(\{\mathit{p}\}) > U(\{\mathit{p}, (1-\delta)\mathit{q} + \delta\mathit{s}\}) > U(\{(1-\delta)\mathit{q} + \delta\mathit{s}\})$$

for all s (which we can do by continuity)

- Define v(s) as  $v(s; p, q, \delta)$ , and show that  $v(s; p, q, \delta)$  doesn't depend on the specifics of the last three parameters.
- Lemma 5 therefore gives

$$U(\{p,q\}) = \max_{w \in \{p,q\}} [u(w) + v(w)] - \max_{z \in \{p,q\}} [v(z)]$$

• Lemma 2 then extends this result to an arbitrary set A

## Discussion: Linearity

Imagine

$$\{p\} \succ \{p,q\} \succ \{q\} \succ \{q,r\} \succ \{r\}$$

- DM can resist q for p and resist r for q.
  - Can they resist r for p?
- Under the GP model, the above implies

$$u(p) > u(q) > u(r)$$
  
 $v(r) > v(q) > v(p)$   
 $u(p) + v(p) > u(q) + v(q) > u(r) + v(r)$ 

• Which in turn implies

$$\{p\} \succ \{p,r\} \succ \{r\}$$

- 'Self Control is Linear'
  - See Noor and Takeoka [2010]

#### Discussion: What is Willpower?

- It seems that the following statement is meaningful:
  - Person A has the same long run preferences as person B
  - Person A has the same temptation as person B
  - Person A has more willpower than person B
- Yet this is not possible in the GP model
- Alternative: Masatlioglu, Nakajima and Ozdenoren [2019]

$$\begin{array}{rcl} U(z) & = & \max_{p \in z} u(p) \\ \text{subject to} & \max_{q \in z} v(q) - v(p) & \leq & w \end{array}$$

## Discussion: Strict Set Betweenness and Random Strolz

- Does  $\{p\} \succ \{p, q\} \succ \{q\}$  imply self control?
- Imagine that you are a Strolz guy with u(p) > u(q), but are not sure that you will be tempted
- Half the time

$$v(p) = v(q)$$

half the time

Implies

$$U(\lbrace p \rbrace) = u(p)$$

$$U(\lbrace p, q \rbrace) = \frac{u(p) + u(q)}{2}$$

$$U(\lbrace q \rbrace) = u(q)$$

Strict set betweenness without self control

## Discussion: Optimism

• Say with probability  $\varepsilon$  won't be tempted so

$$\hat{U}(z) = (1 - \varepsilon)U(z) + \varepsilon \max_{p \in z} u(p)$$

- Can lead to violations of set betweenness.
- Let g = gym, j = jog, t = tv

$$\begin{array}{rcl} u(g) & > & u(j) > u(t) \\ v(g) & < & v(j) < v(t) \\ u(j) + v(j) & > & u(t) + v(t) > u(g) + v(g) \end{array}$$

## Discussion: Optimism

• For  $\varepsilon$  small

$$\{t,j\} \succ \{t,g\}$$

as

$$U(\lbrace t,j\rbrace) = u(j) + v(j) - v(t)$$
  
$$U(\lbrace t,g\rbrace) = u(t)$$

but

$$\{t,j,g\} \succ \{t,j\}$$

as with probability  $\varepsilon$  no temptation and will go to the gym

- Consider choice between menus of drinks cocoa or lemonade
- Must choose between menus now, but your choice from those menus will occur on March 1st
- Which would you prefer?

$$\{c\}, \{I\} \text{ or } \{c, I\}$$
?

 Choice of {c, I} over both {c} and {I} is a violation of set betweenness

- X : set of alternatives
- S : set of states
- $\mu \in \Delta(S)$ : probability distribution over states
- $u: X \times S \rightarrow \mathbb{R}$ : utility function
  - u(x, s) utility of alternative x in state s
- Preference uncertainty driven by uncertainty about s

- Let A be a menu of alternatives
- Choice from A will take place after the state is known
- Value of A before the state is known given by

$$U(A) = \sum_{s \in S} \mu(s) \max_{x \in A} u(x, s)$$

U represents choice between menus

• The 'preference uncertainty' model implies a (potentially strict) preference for larger choice sets

$$A \succ B \Rightarrow A \cup B \succ A$$

• Compare to 'standard' model

$$A \succeq B \Rightarrow A \cup B \sim A$$

And Set Betweenness

$$A \succeq B \Rightarrow A \cup B \preceq A$$

 Preference uncertainty can provide a powerful force that works against a preference for commitment

# Amador Angelitos and Wernig [2005]

- Amador Angelitos and Wernig consider the optimal form of commitment in the face of time inconsistency and a need for flexibility
- Find conditions under which a 'minimum savings rule' is optimal
  - Must save a minimum amount s
  - Free to choose any level of consumption that is consistent with this
- More generally, optimal commitment always exhibits 'bunching at the top'

- Two periods with c consumed in the first period and k consumed in the second
- Total resource constraint is y, B(y) is the budget set
- Utility of time 1 self is given by

$$\theta U(c) + \beta W(k)$$

Utility of time 0 self is given by

$$E\left[\theta U(c) + W(k)\right]$$

•  $\theta$  is an (uncontractible) taste shock, unknown at time 0, distributed according to F

#### A Principle Agent Problem

- Assume distribution of types is represented by continuous  $\theta$  on  $\Theta = [\theta_*, \bar{\theta}]$
- For convenience, assume we are choosing  $u(\theta) = U(c(\theta))$  and  $w(\theta) = W(k(\theta))$  directly
- Value of plan for type  $\theta$  is

$$V( heta) = \max_{ heta' \in \Theta} \left[ rac{ heta}{eta} u( heta') + w( heta') 
ight]$$

Assuming truth telling, and by envelope theorem

$$V'(\theta) = \frac{u(\theta)}{\beta}$$

## A Principle Agent Problem

• Integrating  $V'(\theta^*)$  tells us that

$$V(\theta) = \frac{\theta}{\beta}u(\theta) + w(\theta)$$

$$= \int_{\theta_*}^{\theta} \frac{1}{\beta}u(\theta')d\theta' + \frac{\theta_*}{\beta}u(\theta_*) + w(\theta_*)$$

 As is standard in principle agent problems, this condition plus monotonicity are necessary and sufficient for incentive compatibility

#### The Principle's Problem

• Choose  $\{u, w\}$  to maximize

$$\int (\theta u(\theta) + w(\theta)) f(\theta) d(\theta)$$

subject to

$$\frac{\theta}{\beta}u(\theta) + w(\theta)$$

$$= \int_{\theta_*}^{\theta} \frac{1}{\beta}u(\theta')d\theta' + \frac{\theta_*}{\beta}u(\theta_*) + w(\theta_*)$$

$$C(u(\theta)) + K(w(\theta)) \le y$$

$$u(\theta') \ge u(\theta) \text{ for } \theta' \ge \theta$$

#### The Principle's Problem

- Can use the IC constraint to get rid of w
- Objective function becomes

$$\frac{\theta_*}{\beta}u(\theta_*) + w_* + \frac{1}{\beta}\int_{\theta_*}^{\hat{\theta}} (1 - G(\theta))u(\theta)d\theta \tag{3}$$

subject to

$$W(y - Cu(\theta)) + \frac{\theta}{\beta}u(\theta) - \int_{\theta_*}^{\theta} \frac{1}{\beta}u(\theta')d\theta' - \frac{\theta_*}{\beta}u(\theta_*) - w(\theta_*) \ge 0$$

and monotonicity, where

$$G(\theta) = F(\theta) + \theta(1 - \beta)f(\theta)$$

## Bunching at the Top

• It is always optimal to have some bunching at the top

#### **Theorem**

An optimal allocation  $(w, u^*)$  satisfies  $u^*(\theta) = u^*(\theta_p)$  for  $\theta \ge \theta_p$ , where  $\theta_p$  is the lowest value in  $\Theta$  such that

$$\int_{\theta}^{\theta} (1 - G(\theta')) d(\theta') \le 0$$

for  $\theta \geq \theta_p$ 

It is always optimal to have some bunching at the top

#### Theorem

#### Proof.

The contribution of  $\theta \geq \theta_p$  to the objective function is

$$rac{1}{eta}\int_{ heta_{
ho}}^{ar{ heta}}(1-G( heta))u( heta)d heta$$

rewriting 
$$u(\theta) = u(\theta_p) + \int_{\theta_p}^{\theta} u'(\theta) d(\theta)$$
 gives

$$\frac{1}{\beta}u(\theta_{p})\int_{\theta_{p}}^{\bar{\theta}}(1-G(\theta))u(\theta)d\theta+\int_{\theta_{p}}^{\bar{\theta}}\int_{\theta}^{\bar{\theta}}(1-G(\theta''))u'(\theta')d\theta''d\theta'$$

## Minimal Savings Rule

- It is always optimal for all types above a certain threshold consume the same amount
- This does not imply that a minimum savings rule is necessarily optimal
- For that we need one further condition

$$G(\theta) = F(\theta) + \theta(1 - \beta)f(\theta)$$

is increasing for all  $\theta \leq \theta_p$ 

 If (and only if) this condition is satisfied, a simple minimal savings rule is optimal

- So far, we have assumed that a DM is sophisticated
  - They understand their second stage choice
  - Implemented by the axiom  $x \cup \{p\} \succ x \Leftrightarrow p \rhd q \ \forall \ q \in x$
- What about a DM who is not sophisticated?

Example 1: A DM who ignores temptation

Object	и	V
Salad	4	0
Fish	2	1
Burger	1	4

- Assume these preferences represent choices that the DM will make from the menu
- But they believe that their choices will be governed by u
- Such a DM will prefer  $\{s, b\}$  to  $\{s\}$ , but when faced with the choice from  $\{s, b\}$  will choose b
- Such a DM will violate sophistication
  - Never exhibit a preference for commitment

Example 2: A DM who underestimates temptation

Object	и	V	v'
Salad	5	0	0
Fish	2	1	1
Burger	1	9	5

- Assume that a DM has temptation driven by v, but believes that they have temptation driven by v'
- They are offered the chance to buy a 'commitment contract' where they have to pay \$2 if they eat the burger
- Assume that u(2) = 2, v(2) = 2 the u of money is additive with u of consumption and the v of money is additive with the v of consumption
- Let b + c be the burger with the commitment contract

#### Discussion: Sophistication

• Example 2: A DM who underestimates temptation

Object	и	V	v'
Salad	5	0	0
Fish	2	1	1
Burger	1	9	5
B+C	-1	7	3

• The DM will have preferences

$$\{b+c,s\} \succ \{b,s\}$$

as

$$U({b+c,s}) = u(s) + v'(s) - v'(b+c) = 2$$
  
> 1 = u(b) = U({b,s})

• But the DM will actually choose b+c over s at the second stage as

$$u(b+c) + v(b+c) = 6 > 5 = u(s) + v(s)$$

• Example 2: A DM who underestimates temptation

Object	и	V	v'
Salad	5	0	0
Fish	2	1	1
Burger	1	9	5
B+C	-1	7	3

- End up with lower 'long run' utility
- Also a violation of sophistication as

$$\{b+c,s\} \succ \{b+c\}$$

but b + c will be chosen from the former menu

#### Summary

- Menu preferences allow us to formalize a model of preference for commitment
- We argued that this is a sign that people have problems with temptation
  - Temptation: Preference for Commitment

$$A \succeq B \Rightarrow A \cup B \leq A$$

Preference uncertainty: Preference for Flexibility

$$A \succeq B \Rightarrow A \cup B \succeq A$$

• Compare to 'standard' model

$$A \succ B \Rightarrow A \cup B \sim A$$

 Gul and Pesendorfer provide a model which allows for both temptation and self control

$$U(x) = \max_{p \in x} \left[ u(p) + v(p) \right] - \max_{q \in x} v(q)$$

• Characterized by set betweenness:  $x \succeq y \Rightarrow x \succeq x \cup y \succeq y$